The RSA Trapdoor Permutation

Trapdoor functions (TDF)

Def: a trapdoor func. $X \longrightarrow Y$ is a triple of efficient algs. (G, F, F^{-1})

- $G()$: randomized alg. outputs a key pair (pk, sk)
- $F(pk, \cdot)$: det. alg. that defines a function $X \longrightarrow Y$
- $F^{-1}(sk, \cdot)$: defines a function $Y \longrightarrow X$ that inverts $F(pk, \cdot)$

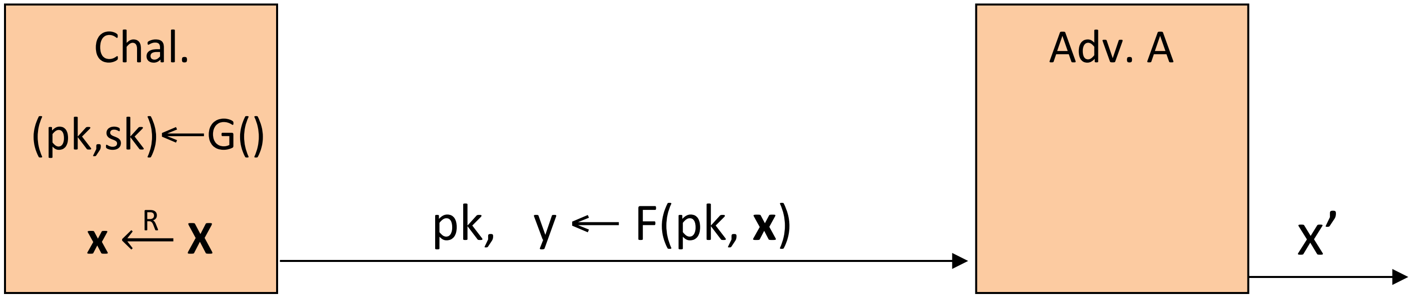
More precisely: $\forall (pk, sk)$ output by G

$$\forall x \in X: \quad F^{-1}(sk, F(pk, x)) = x$$

Secure Trapdoor Functions (TDFs)

(G, F, F^{-1}) is secure if $F(pk, \cdot)$ is a "one-way" function:

can be evaluated, but cannot be inverted without sk



<u>Def</u>: (G, F, F^{-1}) is a secure TDF if for all efficient A:

$$\mathrm{Adv}_{OW}[A,F] = \mathbf{Pr[\ x = x'\ ]} < \text{negligible}$$

Public-key encryption from TDFs

- (G, F, F^{-1}): secure TDF $X \longrightarrow Y$
- (E_s, D_s) : symmetric auth. encryption defined over (K,M,C)
- $H: X \longrightarrow K$ a hash function

We construct a pub-key enc. system (G, E, D):

Key generation G: same as G for TDF

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- $H: X \longrightarrow K$ a hash function

E(pk, m) :

$x \xleftarrow{R} X$, $y \longleftarrow F(pk, x)$

$k \longleftarrow H(x)$, $c \longleftarrow E_s(k, m)$

output (y, c)

D(sk, (y,c)) :

$x \longleftarrow F^{-1}(sk, y)$,

$k \longleftarrow H(x)$, $m \longleftarrow D_s(k, c)$

output m

In pictures:

$$\underbrace{F(pk, x)}_{\text{header}}\ \underbrace{E_s(\ H(x),\ m\ )}_{\text{body}}$$

<u>Security Theorem</u>:

If **(G, F, F^{-1})** is a secure TDF, **(E_s, D_s)** provides auth. enc.

and **H:** $X \longrightarrow K$ is a "random oracle"

then **(G,E,D)** is CCA^{ro} secure.

Incorrect use of a Trapdoor Function (TDF)

Never encrypt by applying F directly to plaintext:

E(pk, m) :	**D(sk, c) :**
output $c \longleftarrow F(pk, m)$	output $F^{-1}(sk, c)$

Problems:

- Deterministic: cannot be semantically secure !!
- Many attacks exist (coming)

The RSA trapdoor permutation

Review: arithmetic mod composites

Let $N = p \cdot q$ where p,q are prime

$$Z_N = \{0,1,2,\ldots,N\text{-}1\} \quad ; \quad (Z_N)^* = \{\text{invertible elements in } Z_N\}$$

<u>Facts</u>: $x \in Z_N$ is invertible $\Leftrightarrow$ $\gcd(x,N) = 1$

– Number of elements in $(Z_N)^*$ is $\varphi(N) = (p-1)(q-1) = N-p-q+1$

<u>Euler's thm</u>: $\boxed{\forall\, x \in (Z_N)^* \ : \ x^{\varphi(N)} = 1}$

The RSA trapdoor permutation

First published: Scientific American, Aug. 1977.

