Program Analysis for Security

John Mitchell

Software bugs are serious problems

07-31-2010, 12:57 PM

911crashes
Junior Member

Join Date: Jul 2010
Posts: 2

[OFFLINE]

calling 911 crashes my HTC evo 4G, every time

I happen to need to call 911 one night, found that my phone crashes every time I dial 911.

my wife's phone does not do that, any thought?

by the way, it is hard to test this problem due to the sensitivity of calling 911 repeatedly.
thanks,

heartboken

Thanks: Isil and Thomas Dillig

Facebook missed a single security check…

Man Finds Easy Hack to Delete Any Facebook Photo Album

Facebook awards him a $12,500 "bug bounty" for his discovery

[PopPhoto.com Feb 10]

App stores

Apps for whatever you're up for.

Stay on top of the news. Stay on top of your finances. Or plan your dream vacation. No matter what you want to do with your iPhone, there's probably an app to help you do it.

Business

iPhone is ready for work. Manage projects, track stocks, monitor finances, and more with these 9-to-5 apps.

View business apps in the App Store ›

Education

Keep up with your studies using intelligent education apps like King of Math and NatureTap.

View education apps in the App Store ›

Entertainment

Kick back and enjoy the show. Or find countless other ways to entertain yourself. These apps offer hours of viewing pleasure.

View entertainment apps in the App Store ›

Family & Kids

Turn every night into family night with interactive apps that are fun for the whole house.

View family and kids apps in the App Store ›

Finance

Create budgets, pay bills, and more with financial apps that take everything into account.

View finance apps in the App Store ›

Food & Drink

Hungry? Thirsty? A little of both? Learn new recipes, drinks, and the secrets behind what makes a great meal.

View food and drink apps in the App Store ›

How can you tell whether software you

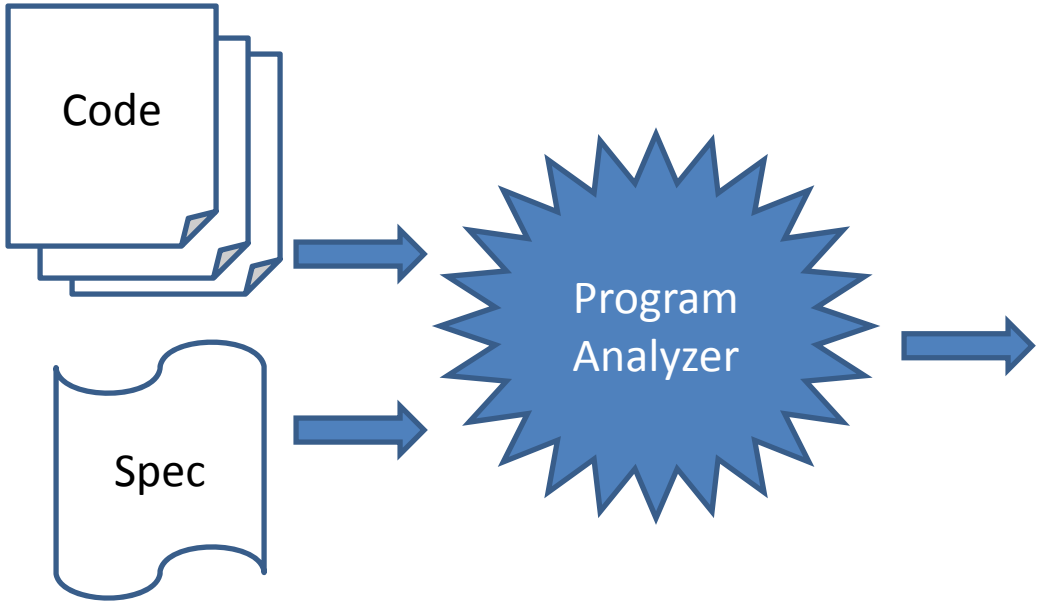
- Develop
- Buy

is safe to install and run?

Two options

- Static analysis
  - Inspect code or run automated method to find errors or gain confidence about their absence
- Dynamic analysis
  - Run code, possibly under instrumented conditions, to see if there are likely problems

Program Analyzers



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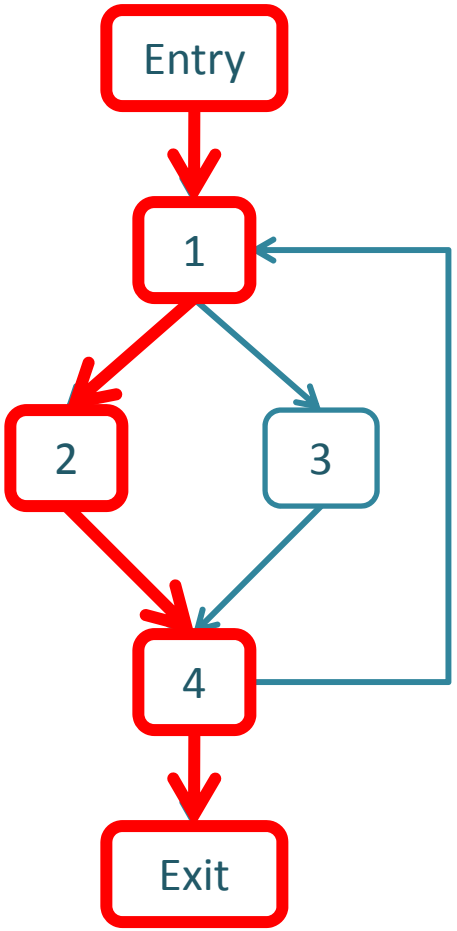
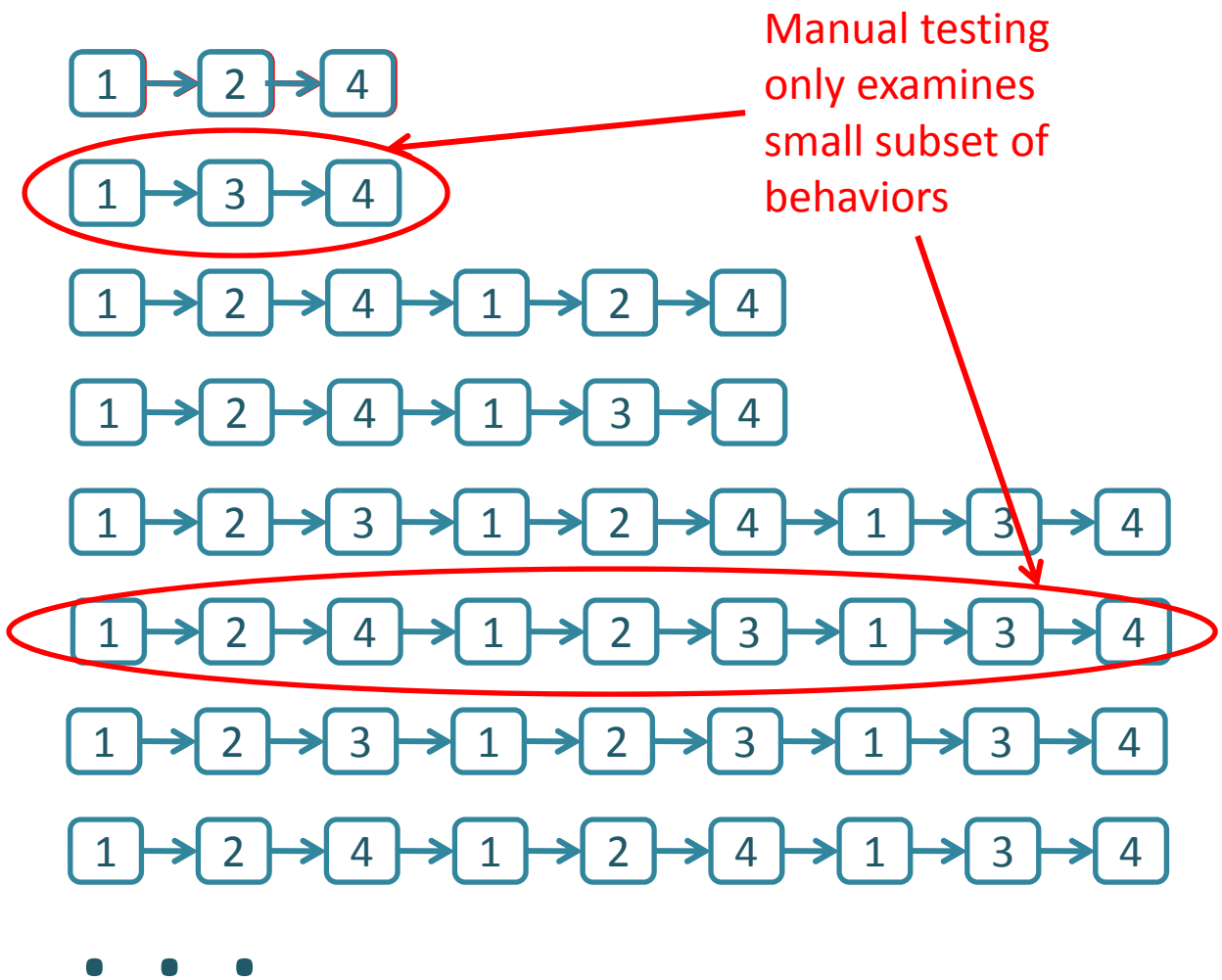
5	dang ptr	8,491
…	…	…
10,502	info leak	10,921



