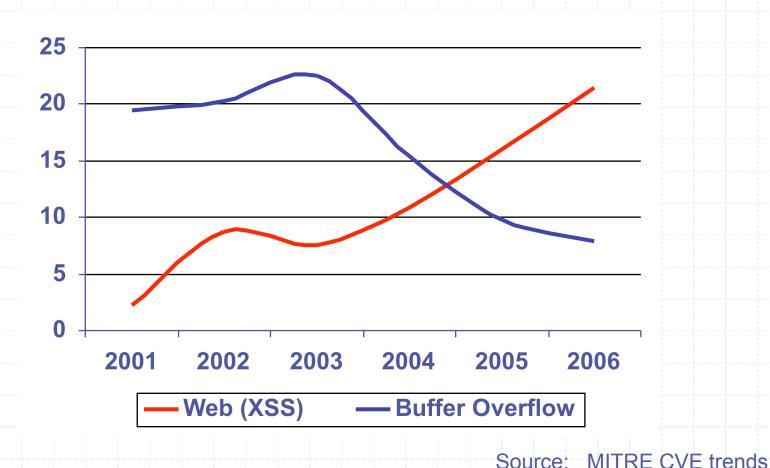
Basic web security model Elie Bursztein CS155

Vulnerability Stats: web is "winning"

Majority of vulnerabilities now found in web software

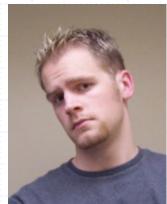


Web security: two sides

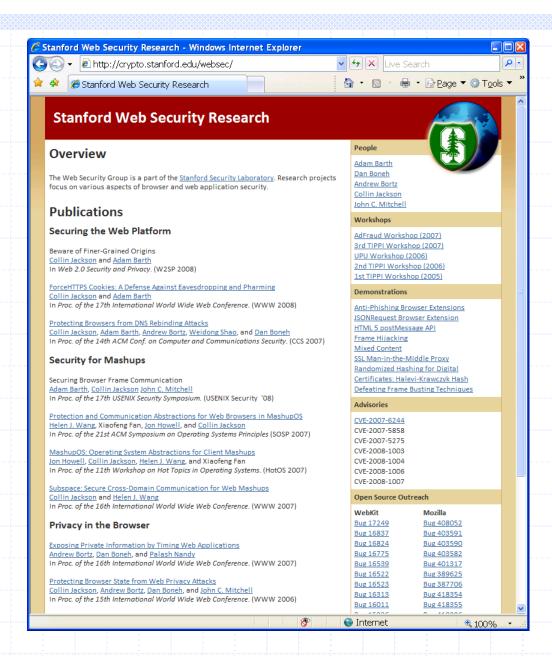
- Web browser: (client side)
 - Attacks target browser security weaknesses
 - Result in:
 - Malware installation (keyloggers, bot-nets)
 - Document theft from corporate network
 - Loss of private data
- Web application code: (server side)
 - Runs at web site: banks, e-merchants, blogs
 - Written in PHP, ASP, JSP, Ruby, ...
 - Many potential bugs: XSS, XSRF, SQL injection
 - Attacks lead to stolen CC#, defaced sites.

Credits





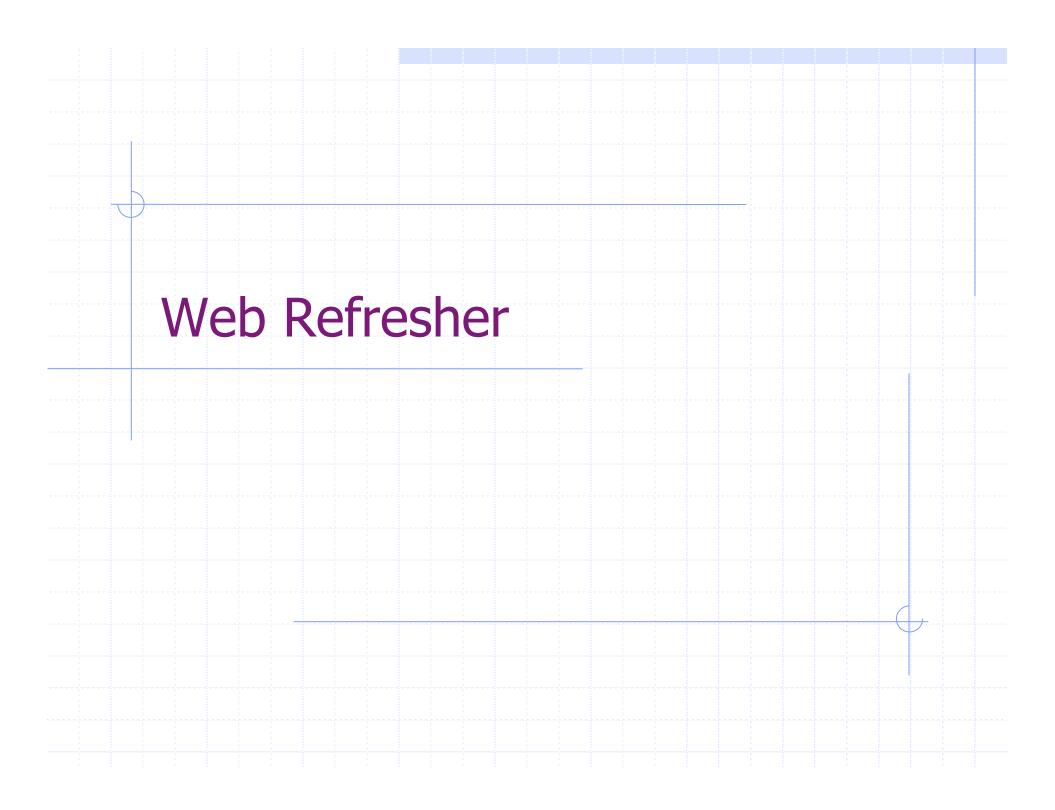
Adam Barth, Collin Jackson, John Mitchell, Dan Boneh and the entire websec team



http://crypto.stanford.edu/websec

Outline

- Web Refresher:
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- Browser security design



HTTP protocol

- ◆ HTTP is
 - widely used
 - Simple
 - Stateless
 - Unencrypted

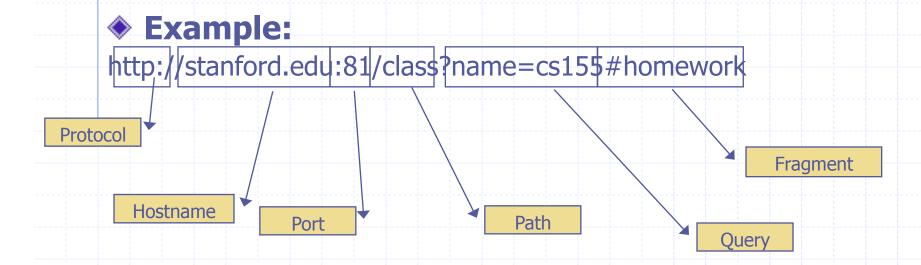






URLs

Global identifiers of network-retrievable documents



- Special characters are encoded as hex:
 - %0A = newline
 - %20 or + = space, %2B = + (special exception)

HTTP Request

Method File **HTTP** version Headers GET /index.html HTTP/1.1 Accept: image/gif, image/x-bitmap, image/jpeg, Accept-Language: en Connection: Keep-Alive User-Agent: Mozilla/1.22 (compatible; MSIE 2.0; Windows 95) Host: www.example.com Referer: http://www.google.com?g=dingbats Blank line Data - none for GET

GET: no side effect.