Very widely used:

- SSL/TLS: certificates and key-exchange
- Secure e-mail and file systems

… many others

The RSA trapdoor permutation

G(): choose random primes $p,q \approx 1024$ bits. Set **N=pq**.

choose integers $\mathbf{e}$, $\mathbf{d}$ s.t. $\mathbf{e \cdot d = 1 \pmod{\varphi(N)}}$

output $pk = (N, e)$, $sk = (N, d)$

F(pk, x): $\mathbb{Z}_N^* \to \mathbb{Z}_N^*$; $\mathbf{RSA(x) = x^e}$ (in Z_N)

$\mathbf{F^{-1}(\,sk, y)} = y^d$; $y^d = \mathbf{RSA(x)^d} = x^{ed} = x^{k\varphi(N)+1} = \left(x^{\varphi(N)}\right)^k \cdot \mathbf{x} = x$

The RSA assumption

RSA_e assumption: RSA with exp. e is a one-way permutation

For all efficient algs. A:

$$\Pr\Big[\ A(N,e,\mathbf{y}) = \mathbf{y}^{1/e}\ \Big] < \text{negligible}$$

where $p,q \xleftarrow{R}$ n-bit primes, $N \leftarrow pq$, $y \xleftarrow{R} Z_N^*$

RSA pub-key encryption (ISO std)

(E_s, D_s): symmetric enc. scheme providing auth. encryption.

H: $Z_N \rightarrow K$ where K is key space of (E_s, D_s)

- **G**(): generate RSA params: pk = (N,e), sk = (N,d)
- **E**(pk, m): (1) choose random x in Z_N

  (2) $y \leftarrow RSA(x) = x^e$, $k \leftarrow H(x)$

  (3) output $(y, E_s(k,m))$
- **D**(sk, (y, c)): output $D_s\big(H\big(RSA^{-1}(y)\big), c\big)$

Textbook RSA is insecure

Textbook RSA encryption:

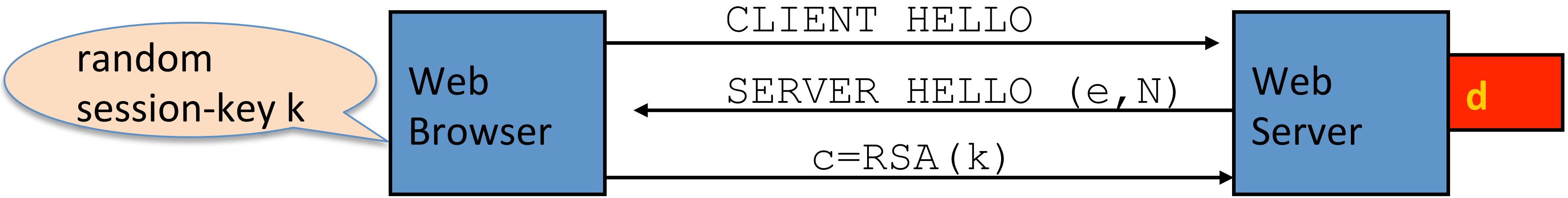
- public key: **(N,e)** Encrypt: $\mathbf{c} \longleftarrow \mathbf{m}^{\mathbf{e}}$ (in Z_N)
- secret key: **(N,d)** Decrypt: $\mathbf{c}^{\mathbf{d}} \longrightarrow \mathbf{m}$

Insecure cryptosystem !!

- Is not semantically secure and many attacks exist

$\Rightarrow$ The RSA trapdoor permutation is not an encryption scheme !