Software

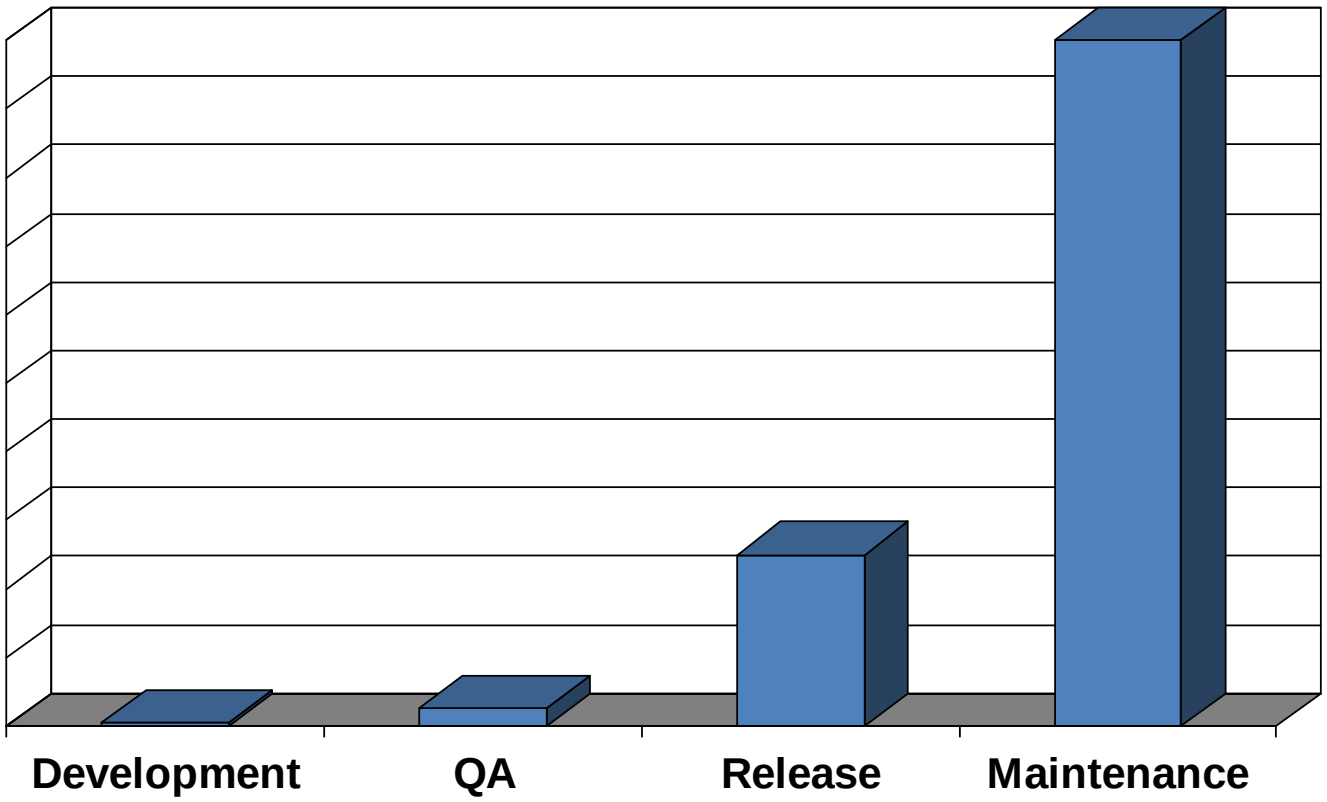


Behaviors

Static vs Dynamic Analysis

- Static
  - Can consider all possible inputs
  - Find bugs and vulnerabilities
  - Can prove absence of bugs, in some cases
- Dynamic
  - Need to choose sample test input
  - Can find bugs vulnerabilities
  - Cannot prove their absence

Cost of Fixing a Defect



Credit: Andy Chou, Coverity

Cost of security or data privacy vulnerability?

Dynamic analysis

- Instrument code for testing
  - Heap memory: Purify
  - Perl tainting (information flow)
  - Java race condition checking
- Black-box testing
  - Fuzzing and penetration testing
  - Black-box web application security analysis

Static Analysis

- Long research history
- Decade of commercial products
  - FindBugs, Fortify, Coverity, MS tools, …

Static Analysis: Outline

- General discussion of static analysis tools
  - Goals and limitations
  - Approach based on abstract states
- More about one specific approach
  - Property checkers from Engler et al., Coverity
  - Sample security checkers results
- Static analysis for of Android apps

Slides from: S. Bugrahe, A. Chou, I&T Dillig, D. Engler, J. Franklin, A. Aiken, …

Static analysis goals

- Bug finding
  - Identify code that the programmer wishes to modify or improve
- Correctness
  - Verify the absence of certain classes of errors

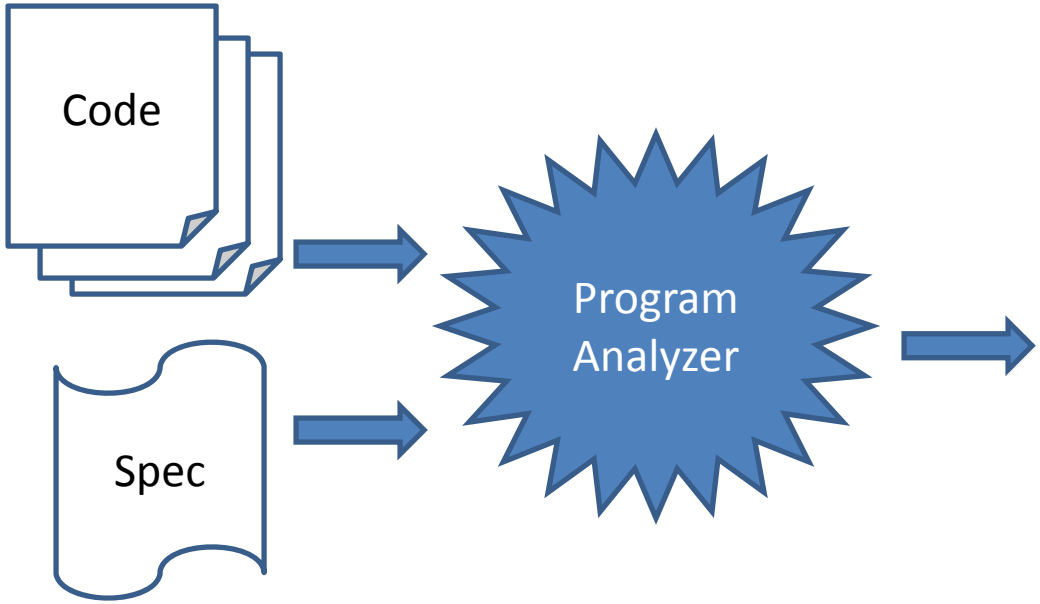
Soundness, Completeness

Property	Definition
Soundness	"Sound for reporting correctness" Analysis says no bugs $\rightarrow$ No bugs or equivalently There is a bug $\rightarrow$ Analysis finds a bug
Completeness	"Complete for reporting correctness" No bugs $\rightarrow$ Analysis says no bugs