POST: possible side effect.

HTTP Response

HTTP version Status code

Reason phrase

Headers

HTTP/1.0 200 OK

Date: Sun, 21 Apr 1996 02:20:42 GMT

Server: Microsoft-Internet-Information-Server/5.0

Connection: keep-alive Content-Type: text/html

Last-Modified: Thu, 18 Apr 1996 17:39:05 GMT

Set-Cookie: ...

Content-Length: 2543

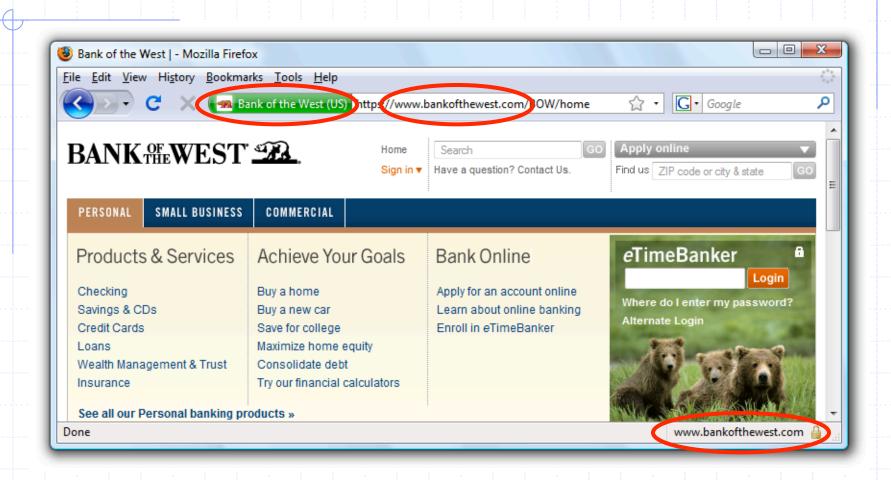
<HTML> Some data... blah, blah, blah </HTML>

Cookies

Data

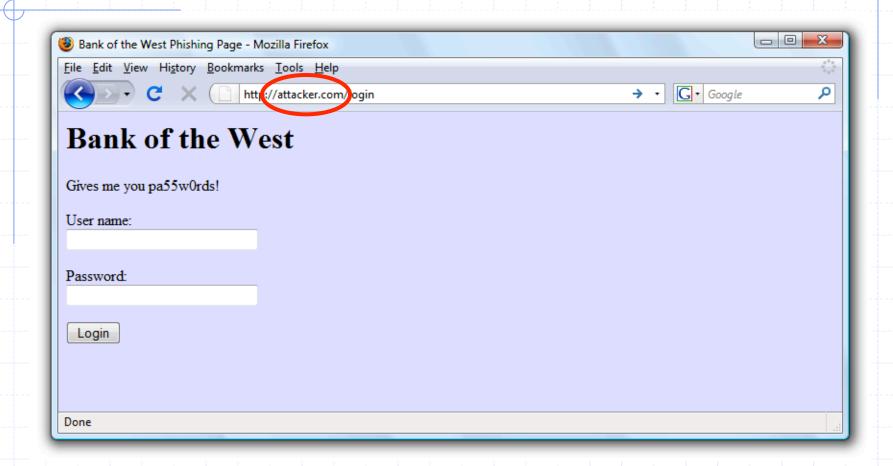


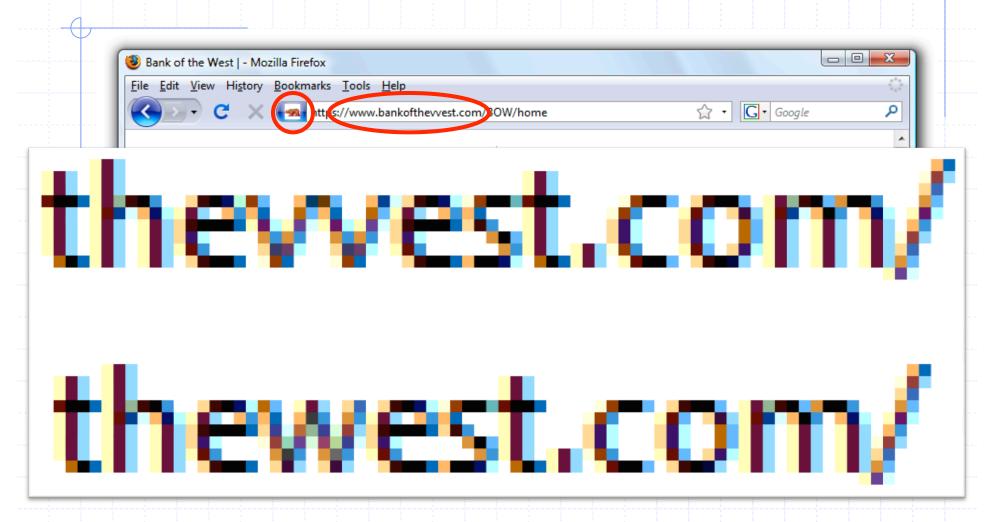
When is it safe to type my password?

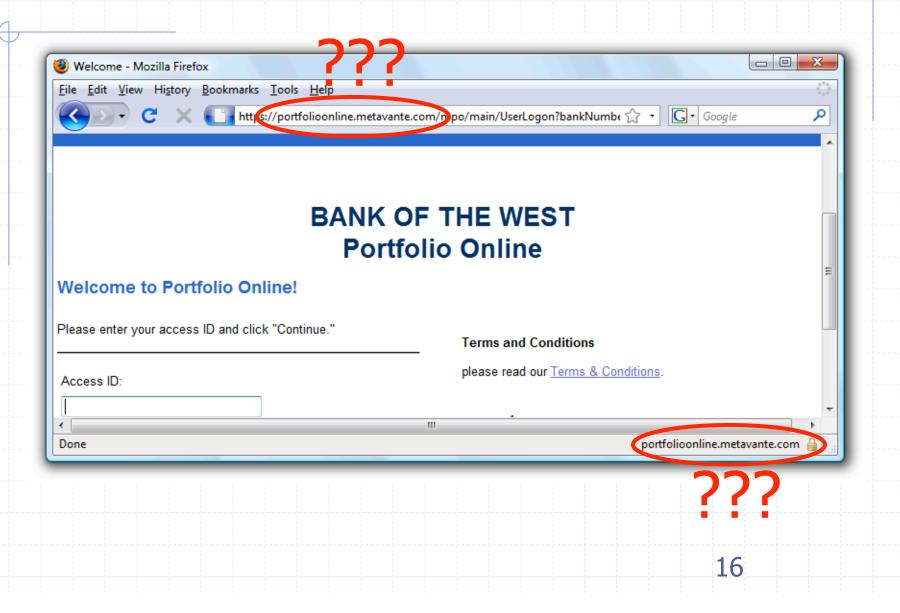


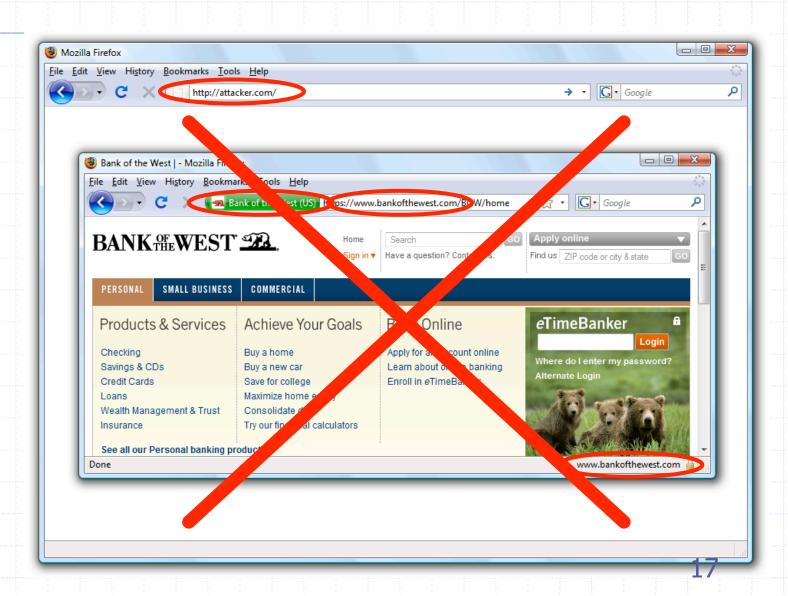
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How does the browser isolate different sites?