A simple attack on textbook RSA



Suppose k is 64 bits: $k \in \{0,\ldots,2^{64}\}$. Eve sees: $c = k^e$ in Z_N

If $\mathbf{k = k_1 \cdot k_2}$ where $k_1, k_2 < 2^{34}$ (prob. ≈20%) then $\mathbf{c/k_1^e = k_2^e}$ in Z_N

Step 1: build table: $c/1^e, c/2^e, c/3^e, \ldots, c/2^{34e}$. time: 2^{34}

Step 2: for $k_2 = 0,\ldots, 2^{34}$ test if k_2^e is in table. time: 2^{34}

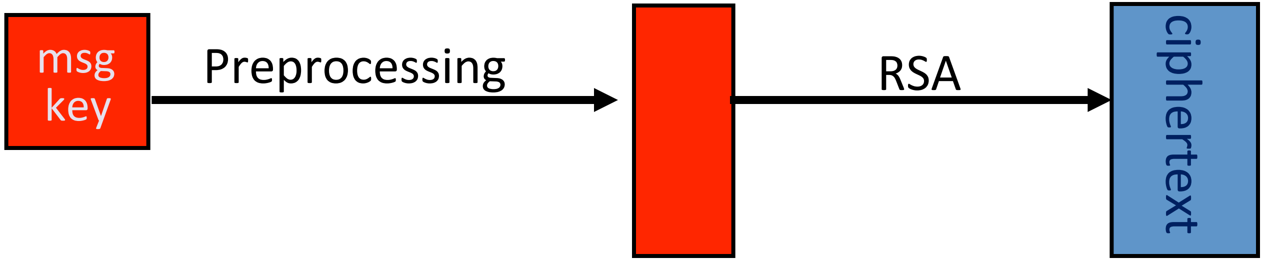
Output matching (k_1, k_2). Total attack time: $\approx 2^{40} << 2^{64}$

RSA in practice

RSA encryption in practice

Never use textbook RSA.

RSA in practice (since ISO standard is not often used) :

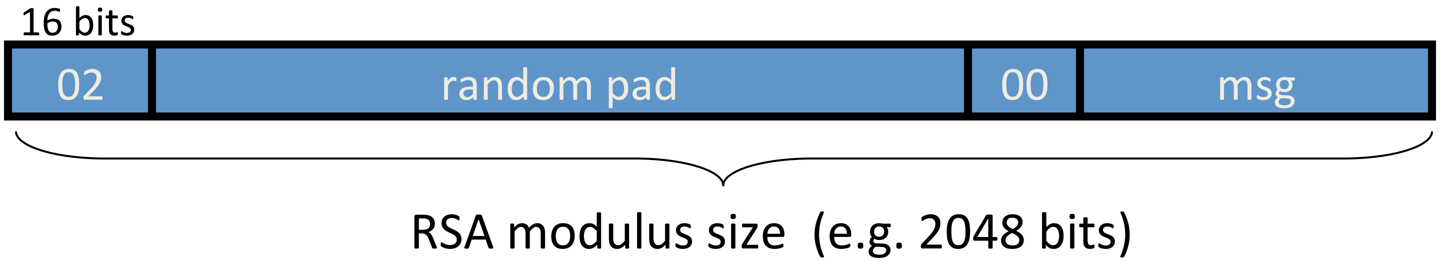


Main questions:

- How should the preprocessing be done?
- Can we argue about security of resulting system?

PKCS1 v1.5

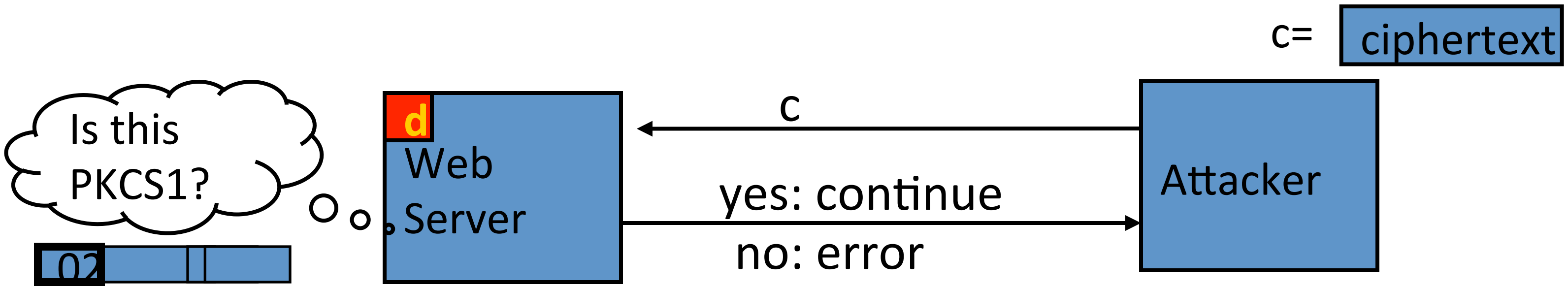
PKCS1 mode 2: (encryption)



- Resulting value is RSA encrypted
- Widely deployed, e.g. in HTTPS

Attack on PKCS1 v1.5 (Bleichenbacher 1998)