Recall: $A \rightarrow B$ is equivalent to $(\neg B) \rightarrow (\neg A)$

	Complete	Incomplete
Sound	Reports all errors Reports no false alarms **Undecidable**	Reports all errors May report false alarms **Decidable**
Unsound	May not report all errors Reports no false alarms **Decidable**	May not report all errors May report false alarms **Decidable**

Sound Program Analyzer



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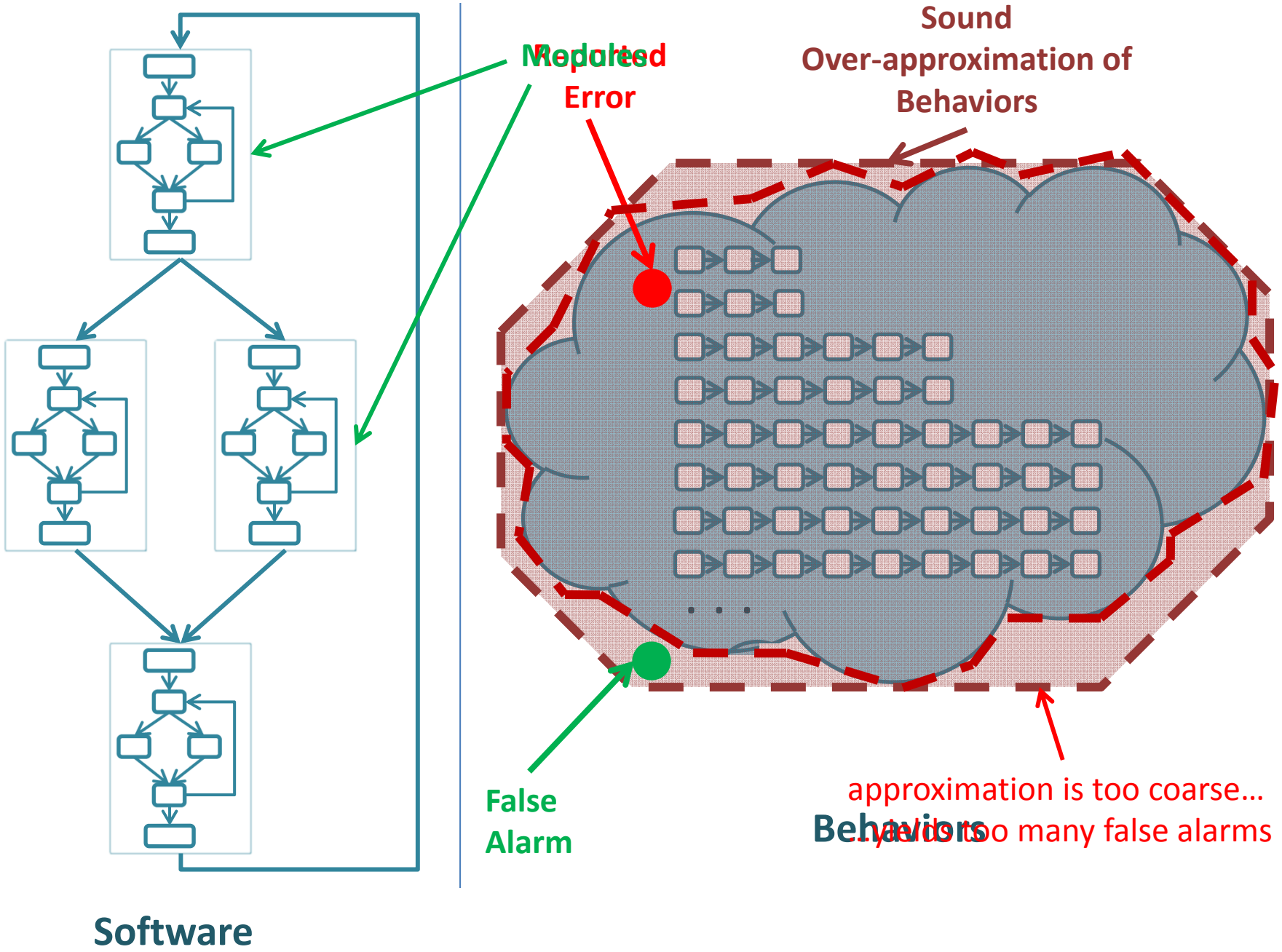
false alarm

false alarm

Analyze large code bases

Sound: may report many warnings

May emit false alarms

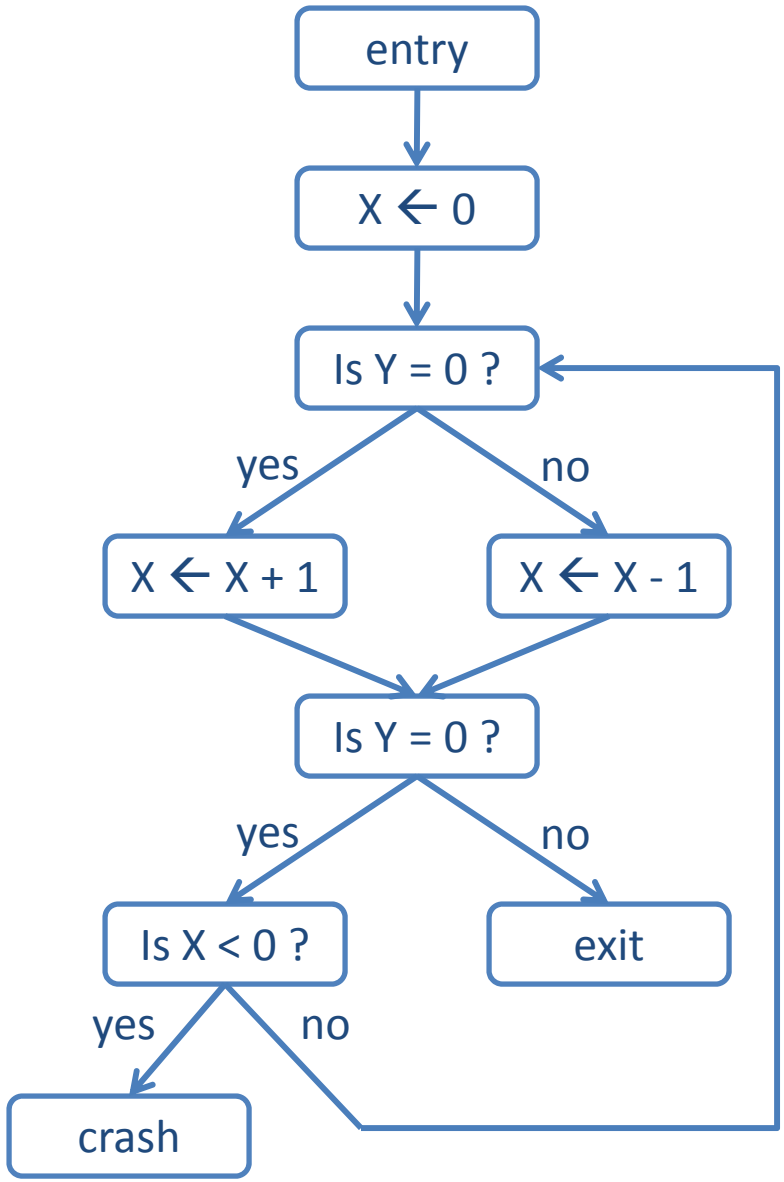

Sound
Over-approximation of
Behaviors
Reported
Error
False
Alarm
approximation is too coarse…
…yields too many false alarms
Behaviors
Software