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Policy Goals

Safe to visit an evil web site



- Safe to visit two pages at the same time
 - Address bar distinguishes them
- Allow safe delegation





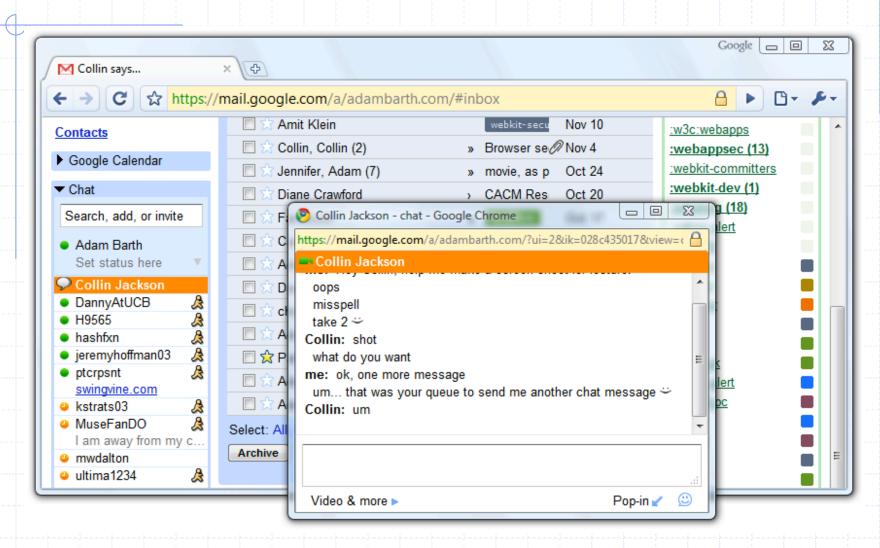
Components of browser security policy

- Frame to Frame relationships
 - canScript(A,B)
 - Can Frame A execute a script that reads or writes DOM elements of Frame B?
 - canNavigate(A,B)
 - Can Frame A change the origin of content for Frame B?
- Frame to cookie relationships
 - readCookie(A,S), writeCookie(A,S)
 - Can Frame A read/write cookies from origin S?
- SecurityIndicator (W) [ssl lock icon]
 - Is the security indicator displayed for window W?

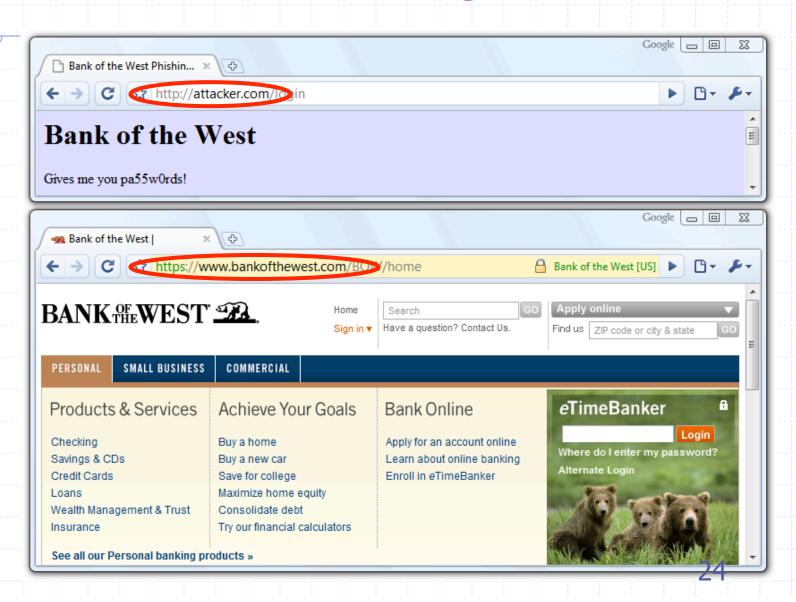
Popup windows

- With hyperlinks
 click here
- With JavaScript
 mywin = window.open("http://www.b.com", "foo", "width=10,height=10")
 - Navigating named window re-uses existing one
 - Can access properties of remote window: mywin.document.body mywin.location = "http://www.c.com";

Windows Interact



Are all interactions good?

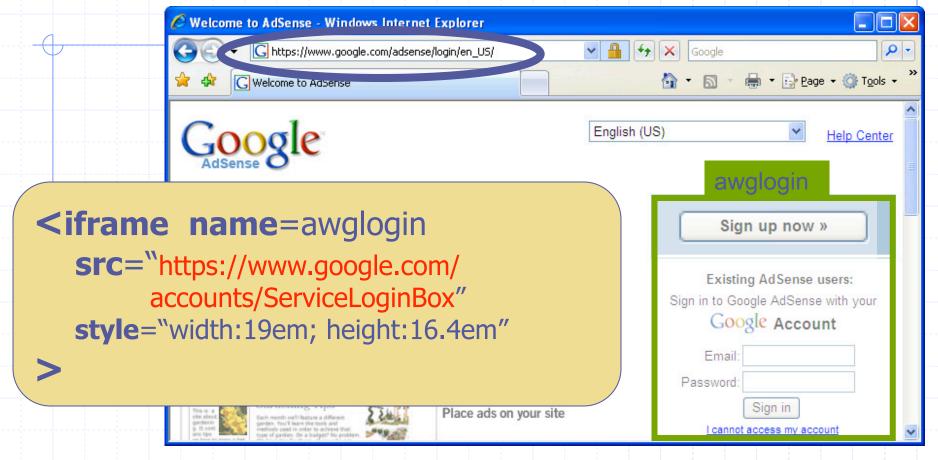


Frames

- Modularity
 - Brings together content from multiple sources
 - Client-side aggregation
- Delegation
 - Frame can draw only on its own rectangle