PKCS1 used in HTTPS:



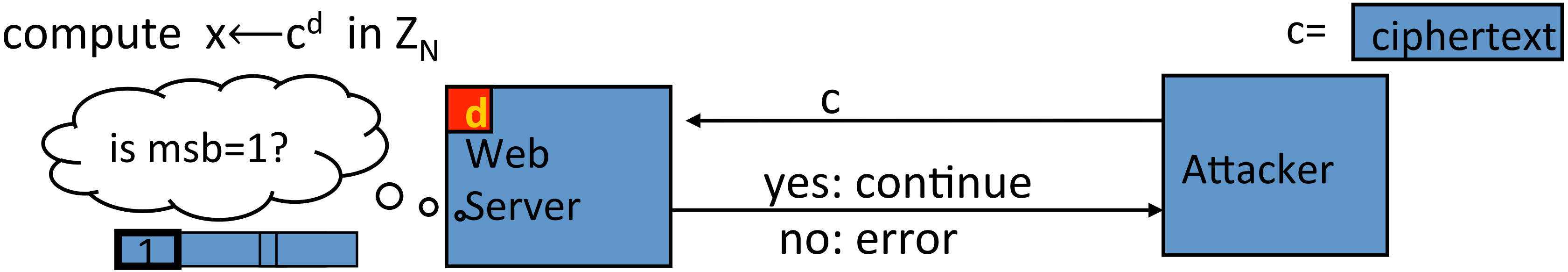
$\Rightarrow$ attacker can test if 16 MSBs of plaintext = '02'

Chosen-ciphertext attack: to decrypt a given ciphertext c do:

– Choose $r \in Z_N$. Compute $c' \leftarrow r^e \cdot c \;=\; (r \cdot \mathrm{PKCS1}(m))^e$

– Send c' to web server and use response

Baby Bleichenbacher



Suppose N is $N = 2^n$ (an invalid RSA modulus). Then:

- Sending c reveals msb(x)
- Sending $\mathbf{2^e \cdot c = (2x)^e}$ in Z_N reveals msb(2x mod N) = $\mathbf{msb_2(x)}$
- Sending $\mathbf{4^e \cdot c = (4x)^e}$ in Z_N reveals msb(4x mod N) = $\mathbf{msb_3(x)}$

… and so on to reveal all of x

HTTPS Defense (RFC 5246)

Attacks discovered by Bleichenbacher and Klima et al. … can be avoided by treating incorrectly formatted message blocks … in a manner indistinguishable from correctly formatted RSA blocks.
In other words:

1. *Generate a string **R** of 46 random bytes*

2. *Decrypt the message to recover the plaintext M*

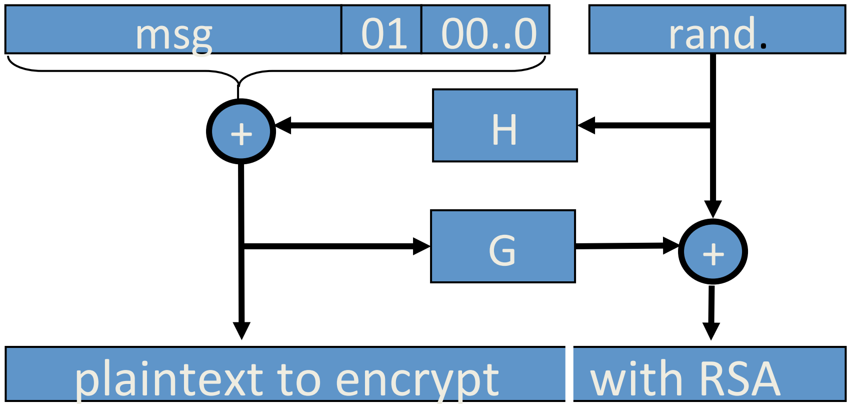
3. *If the PKCS#1 padding is not correct*

   *pre_master_secret = **R***

PKCS1 v2.0: OAEP

New preprocessing function: OAEP [BR94]

check pad
on decryption.
reject CT if invalid.