Outline

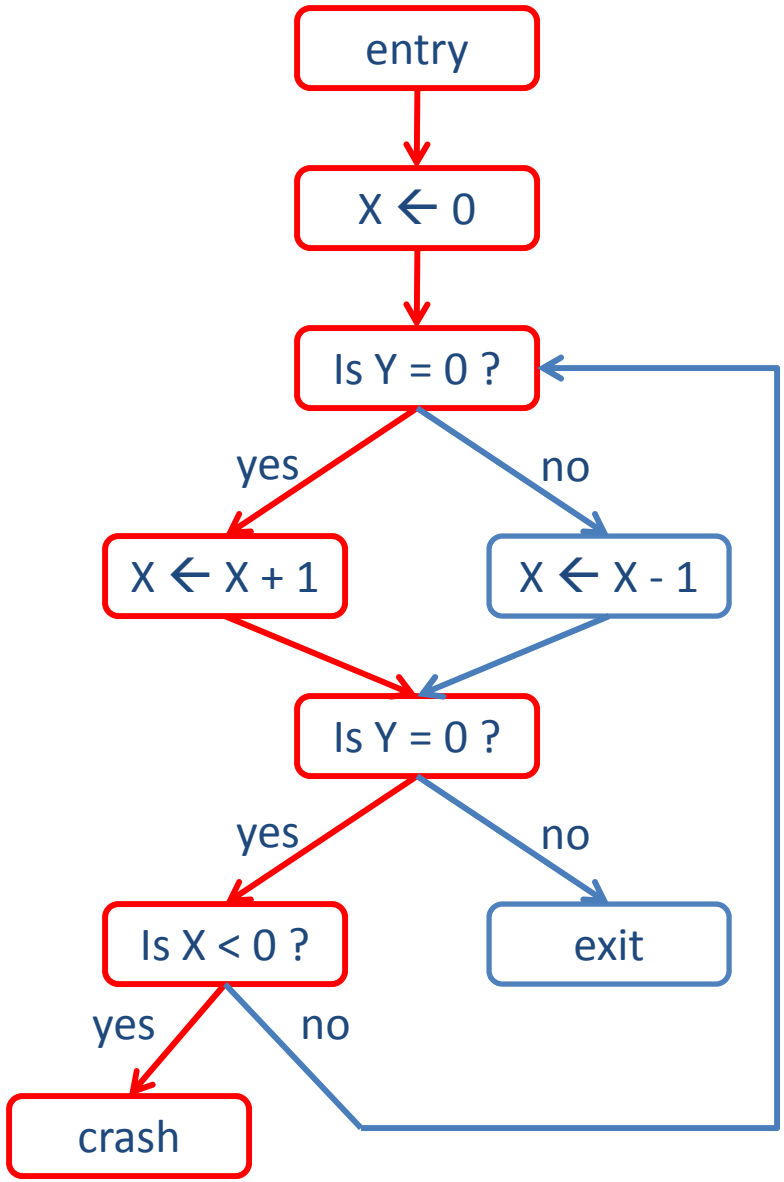
- General discussion of tools
  - Goals and limitations
  - ➡ Approach based on abstract states
- More about one specific approach
  - Property checkers from Engler et al., Coverity
  - Sample security-related results
- Static analysis for Android malware
  - …

Slides from: S. Bugrahe, A. Chou, I&T Dillig, D. Engler, J. Franklin, A. Aiken, …

Does this program ever crash?



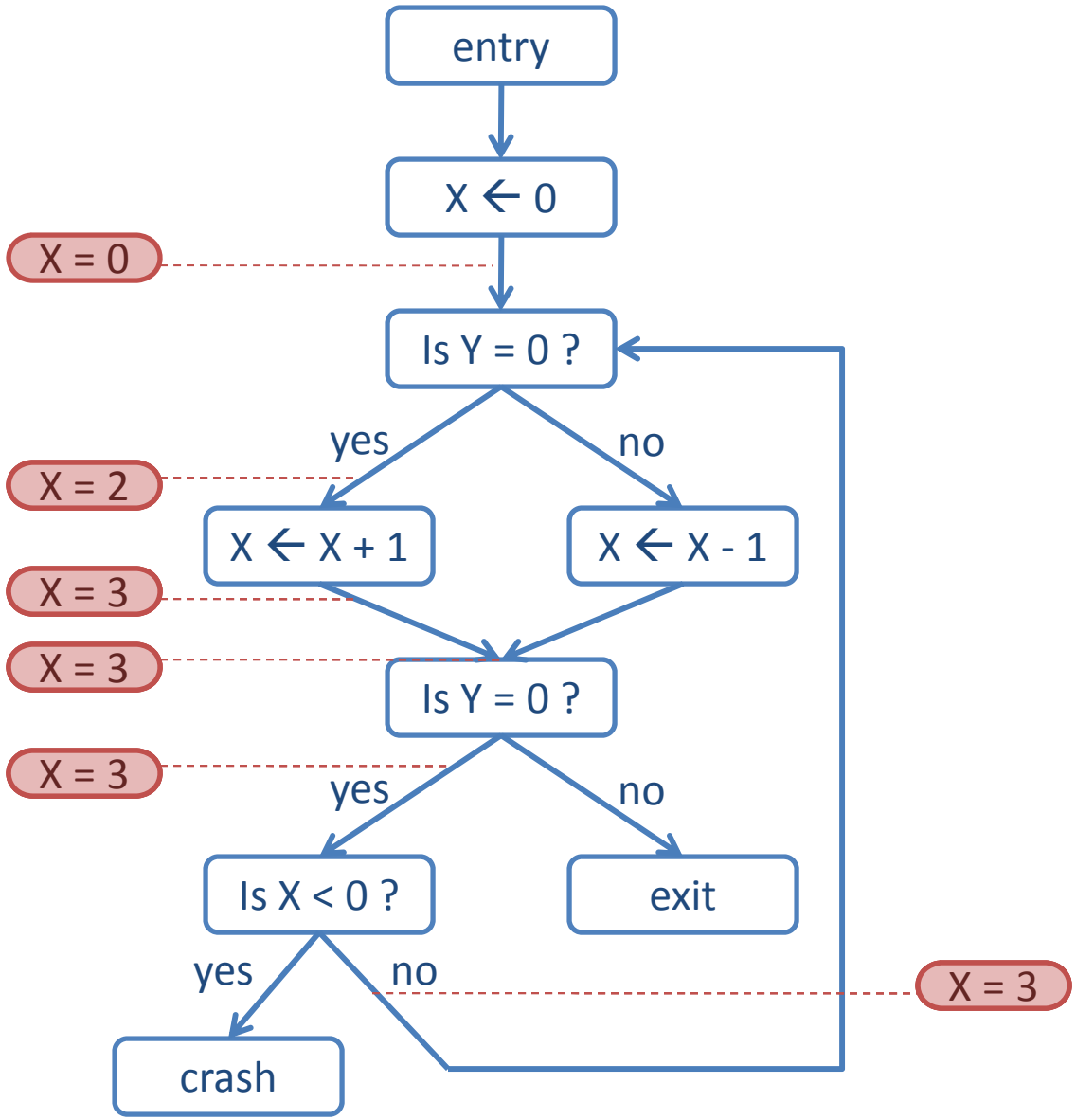
Does this program ever crash?



infeasible path!