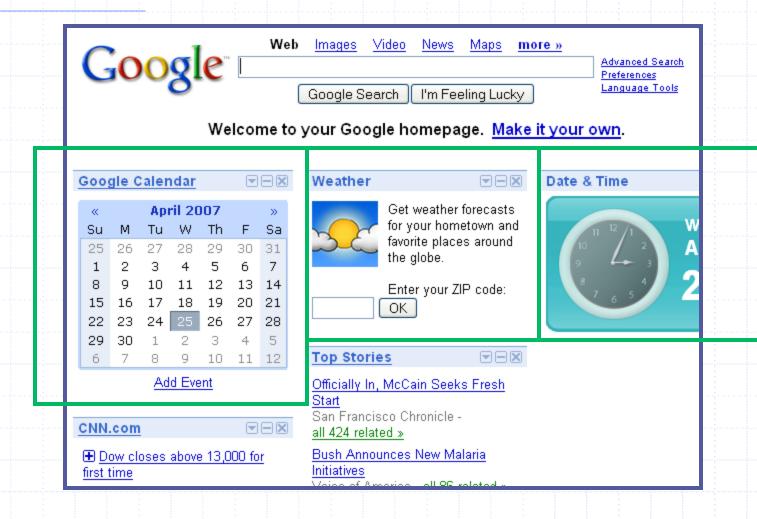


Frames and iFrames

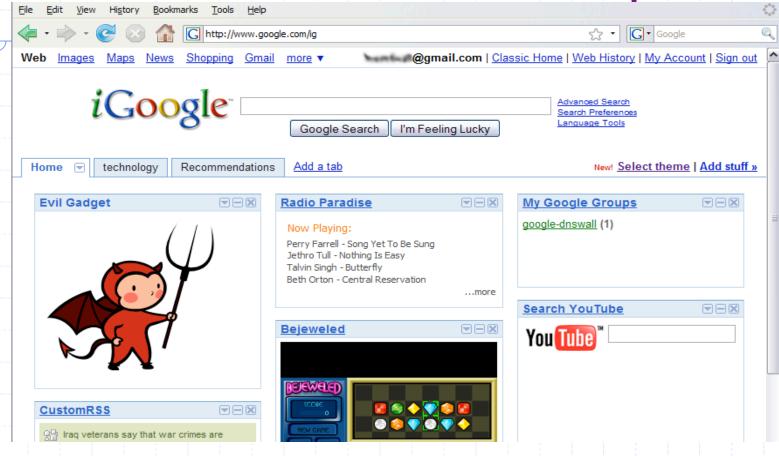


Address bar says nothing about origin of embedded content frames (ads), scripts, flash objects, CSS

Masups: lots of frames (gadgets)

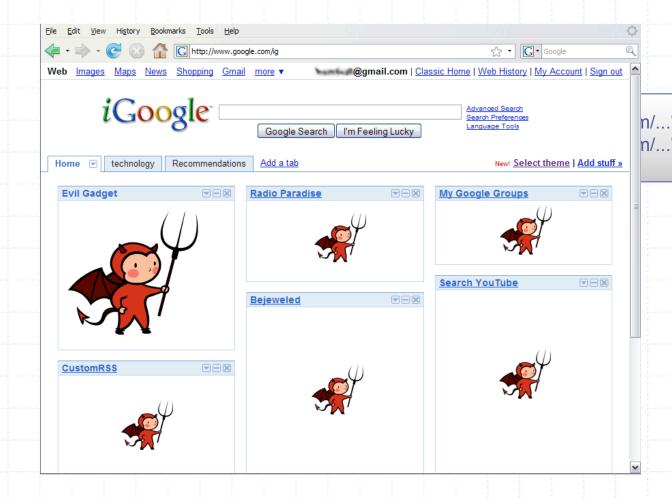


Need for isolation - mashups

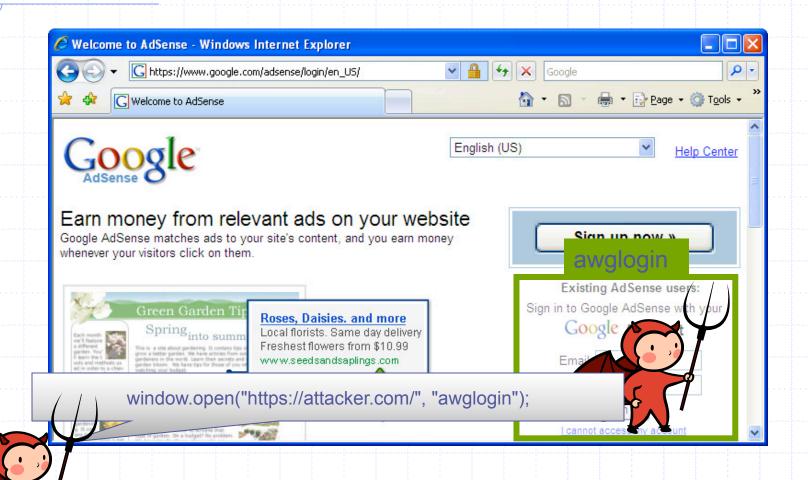


Malicious gadget should not affect other gadgets

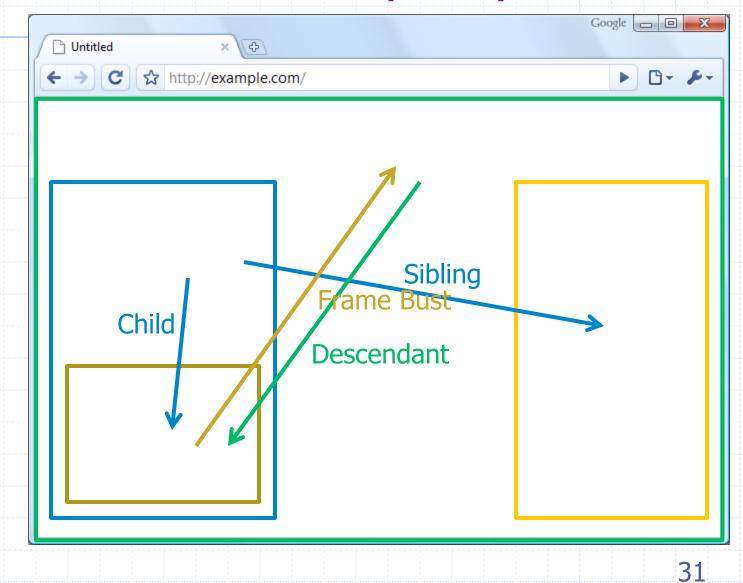
Window Policy Anomaly



A Guninski Attack



What should the policy be?



Legacy Browser Behavior

Browser	Policy
IE 6 (default)	Permissive
IE 6 (option)	Child
(no Flash)	Descendant
IE7 (with Flash)	Permissive
Firefox 2	Window
Safari 3	Permissive
Opera 9	Window
? HTML 5	Child

Adoption of Descendant Policy

	Browser	Policy
	IE7 (no Flash)	Descendant
	IE7 (with Flash)	Descendant
	Firefox 3	Descendant
	Safari 3	Descendant
	Opera 9	(many policies)
?	HTML 5	Descendant

Library import

<script src=https://seal.verisign.com/getseal?
host_name=a.com></script>







- Script has privileges of imported page, NOT source server.
- Can script other pages in this origin, load more scripts
- Other forms of importing









Pages can embed content from many sources (example)

```
Frames:
             <iframe src="//site.com/frame.html" > </iframe>
  Scripts:
                <script src="//site.com/script.js" > </script>
CSS:
k rel="stylesheet" type="text /css" href="//site/com/theme.css" />
Objects (flash): [using swfobject.js script]
   <script>
     var so = new SWFObject(\'//site.com/flash.swf', ...);
     so.addParam('allowscriptaccess', 'always');
     so.write('flashdiv');
   </script>
```

