

Control Hijacking Attacks

Buffer overflows and
format string bugs

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Buffer overflows

- ◆ Extremely common bug.
 - First major exploit: 1988 Internet Worm, fingerd.
- ◆ 15 years later: ≈ 50% of all CERT advisories:
 - 1998: 9 out of 13
 - 2001: 14 out of 37
 - 2003: 13 out of 28
- ◆ Often lead to total compromise of host.
- ◆ Developing buffer overflow attacks:
 - Locate buffer overflow within an application.
 - Design an exploit.

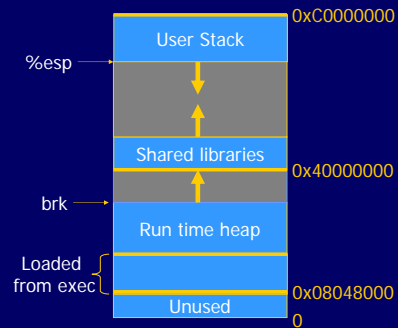
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What is needed

- ◆ Understanding C functions and the stack.
- ◆ Some familiarity with machine code.
- ◆ Know how systems calls are made.
- ◆ The `exec()` system call.
- ◆ Attacker needs to know which CPU and OS are running on the target machine.
 - Our examples are for x86 running Linux.
 - Details vary slightly between CPUs and OSs:
 - Little endian vs. big endian (x86 vs. Motorola)
 - Stack Frame structure (Linux vs. Windows)
 - Stack growth direction.

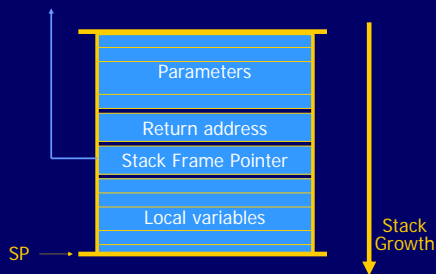
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Linux process memory layout



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Stack Frame



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What are buffer overflows?

- ◆ Suppose a web server contains a function:


```
void func(char *str) {
    char buf[128];
    strcpy(buf, str);
    do-something(buf);
}
```
- ◆ When the function is invoked the stack looks like:

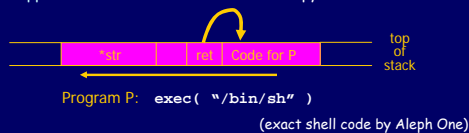
The diagram shows the stack layout for the `func` function call. The stack grows downwards. From top to bottom, the stack contains: `buf`, `str`, `ret-addr`, and `str`. The `top of stack` is indicated by a horizontal line above the `buf` segment.
- ◆ What if `*str` is 136 bytes long? After `strcpy`:

The diagram shows the stack layout after a buffer overflow. The `buf` segment is now 136 bytes long, which has overwritten the `ret-addr` and the second `str` segment. The `top of stack` is indicated by a horizontal line above the `buf` segment.

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Basic stack exploit

- ◆ Main problem: no range checking in strcpy().
- ◆ Suppose *str is such that after strcpy stack looks like:



- ◆ When func() exits, the user will be given a shell !!
- ◆ Note: attack code runs *in stack*.
- ◆ To determine ret guess position of stack when func() is called.

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Some unsafe C lib functions

```
strcpy (char *dest, const char *src)
strcat (char *dest, const char *src)
gets (char *s)
scanf (const char *format, ...)
printf (const char *format, ...)
```

⋮

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Exploiting buffer overflows

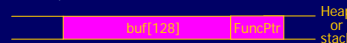
- ◆ Suppose web server calls func() with given URL.
 - Attacker sends a 200 byte URL. Gets shell on web server.
- ◆ Some complications:
 - Program P should not contain the '\0' character.
 - Overflow should not crash program before func() exists.