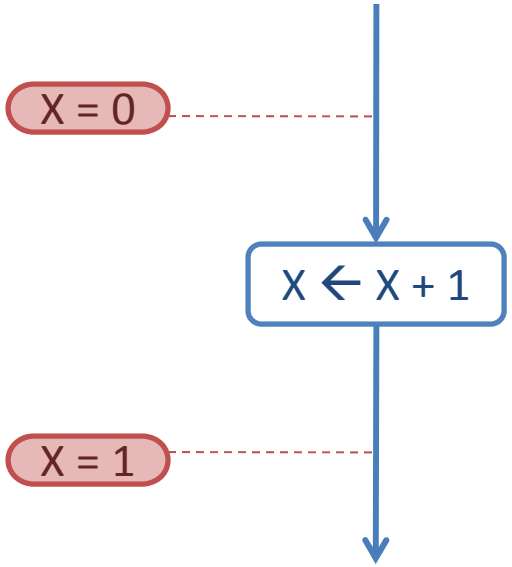
… program will never crash

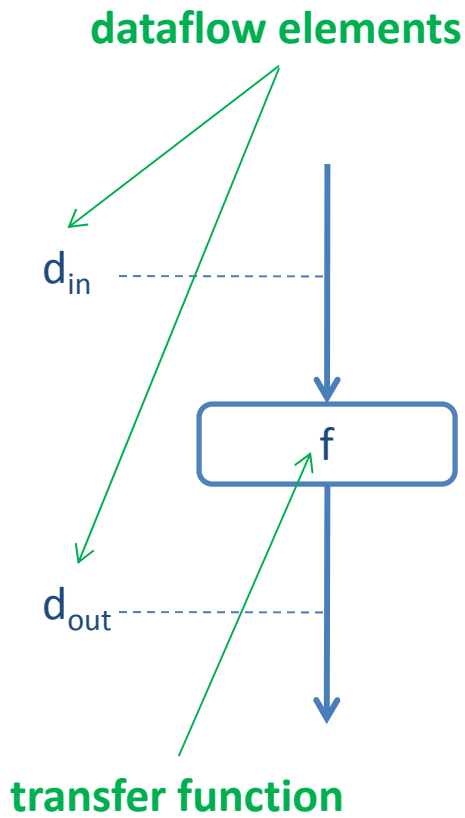
Try analyzing without approximating…



non-termination!