Cross-origin Interaction

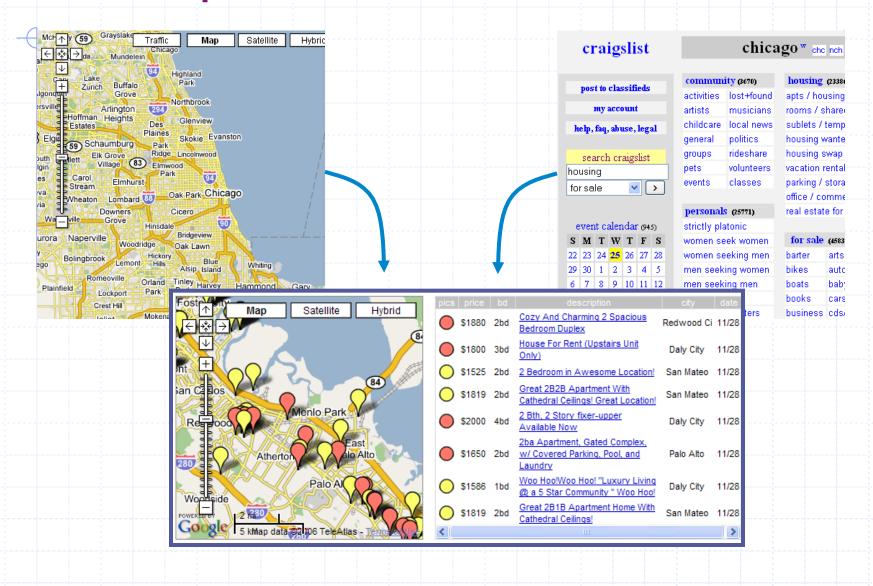
Sites often need to communicate:

Google AdSense:

<script src="http://googlesyndication.com/show_ads.js">

- Mashups
- Gadget aggregators (e.g. iGoogle or live.com)
- Primary method: script inclusion; site A does:
- <script src=//siteB.com/script.js>
- Script from B runs in A's origin: full control over A's DOM
- Note: to communicate with B, site A gives B full control !!

Mashups

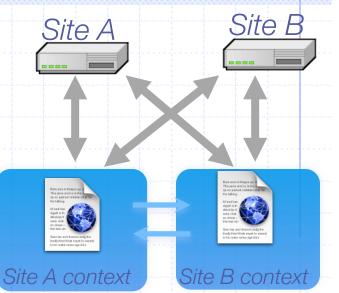


Need for isolation: embedded content



3rd party ad should not read/write enclosing DOM

Recent Developments



Cross-origin network requests

Access-Control-Allow-Origin: < list of domains>

Access-Control-Allow-Origin: *

Cross-origin client side communication

Client-side messaging via navigation (older browsers)

postMessage (newer browsers)

window.postMessage

- New API for inter-frame communication
 - Supported in latest betas of many browsers



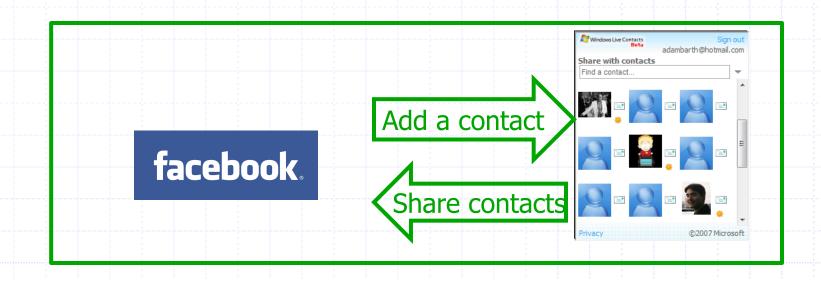








A network-like channel between frames



postMessage syntax

```
window.addEventListener("message", function (e) {
  if (e.origin == "http://a.com") {
    ... e.data ... }
}, false);
```

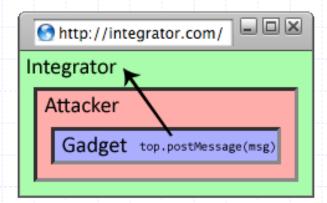


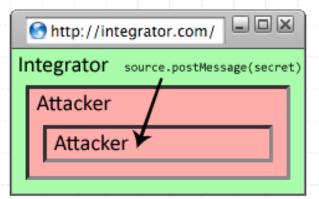
Attack at dawn!



Why include "targetOrigin"?

- What goes wrong?
 frames[0].postMessage("Attack at dawn!");
- Messages sent to frames, not principals
- When would this happen?





Data export

Many ways to send information to other origins

```
<form action="http://www.bank.com/">
  <input name="data" type="hidden" value="hello">
  </form>
```

-
- No user involvement required
- Cannot read back response
- Read response only from your origin
- Some port are restricted (SMTP)

Same Origin Requests with XMLHttpRequest

```
<script>
var xhr = new XMLHttpRequest();
                                             prepare request
xhr.open("POST", "http://www.example.com:81/foo/
  example.cgi", true); // asynchronous
xhr.send("Hello world!");
xhr.onload = function() {
 if (xhr.status == 200) {
  alert(xhr.responseText);
                                read response
</script>
```

Sending a Cross-Domain GET

- Data must be URL encoded
- Browser sends:

GET file.cgi?foo=1&bar=x%20y HTTP/1.1 Host: othersite.com

- ⇒ Any web page can send info to any site
- Denial of Service (DoS) using GET:
 - a popular site can DoS another site [Puppetnets '06]

Sending a Cross-Domain POST

```
<form method="POST" action="http://othersite.com/file.cgi" encoding="text/
plain">
```

<input type="hidden" name="Hello world" value="4">

</form>

<script>document.forms[0].submit()</script>

submit post

- Hidden iframe can do this in background
 - ⇒ user visits a malicious page, browser submits form on behalf of user
 - ⇒ e.g. page re-programs user's home router (XSRF)
- Can't send to some restricted ports, like 25 (SMTP)

Cookie Security How to make HTTP statefull securely?

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Same origin policy: "high level"

Review: Same Origin Policy (SOP) for DOM:

 Origin A can access origin B's DOM if match on (scheme, domain, port)

Today: Same Original Policy (SOP) for cookies:

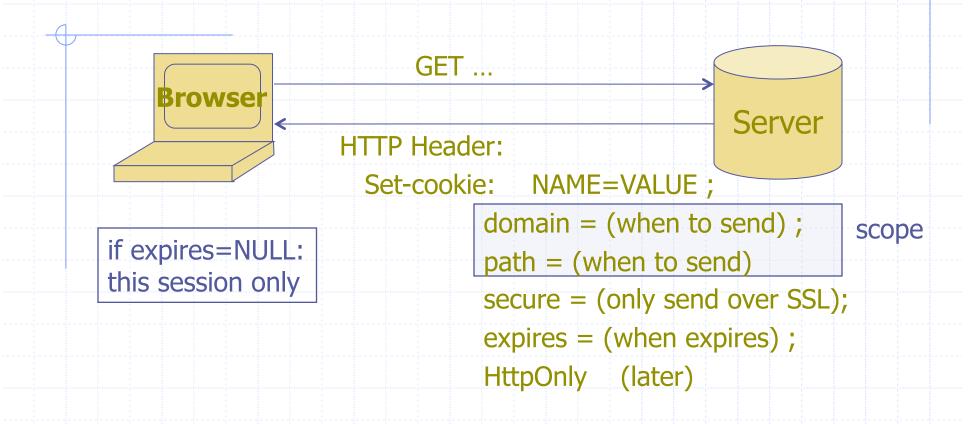
Generally speaking, based on:

([scheme], domain, path)

optional

scheme://domain:port/path?params

Setting/deleting cookies by server



- Delete cookie by setting "expires" to date in past
- Default scope is domain and path of setting URL