- ◆ Sample remote buffer overflows of this type:
 - (Old) overflow in MIME type field in MS Outlook.
 - (Old) overflow in Symantec Virus Detection
 - Set test = CreateObject("Symantec.SymVAFileQuery.1")
 - test.GetPrivateProfileString "file", [long string]

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Control hijacking opportunities

- ◆ Stack smashing attack:
 - Override return address in stack activation record by overflowing a local buffer variable.
- ◆ Function pointers: (e.g. PHP 4.0.2, MS MediaPlayer Bitmaps)
 - Overflowing buf will override function pointer.
- ◆ Longjmp buffers: longjmp(pos) (e.g. Perl 5.003)
 - Overflowing buf next to pos overrides value of pos.



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Other types of overflow attacks

- ◆ Integer overflows: (e.g. MS DirectX MIDI Lib) Phrack60


```
void func(int a, char v) {
    char buf[128];
    init(buf);
    buf[3*a+1] = v;
}
```

 - Problem: $3*a+1$ can point to 'ret-addr' on stack.
- ◆ Double free: double free space on heap.
 - Can cause memory mgr to write data to specific locations.
 - Examples: CVS server

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Finding buffer overflows

- ◆ To find overflow:
 - Run web server on local machine.
 - Issue requests with long tags. All long tags end with "\$\$\$\$\$\$".
 - If web server crashes, search core dump for "\$\$\$\$\$\$" to find overflow location.
- ◆ Some automated tools exist. (e.g. eEye Retina).
- ◆ Then use disassemblers and debuggers (e.g. IDA-Pro) to construct exploit.

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Preventing overflow attacks

- ◆ **Main problem:**
 - strcpy(), strcat(), sprintf() have no range checking.
 - "Safe" versions strncpy(), strncat() are misleading
 - strncpy() may leave buffer unterminated.
 - strncpy(), strncat() encourage off by 1 bugs.
- ◆ **Defenses:**
 - Type safe languages (Java, ML). Legacy code?
 - Mark stack as non-execute. Random stack location.
 - Static source code analysis.
 - Run time checking: StackGuard, Libsafe, SafeC, (Purify).
 - Many more ... (covered later in course)

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Marking stack as non-execute

- ◆ Basic stack exploit can be prevented by marking stack segment as non-executable.
 - NX-bit on AMD Athlon 64, XD-bit on Intel P4 "Prescott".
 - NX bit in every Page Table Entry (PTE)
 - Support in SP2. Code patches exist for Linux, Solaris.
- ◆ **Limitations:**
 - Does not defend against 'return-to-libc' exploit.
 - Overflow sets ret-addr to address of libc function.
 - Does not block more general overflow exploits:
 - Overflow on heap: overflow buffer next to func pointer.
 - Some apps need executable stack (e.g. LISP interpreters).

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Static source code analysis

- ◆ Statically check source to detect buffer overflows.
 - Several consulting companies.
- ◆ Can we automate the review process?
- ◆ Several tools exist:
 - Coverity (Engler et al.): Test trust inconsistency.
 - Microsoft program analysis group:
 - PREFIX: looks for fixed set of bugs (e.g. null ptr ref)
 - PREFIX: local analysis to find idioms for prog errors.
 - Berkeley: Wagner, et al. Test constraint violations.
- ◆ Find lots of bugs, but not all.

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Run time checking: StackGuard

- ◆ Many many run-time checking techniques ...
 - Here, only discuss methods relevant to overflow protection.
- ◆ **Solutions 1: StackGuard (WireX)**
 - Run time tests for stack integrity.
 - Embed "canaries" in stack frames and verify their integrity prior to function return.



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Canary Types

- ◆ **Random canary:**
 - Choose random string at program startup.
 - Insert canary string into every stack frame.
 - Verify canary before returning from function.
 - To corrupt random canary, attacker must learn current random string.
- ◆ **Terminator canary:**
 - Canary = 0, newline, linefeed, EOF
 - String functions will not copy beyond terminator.
 - Hence, attacker cannot use string functions to corrupt stack.