$\in\{0,1\}^{n-1}$

Thm [FOPS'01] : RSA is a trap-door permutation $\Rightarrow$
RSA-OAEP is CCA secure when H,G are *random oracles*

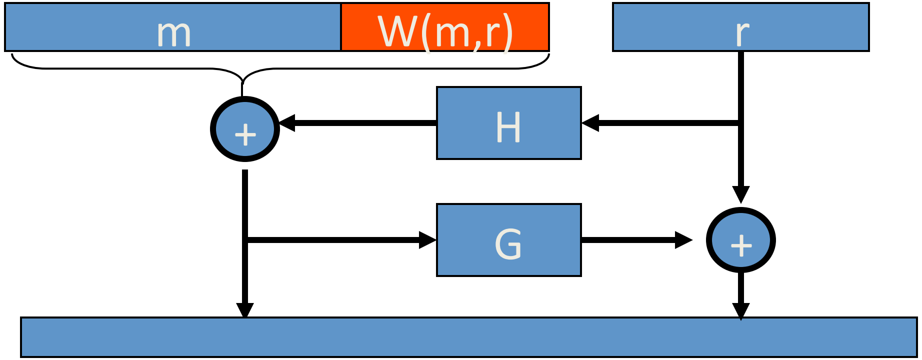
in practice: use SHA-256 for H and G

OAEP Improvements

OAEP+: [Shoup'01]

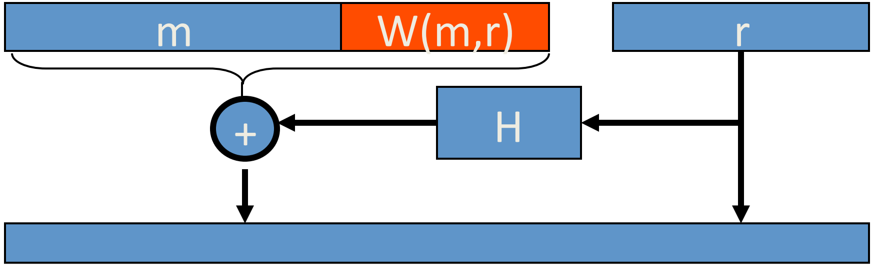
∀ trap-door permutation F
F-OAEP+ is CCA secure when
H,G,W are *random oracles*.

During decryption validate W(m,r) field.



SAEP+: [B'01]

RSA (e=3) is a trap-door perm ⇒
RSA-SAEP+ is CCA secure when
H,W are *random oracle*.



Subtleties in implementing OAEP [M '00]

```
OAEP-decrypt(ct):
        error = 0;
        ........
        if ( RSA^-1(ct) > 2^(n-1) )
              { error =1;  goto exit; }
        ........
        if ( pad(OAEP^-1(RSA^-1(ct))) != "01000" )
              { error = 1;  goto exit; }
```

Problem: timing information leaks type of error

$\Rightarrow$ Attacker can decrypt any ciphertext

Lesson: Don't implement RSA-OAEP yourself !

Is RSA a one-way function?

Is RSA a one-way permutation?

To invert the RSA one-way func. (without d) attacker must compute:

$$x \quad \text{from} \quad c = x^e \pmod{N}.$$

How hard is computing e'th roots modulo N ??

Best known algorithm:

- Step 1: factor N (hard)
- Step 2: compute e'th roots modulo p and q (easy)

Shortcuts?

Must one factor N in order to compute e'th roots?

To prove no shortcut exists show a reduction:

- Efficient algorithm for e'th roots mod N

  $\Rightarrow$ efficient algorithm for factoring N.

- Oldest problem in public key cryptography.

Some evidence no reduction exists: (BV'98)

- "Algebraic" reduction $\Rightarrow$ factoring is easy.

How **not** to improve RSA's performance

To speed up RSA decryption use small private key d ($d \approx 2^{128}$)

$$\boxed{c^d = m \pmod{N}}$$

Wiener'87: if $d < N^{0.25}$ then RSA is insecure.

BD'98: if $d < N^{0.292}$ then RSA is insecure (open: $d < N^{0.5}$)

<u>Insecure:</u> priv. key d can be found from (N,e)

Wiener's attack

Recall: $e \cdot d = 1 \pmod{\varphi(N)} \Rightarrow \exists\, k \in Z: \quad e \cdot d = k \cdot \varphi(N) + 1$

$$\left| \frac{e}{\varphi(N)} - \frac{k}{d} \right| = \frac{1}{d \cdot \varphi(N)} \leq \frac{1}{\sqrt{N}}$$

Wiener's attack

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$$\left| \frac{e}{\varphi(N)} - \frac{k}{d} \right| = \frac{1}{d\cdot\varphi(N)} \le \frac{1}{\sqrt{N}}$$