Scope setting rules (write SOP)

domain: any domain-suffix of URL-hostname, except TLD

example: host = "login.site.com"

allowed domains

login.site.com

disallowed domains

user.site.com othersite.com .com

 ⇒ login.site.com can set cookies for all of .site.com but not for another site or TLD
 Problematic for sites like .stanford.edu

path: can be set to anything

Cookies are identified by (name,domain,path)

```
cookie 1
name = userid
value = test
domain = login.site.com
path = /
secure
```

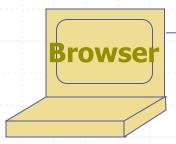
```
cookie 2
name = userid
value = test123
domain = .site.com
path = /
secure
```

distinct cookies

Both cookies stored in browser's cookie jar;
 both are in scope of login.site.com

Reading cookies on server

(read SOP)



GET //URL-domain/URL-path Cookie: NAME = VALUE

Server

Browser sends all cookies in URL scope:

- cookie-domain is domain-suffix of URL-domain, and
- cookie-path is prefix of URL-path, and
- [protocol=HTTPS if cookie is "secure"]

Goal: server only sees cookies in its scope

Examples

both set by login.site.com

```
cookie 1
name = userid
value = u1
domain = login.site.com
path = /
secure
```

```
cookie 2
name = userid
value = u2
domain = .site.com
path = /
non-secure
```

http://checkout.site.com/

http://login.site.com/

https://login.site.com/

cookie: userid=u2

cookie: userid=u2

cookie: userid=u1; userid=u2

(arbitrary order)

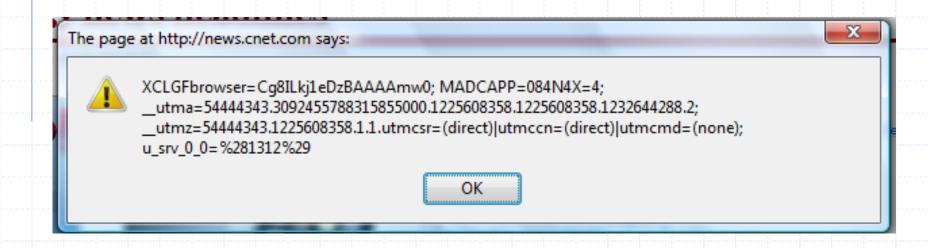
Client side read/write: document.cookie

- Setting a cookie in Javascript: document.cookie = "name=value; expires=...;"
- Reading a cookie: alert(document.cookie) prints string containing all cookies available for document (based on [protocol], domain, path)
- Deleting a cookie: document.cookie = "name=; expires= Thu, 01-Jan-70"

document.cookie often used to customize page in Javascript

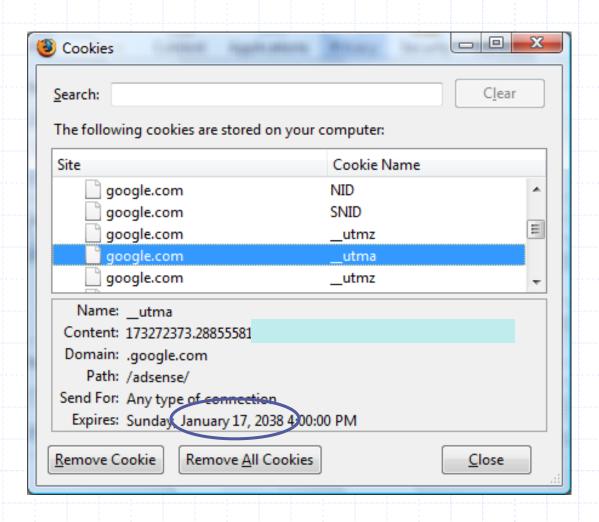
. Javascript URL

javascript: alert(document.cookie)



Displays all cookies for current document

Viewing/deleting cookies in Browser UI



Cookie protocol problems

Server is blind:

- Does not see cookie attributes (e.g. secure)
- Does not see which domain set the cookie

Server only sees: Cookie: NAME=VALUE

Interaction with the DOM SOP

Cookie SOP: path separation

x.com/A does not see cookies of x.com/B

Not a security measure:

DOM SOP: x.com/A has access to DOM of x.com/B

<iframe src="x.com/B"></iframe>
alert(frames[0].document.cookie);

Path separation is done for efficiency not security:

x.com/A is only sent the cookies it needs

HttpOnly Cookies

IE6 SP1, FF2.0.0.5

(not Safari)



- Cookie sent over HTTP(s), but not accessible to scripts
 - cannot be read via document.cookie
 - Also blocks access from XMLHttpRequest headers
 - Helps prevent cookie theft via XSS
 - ... but does not stop most other risks of XSS bugs.

Browser security design

How to build a secure browser?

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Approach

Fact: Browsers will always have bugs

Goal: Reduce the harm

Frequency of interactions with attacker



Percentage of time vulnerability is unpatched



Damage if attack works



Outline

Frequency of interactions with attacker



Percentage of time vulnerability is unpatched



Damage if attack works

1. Preventing the Introduction

2. Vulnerability Response

3. Failure Containment

Frequency of interactions with attacker



Percentage of time vulnerability is unpatched

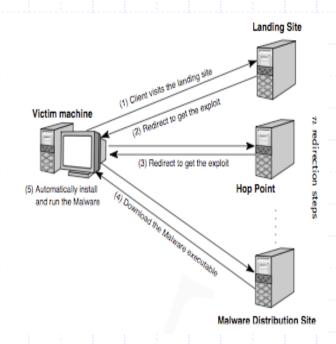


Damage if attack works

PREVENTING THE INTRODUCTION

Drive-by downloads

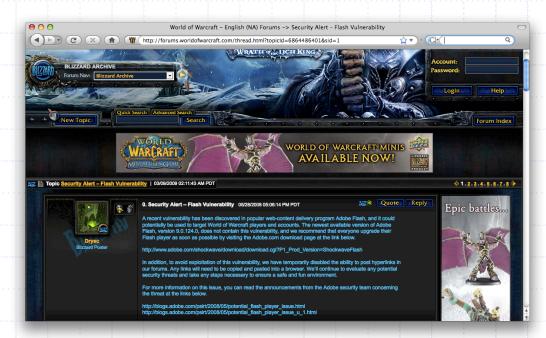
- Silently installs software when web page is loaded
- Increase exposure by compromising other sites and insert code into them
- Sites owners unaware they are participating in an attack



Provos et al. "All your iFRAMES Point to Us"

World of Warcraft keylogger

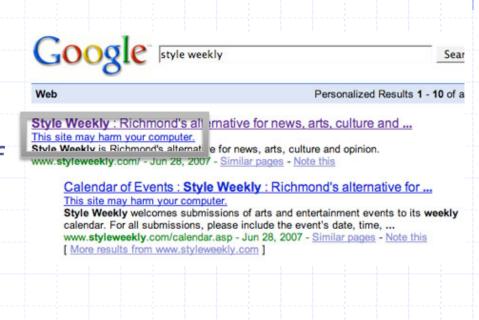
- Flash Player exploit used to install keylogger
- Links to malicious SWF posted on forums