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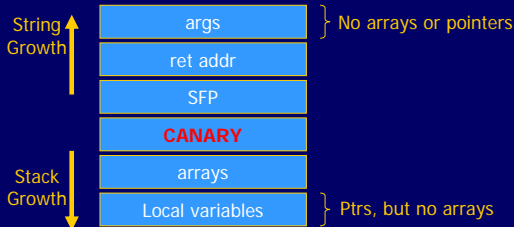
StackGuard (Cont.)

- ◆ StackGuard implemented as a GCC patch.
 - Program must be recompiled.
- ◆ Minimal performance effects: 8% for Apache.
- ◆ Newer version: PointGuard.
 - Protects function pointers and setjmp buffers by placing canaries next to them.
 - More noticeable performance effects.
- ◆ **Note:** Canaries don't offer fullproof protection.
 - Some stack smashing attacks can leave canaries untouched.

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StackGuard variants - ProPolice

- ◆ ProPolice (IBM) - gcc 3.4.1. (-fstack-protector)
 - Rearrange stack layout to prevent ptr overflow.



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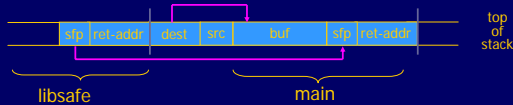
Windows XP SP2 /GS

- ◆ Non executable stack.
- ◆ Compiler /GS option:
 - Combination of ProPolice and Random canary.
 - Triggers UnHandledException in case of Canary mismatch to shutdown process.
- ◆ Litchfield vulnerability report.
 - Overflow overwrites exception handler.
 - Redirects exception to attack code.

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Run time checking: Libsafe

- ◆ Solutions 2: Libsafe (Avaya Labs)
 - Dynamically loaded library.
 - Intercepts calls to strcpy (dest, src)
 - Validates sufficient space in current stack frame: $|frame_pointer - dest| > strlen(src)$
 - If so, does strcpy.
 - Otherwise, terminates application.



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More methods ...

- ◆ StackShield
 - At function prologue, copy return address RET and SFP to "safe" location (beginning of data segment)
 - Upon return, check that RET and SFP is equal to copy.
 - Implemented as assembler file processor (GCC)
- ◆ Randomization:
 - PaX ASLR: Randomize location of libc.
 - Attacker cannot jump directly to exec function.
 - Instruction Set Randomization (ISR)
 - Attacker cannot execute its own code.

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Format string bugs

Format string problem

```
int func(char *user) {
    fprintf(stdout, user);
}
```

- Problem: what if user = "%s%s%s%s%s%s%s" ??
- Most likely program will crash: DoS.
 - If not, program will print memory contents. Privacy?
 - Full exploit using user = "%n"

Correct form:

```
int func(char *user) {
    fprintf(stdout, "%s", user);
}
```

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History

- ◆ Danger discovered in June 2000.
- ◆ Examples:
 - wu-ftpd 2.* : remote root.
 - Linux rpc.statd: remote root
 - IRIX telnetd: remote root
 - BSD chpass: local root



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Vulnerable functions

Any function using a format string.

Printing:

```
printf, fprintf, sprintf, ...  
vprintf, vfprintf, vsprintf, ...
```

Logging:

```
syslog, err, warn
```

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Exploit

- ◆ Dumping arbitrary memory:
 - Walk up stack until desired pointer is found.
 - `printf("%08x.%08x.%08x.%08x|%s|")`
- ◆ Writing to arbitrary memory:
 - `printf("hello %n", &temp) -- writes '6' into temp.`
 - `printf("%08x.%08x.%08x.%08x.%n")`

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Overflow using format string

```
char errormsg[512], outbuf[512];  
sprintf( errormsg, "Illegal command: %400s", user);  
...  
sprintf( outbuf, errormsg );
```

- ◆ What if `user = "%500d <nops> <shellcode>"`
 - Bypass "%400s" limitation.
 - Will overflow outbuf.

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