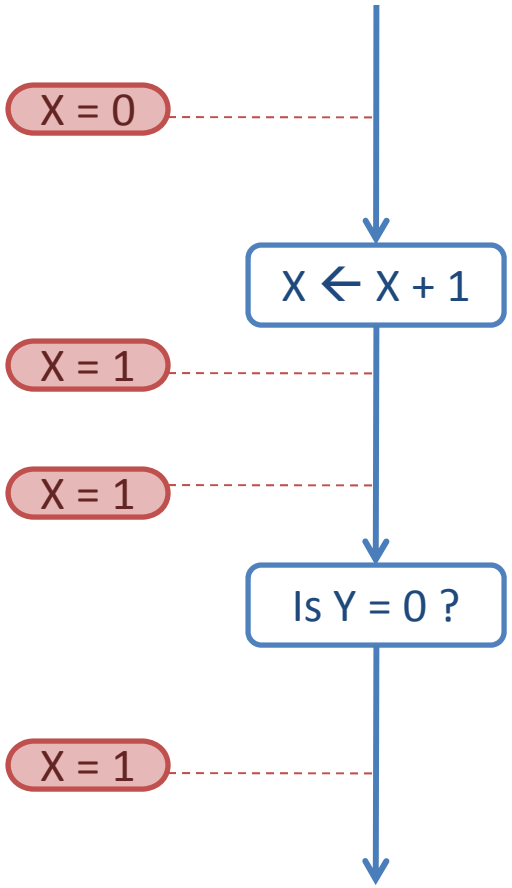
… therefore, need to approximate

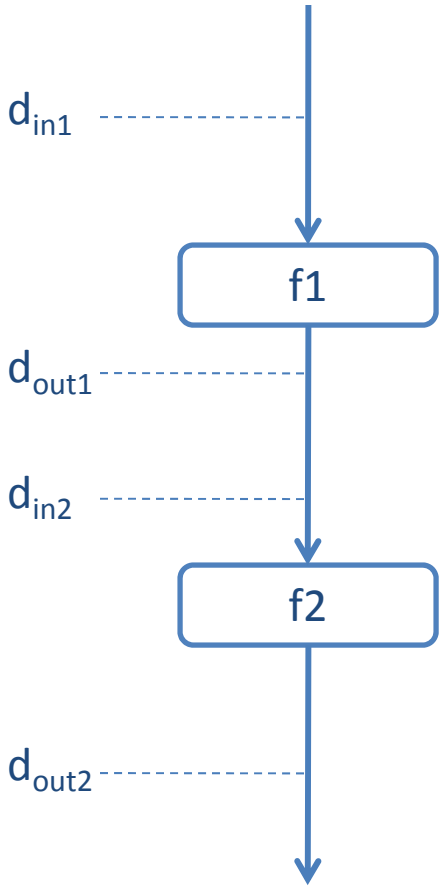




$$d_{out} = f(d_{in})$$

dataflow equation

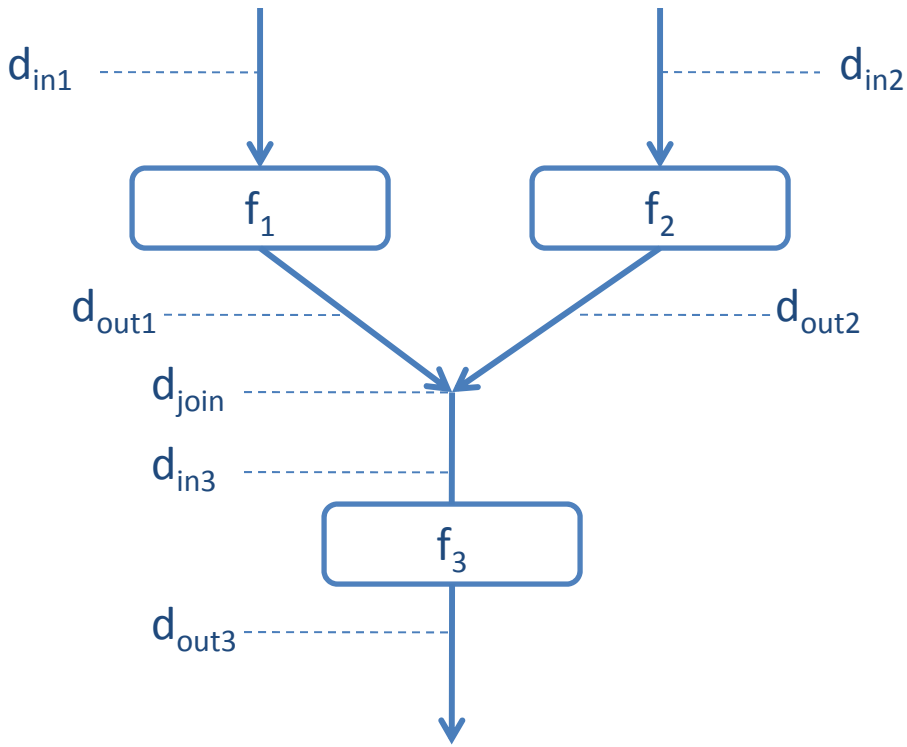




$$d_{out1} = f_1(d_{in1})$$

$$d_{out1} = d_{in2}$$

$$d_{out2} = f_2(d_{in2})$$



$$d_{out1} = f_1(d_{in1})$$

$$d_{out2} = f_2(d_{in2})$$

$$d_{join} = d_{out1} \sqcup d_{out2}$$

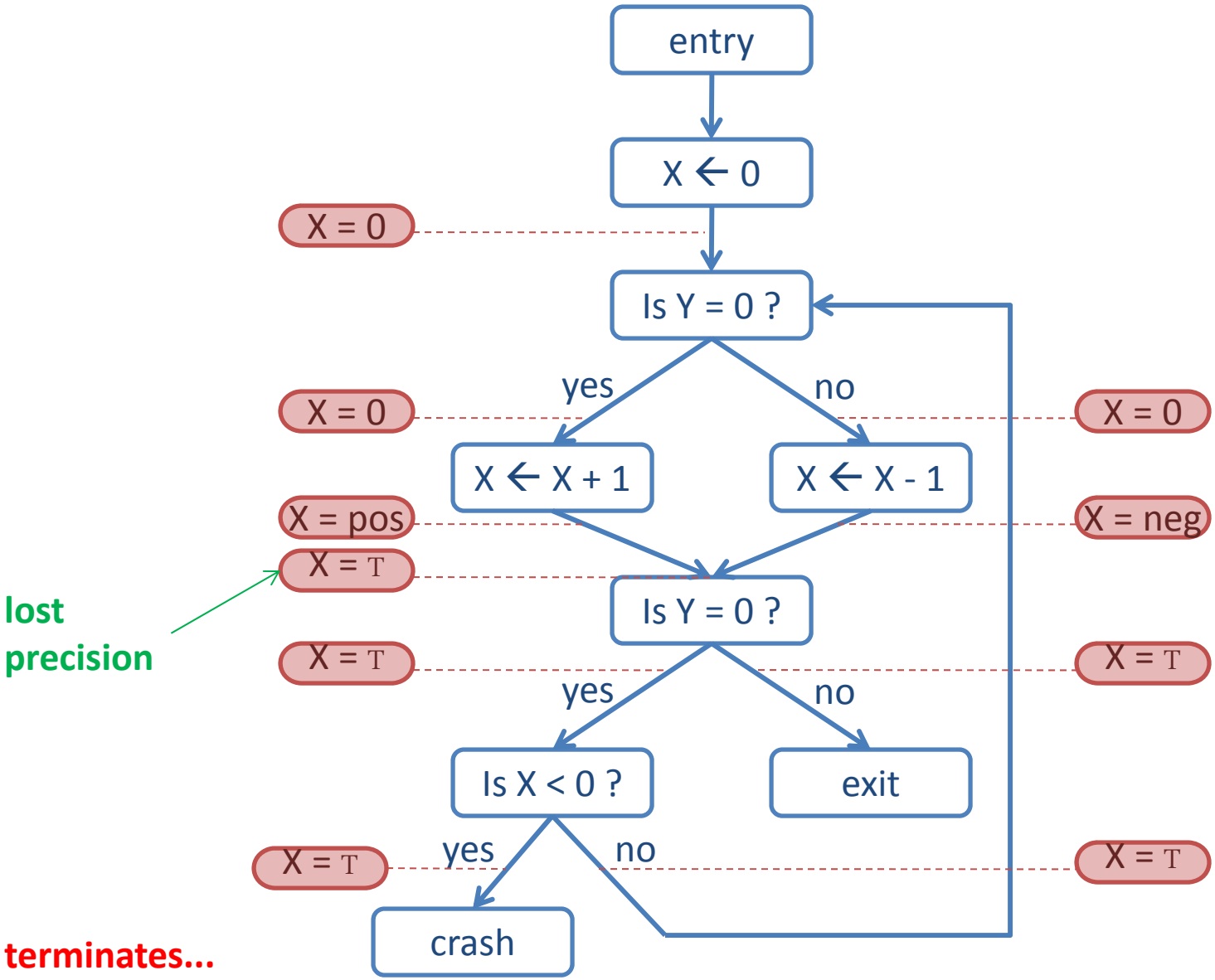
$$d_{join} = d_{in3}$$

$$d_{out3} = f_3(d_{in3})$$

least upper bound operator
Example: union of possible values

What is the space of dataflow elements, Δ?
What is the least upper bound operator, $\sqcup$?

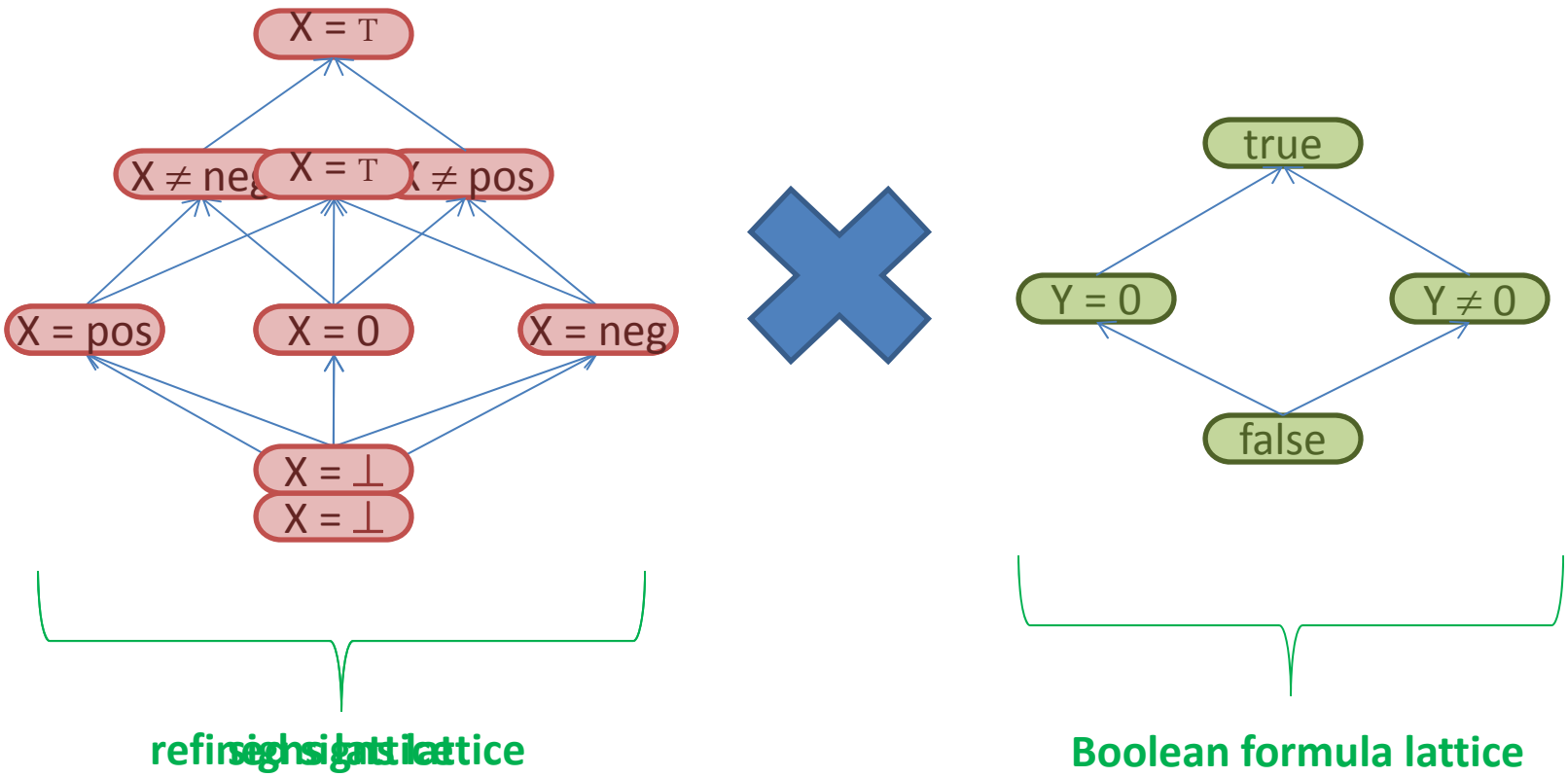
Try analyzing with “signs” approximation…