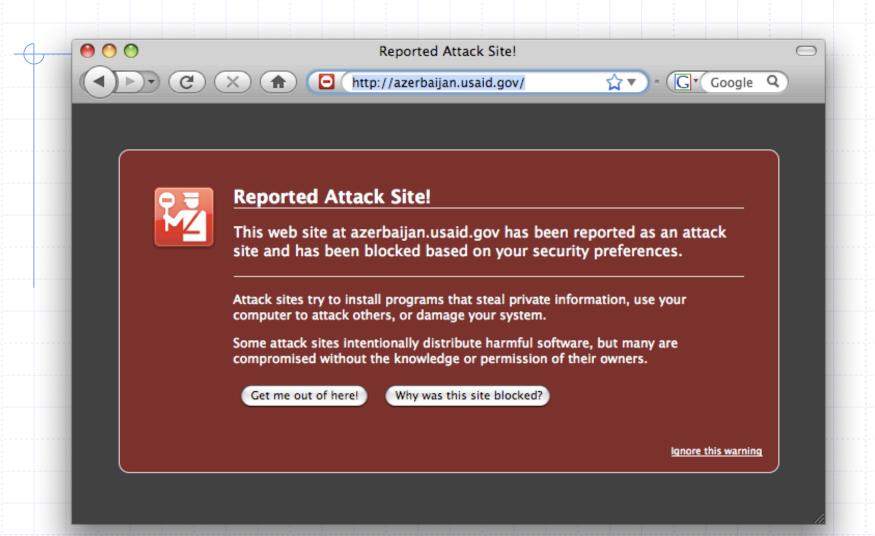
"Solution": Disable hyperlinks on forum

Scaling it up to the entire web

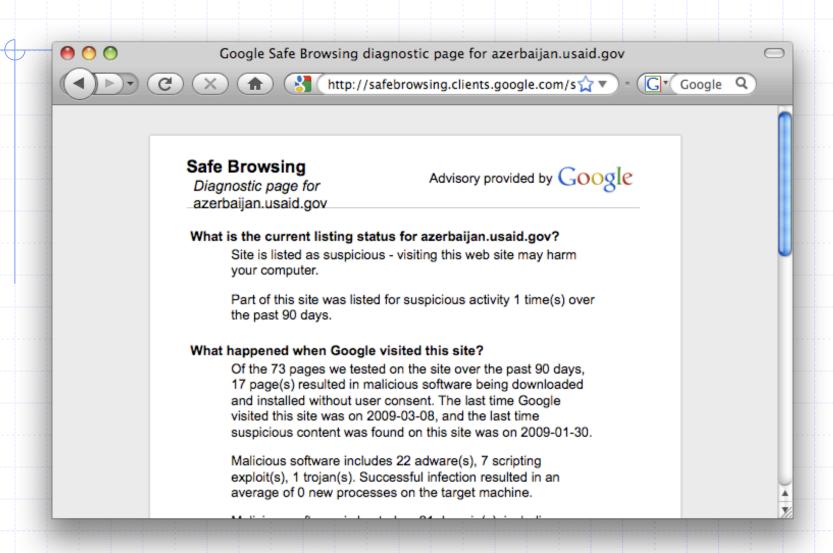
- 1.3% of the incoming search queries to Google returned at a least one malware site
- Visit sites with an army of browsers in VMs, check for changes to local system
- Indicate potentially harmful sites in search results



Now do it in the browser

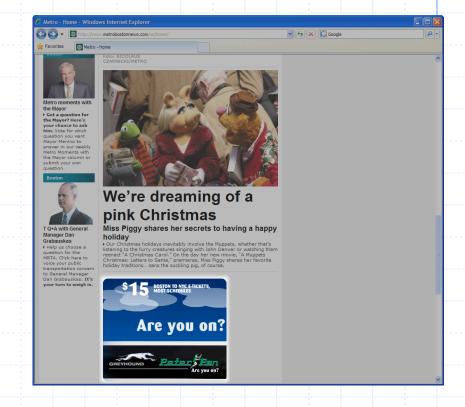


Helping the webmaster out



Introductions are easy

- Impressions are cheap (\$1 = 2000)
- Ad that is harmless today may be malicious tomorrow
- Possible mitigations:
 <iframe</p>
 security=restricted>
 <iframe sandbox>



Frequency of interactions with attacker



Percentage of time vulnerability is unpatched



Damage if attack works

VULNERABILITY RESPONSE

Closing the vulnerability window

Discovery

Publication

Patch available

Patch deployed

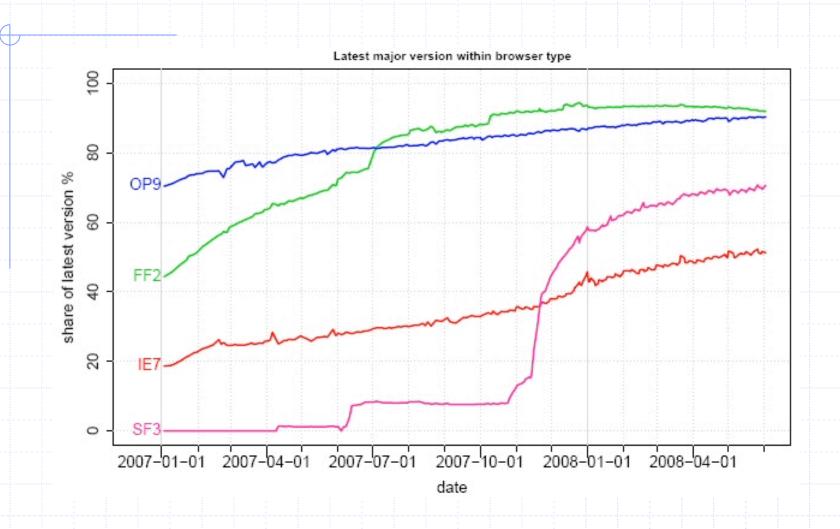
- Delay publication
 - Coordinate with security researchers
 - Offer prizes for responsibly disclosed security bugs
- Make patch available faster
- Deploy patch faster

Obstacles to patch deployment

- Interrupts work flow
- Requires adminstrator privileges
- Risk of breaking things
- Separate update mechanisms
- Silent approach:GoogleUpdate.exe



Getting better, but not fast enough



Frei et al. Examination of vulnerable online Web browser populations and the "insecurity iceberg"