$\varphi(N) = N\text{-}p\text{-}q\text{+}1 \Rightarrow |N - \varphi(N)| \le p+q \le 3\sqrt{N}$

Wiener's attack

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$$\left| \frac{e}{\varphi(N)} - \frac{k}{d} \right| = \frac{1}{d \cdot \varphi(N)} \leq \frac{1}{\sqrt{N}}$$

$\varphi(N) = N - p - q + 1 \quad \Rightarrow \quad |N - \varphi(N)| \leq p + q \leq 3\sqrt{N}$

$d \leq N^{0.25}/3 \quad \Rightarrow \quad \left| \frac{e}{N} - \frac{k}{d} \right| \leq \left| \frac{e}{N} - \frac{e}{\varphi(N)} \right| + \overbrace{\left| \frac{e}{\varphi(N)} - \frac{k}{d} \right|}^{\leq 1/\sqrt{N}} \leq \frac{1}{2d^2}$

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$\varphi(N) = N - p - q + 1 \;\Rightarrow\; |N - \varphi(N)| \le p+q \le 3\sqrt{N}$

$d \le N^{0.25}/3 \;\Rightarrow\;$

$$\left|\frac{e}{N} - \frac{k}{d}\right| \le \underbrace{\left|\frac{e}{N} - \frac{e}{\varphi(N)}\right|}_{\le \frac{3\sqrt{N}}{N}\cdot\frac{e}{\varphi(N)} \le \frac{3}{\sqrt{N}} \le \frac{1}{2d^2} - \frac{1}{\sqrt{N}}} + \overbrace{\left|\frac{e}{\varphi(N)} - \frac{k}{d}\right|}^{\le 1/\sqrt{N}} \le \frac{1}{2d^2}$$

$$\frac{1}{2d^2} - \frac{1}{\sqrt{N}} \ge \frac{3}{\sqrt{N}}$$

$\Rightarrow$ Continued fraction expansion of e/N gives k/d.

$e\cdot d = 1 \pmod{k} \;\Rightarrow\; \gcd(d,k)=1 \;\Rightarrow\;$ can find d from k/d

RSA With Low public exponent

To speed up RSA encryption use a small e: $c = m^e \pmod{N}$

- Minimum value: **$e=3$** $(\gcd(e, \varphi(N)) = 1)$
- Recommended value: **$e=65537=2^{16}+1$**

Encryption: 17 multiplications

Asymmetry of RSA: fast enc. / slow dec.

- ElGamal: approx. same time for both.

Key lengths

Security of public key system should be comparable to security of symmetric cipher:

Cipher key-size	RSA Modulus size
80 bits	1024 bits
128 bits	3072 bits
256 bits (AES)	**15360** bits

Implementation attacks

Timing attack: [Kocher et al. 1997] , [BB’04]

The time it takes to compute $c^d \pmod N$ can expose d

Power attack: [Kocher et al. 1999)

The power consumption of a smartcard while it is computing $c^d \pmod N$ can expose d.

Faults attack: [BDL’97]

A computer error during $c^d \pmod N$ can expose d.

A common defense:: check output. 10% slowdown.

An Example Fault Attack on RSA (CRT)

A common implementation of RSA decryption: $x = c^d$ in Z_N

$$\left.\begin{array}{ll} \text{decrypt mod p:} & x_p = c^d \text{ in } Z_p \\ \text{decrypt mod q:} & x_q = c^d \text{ in } Z_q \end{array}\right\} \text{combine to get } x = c^d \text{ in } Z_N$$

Suppose error occurs when computing x_q, but no error in x_p

Then: output is x' where $x' = c^d$ in Z_p but $x' \neq c^d$ in Z_q

$\Rightarrow$ $(x')^e = c$ in Z_p but $(x')^e \neq c$ in Z_q $\Rightarrow$ $\gcd\left((x')^e - c , N\right) = p$

RSA Key Generation Trouble [Heninger et al./Lenstra et al.]

OpenSSL RSA key generation (abstract):

```
prng.seed(seed)
p = prng.generate_random_prime()
prng.add_randomness(bits)
q = prng.generate_random_prime()
N = p*q
```

Suppose poor entropy at startup:

- Same p will be generated by multiple devices, but different q
- N_1 , N_2 : RSA keys from different devices $\Rightarrow$ $\gcd(N_1,N_2) = p$

RSA Key Generation Trouble [Heninger et al./Lenstra et al.]

Experiment: factors 0.4% of public HTTPS keys !!

Lesson:

- Make sure random number generator is properly seeded when generating keys

THE END