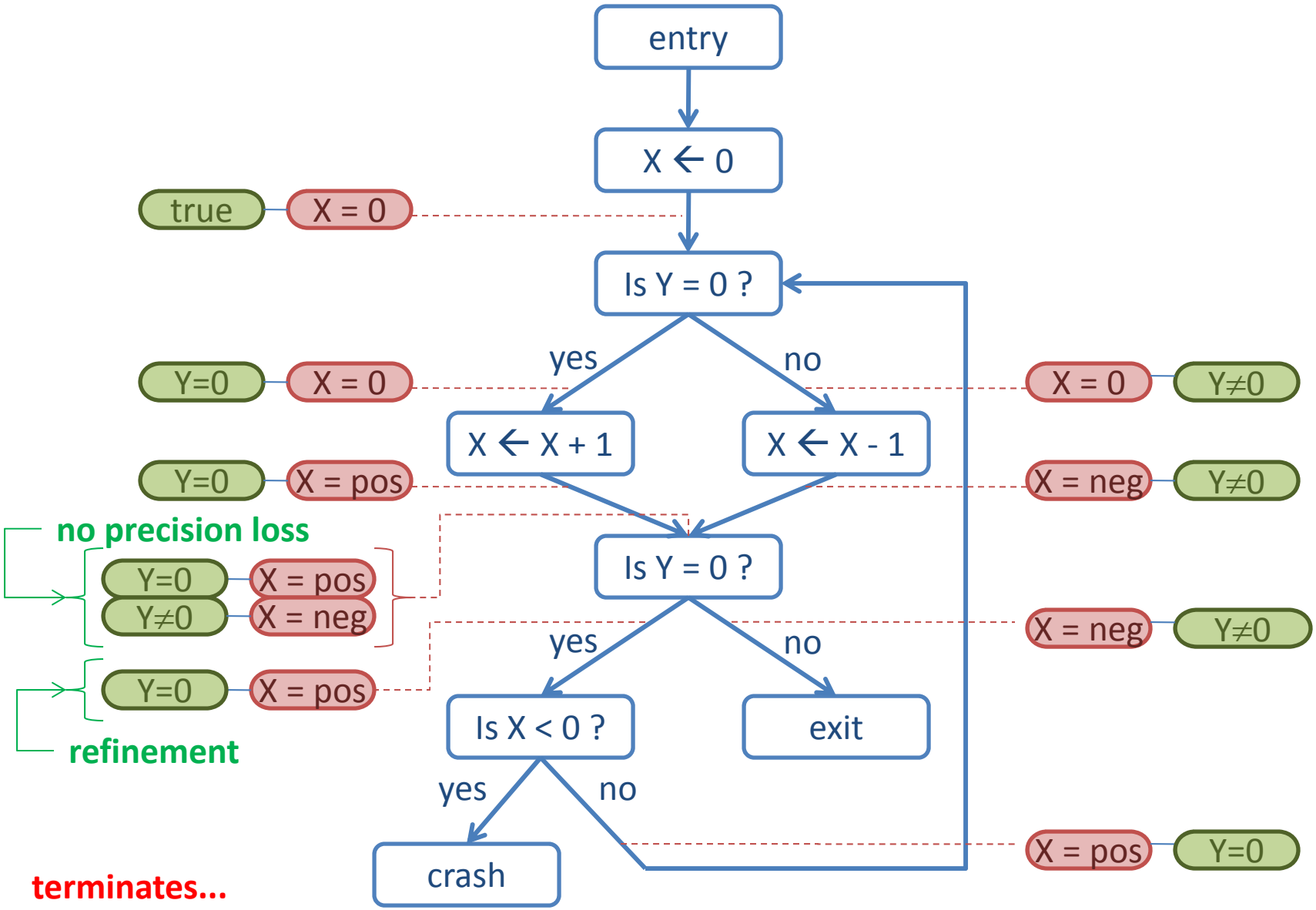
terminates...

… but reports false alarm

… therefore, need more precision


X = T
X ≠ neg
X = T
X ≠ pos
X = pos
X = 0
X = neg
X = ⊥
X = ⊥
true
Y = 0
Y ≠ 0
false
sign lattice
Boolean formula lattice

Try analyzing with “path-sensitive signs” approximation…



terminates...

… no false alarm

… soundly proved never crashes

Outline

- General discussion of tools
  - Goals and limitations
  - Approach based on abstract states
- ★ More about one specific approach
  - Property checkers from Engler et al., Coverity
  - Sample security-related results
- Static analysis for Android malware
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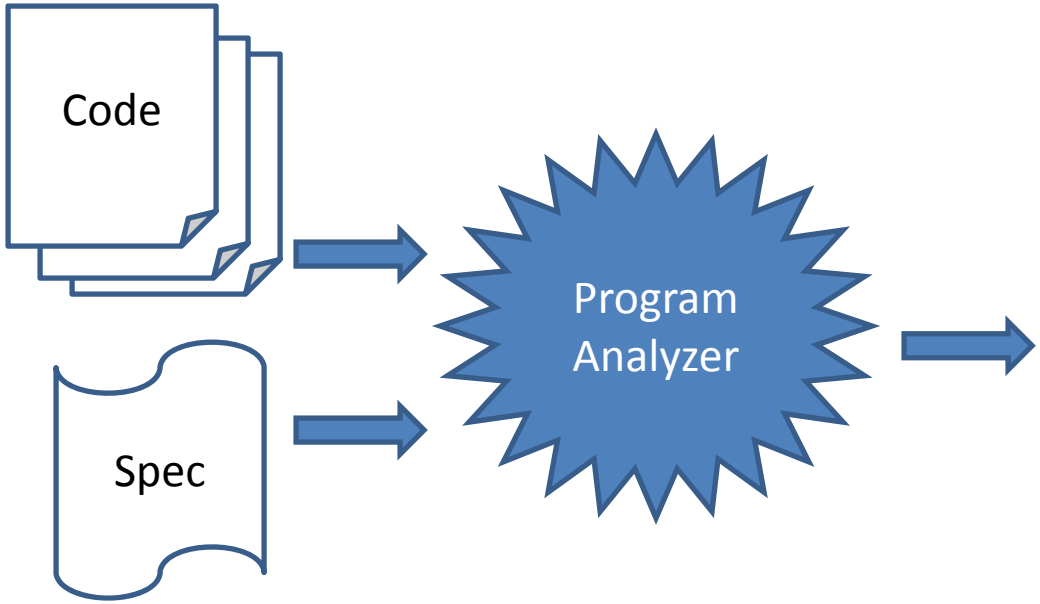
Unsound Program Analyzer