Frequency of interactions with attacker

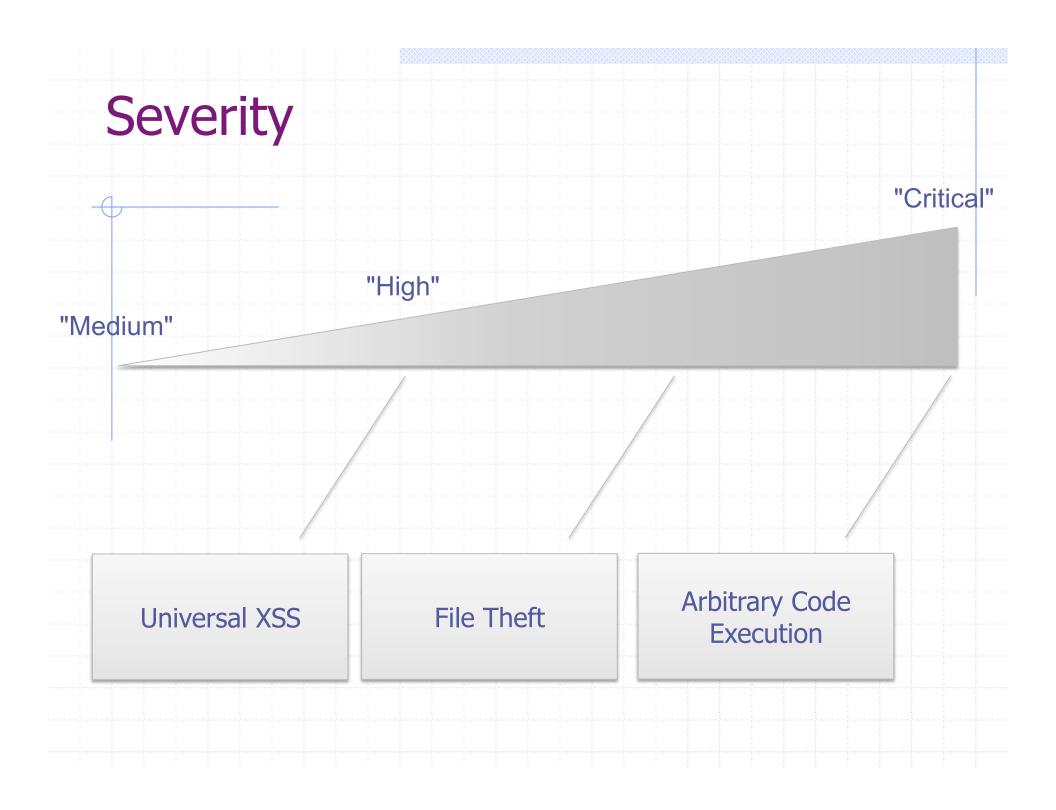


Percentage of time vulnerability is unpatched



Damage if attack works

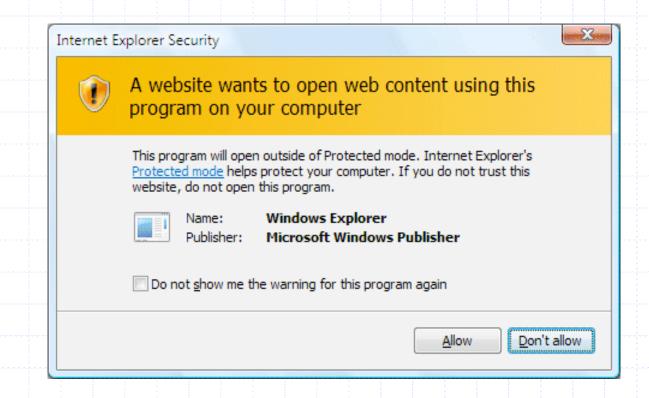
FAILURE CONTAINMENT



Protected Mode IE

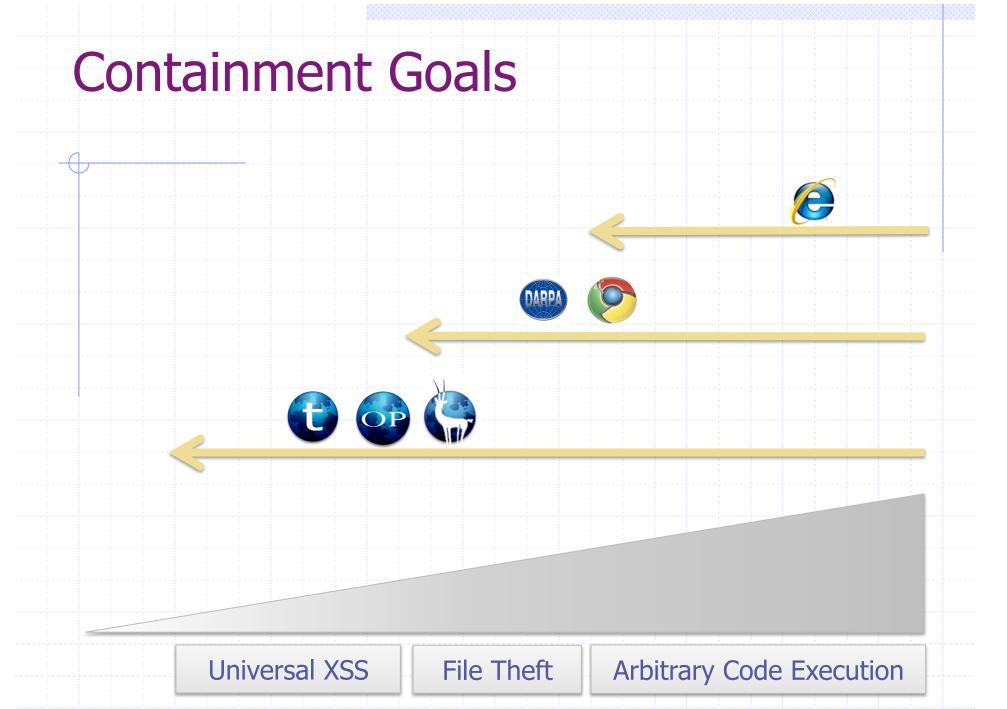


- ◆ IE7 in Vista is a "low rights" process
- Can prompt user to get more privileges



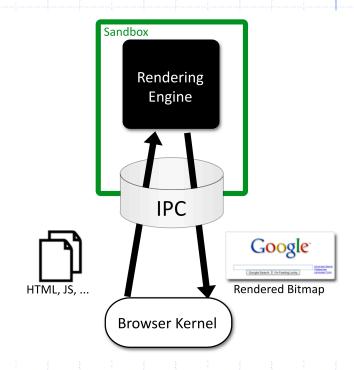
IE7 Containment Goals

- Arbitrary code execution won't let attacker:
 - Install software
 - Copy files to startup folder
 - Change homepage or search provider setting
- Can we do more?



Chromium Security Architecture

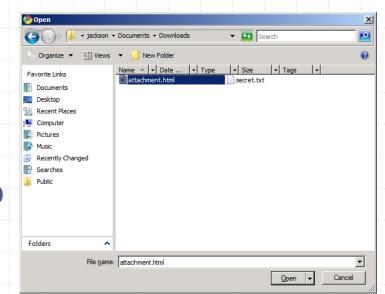
- Browser ("kernel")
 - Full privileges (file system, networking)
 - Coarse-grained security policies protect local system
- Rendering engine
 - Sandboxed
 - Fine-grained same origin policy enforcement
- One process per plugin
 - Sandboxing optional



Barth et al. "The Security Architecture of the Chromium Browser"

Preventing File Theft

- File Downloads.
 - Renderer can only write files to My Documents\Downloads
- File Uploads.
 - Renderer is granted ability to upload file using browser kernel's file picker.
- Network Requests.
 - Can only request web-safe schemes (http, https, ftp)
 - Dedicated renderers for file://



Task Allocation



Rendering Engine

HTML parsing
CSS parsing
Image decoding
JavaScript interpreter
Regular expressions
Layout
Document Object Model
Rendering
SVG
XML parsing
XSLT

Browser Kernel

Cookie database
History database
Password database
Window management
Location bar
Safe Browsing blacklist
Network stack
SSL/TLS
Disk cache
Download manager
Clipboard



Both

URL parsing Unicode parsing

Is the "kernel" too complex?

◆ Total CVEs:

	Browser	Renderer	Unclassified
Internet Explorer	4	10	5
Firefox	17	40	3
Safari	. 12	37	1

Arbitrary code execution vulnerabilities

		Browser	Renderer	Unclassified
[Internet Explorer	1	9	5
-	Firefox	5	19	0
~-	\mathbf{Safari}	5	10	0

Another approach: Cookie Blocking

- Block the "Cookie" header for cross-domain resource loads
- Third-party cookie blocking already does this for privacy
- Third-party frames are ok
- Cross-subdomain might be ok



Open question: How many sites does this break compared to content type filtering?

Conclusion

Frequency of interactions with attacker



Percentage of time vulnerability is unpatched



Damage if attack works

1. Preventing the Introduction

2. Vulnerability Response

3. Failure Containment