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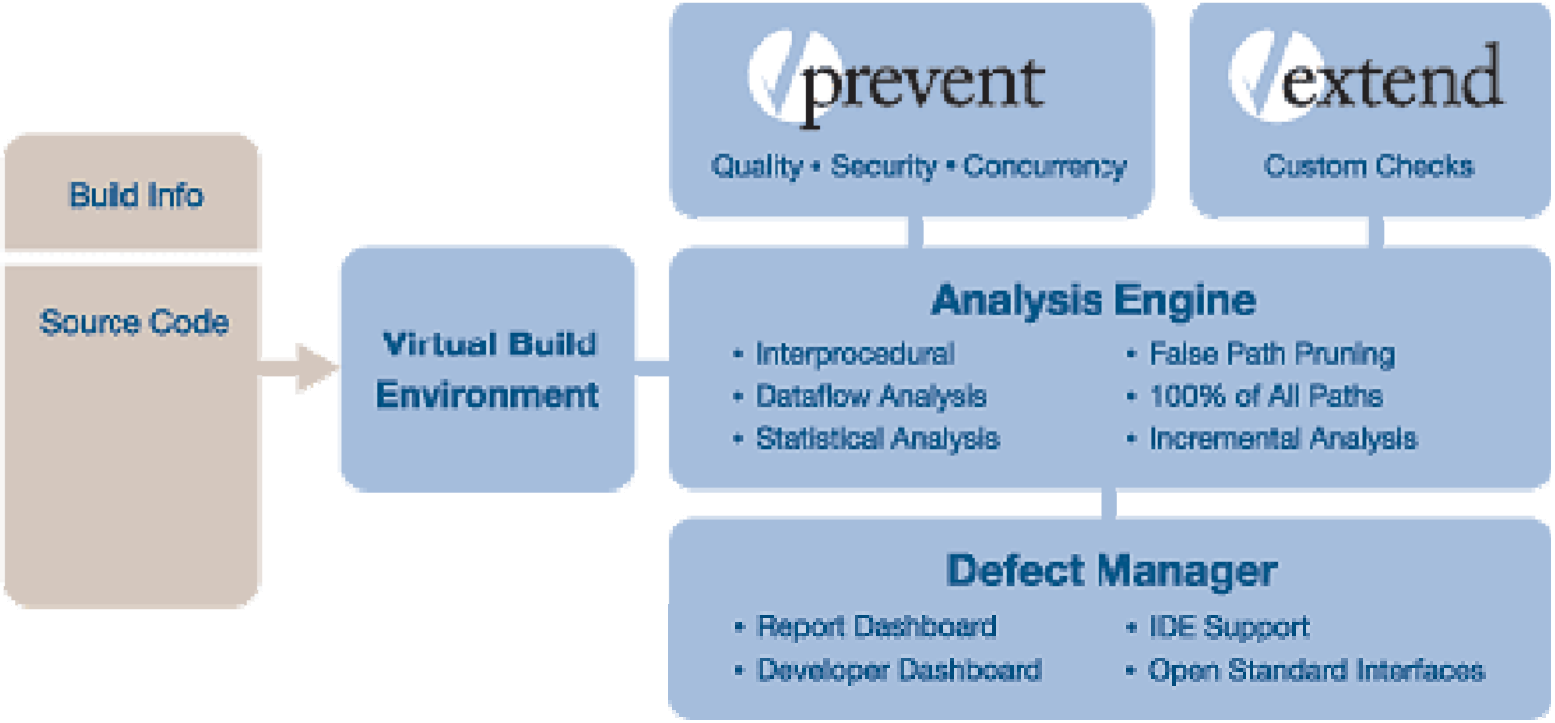
false alarm

false alarm

analyze large code bases

Not sound: may miss some bugs

may emit false alarms


Build Info
Source Code
Virtual Build Environment
prevent
Quality • Security • Concurrency
extend
Custom Checks
Analysis Engine
• Interprocedural
• Dataflow Analysis
• Statistical Analysis
• False Path Pruning
• 100% of All Paths
• Incremental Analysis
Defect Manager
• Report Dashboard
• Developer Dashboard
• IDE Support
• Open Standard Interfaces

Demo

- Coverity video: http://youtu.be/_Vt4niZfNeA
- Observations
  - Code analysis integrated into development workflow
  - Program context important: analysis involves sequence of function calls, surrounding statements
  - This is a sales video: no discussion of false alarms

Bugs to Detect

Some examples

- Crash Causing Defects
- Null pointer dereference
- Use after free
- Double free
- Array indexing errors
- Mismatched array new/delete
- Potential stack overrun
- Potential heap overrun
- Return pointers to local variables
- Logically inconsistent code
- Uninitialized variables
- Invalid use of negative values
- Passing large parameters by value
- Underallocations of dynamic data
- Memory leaks
- File handle leaks
- Network resource leaks
- Unused values
- Unhandled return codes
- Use of invalid iterators

Slide credit: Andy Chou

Example: Check for missing optional args

- **Prototype for open() syscall:**

```
int open(const char *path, int oflag, /* mode_t mode */...);
```

- **Typical mistake:**

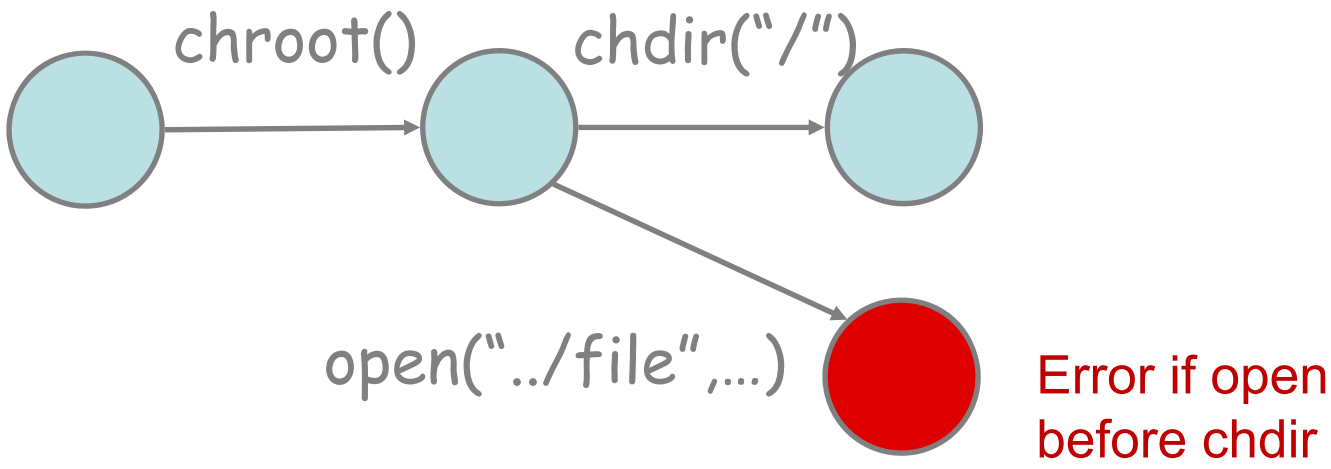
```
fd = open("file", O_CREAT);
```

- **Result: file has random permissions**

- **Check: Look for oflags == O_CREAT without mode argument**

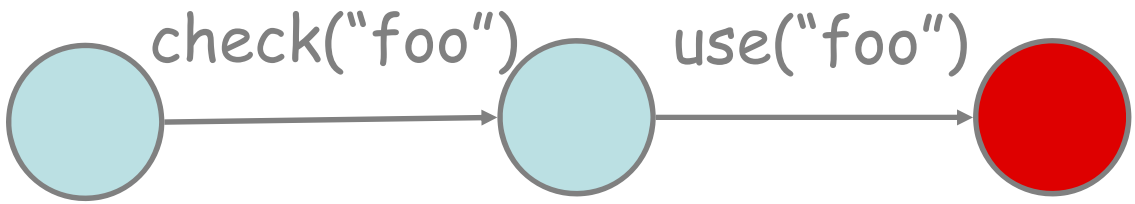
Example: Chroot protocol checker

- **Goal: confine process to a "jail" on the filesystem**
  - chroot() changes filesystem root for a process
- **Problem**
  - chroot() itself does not change current working directory

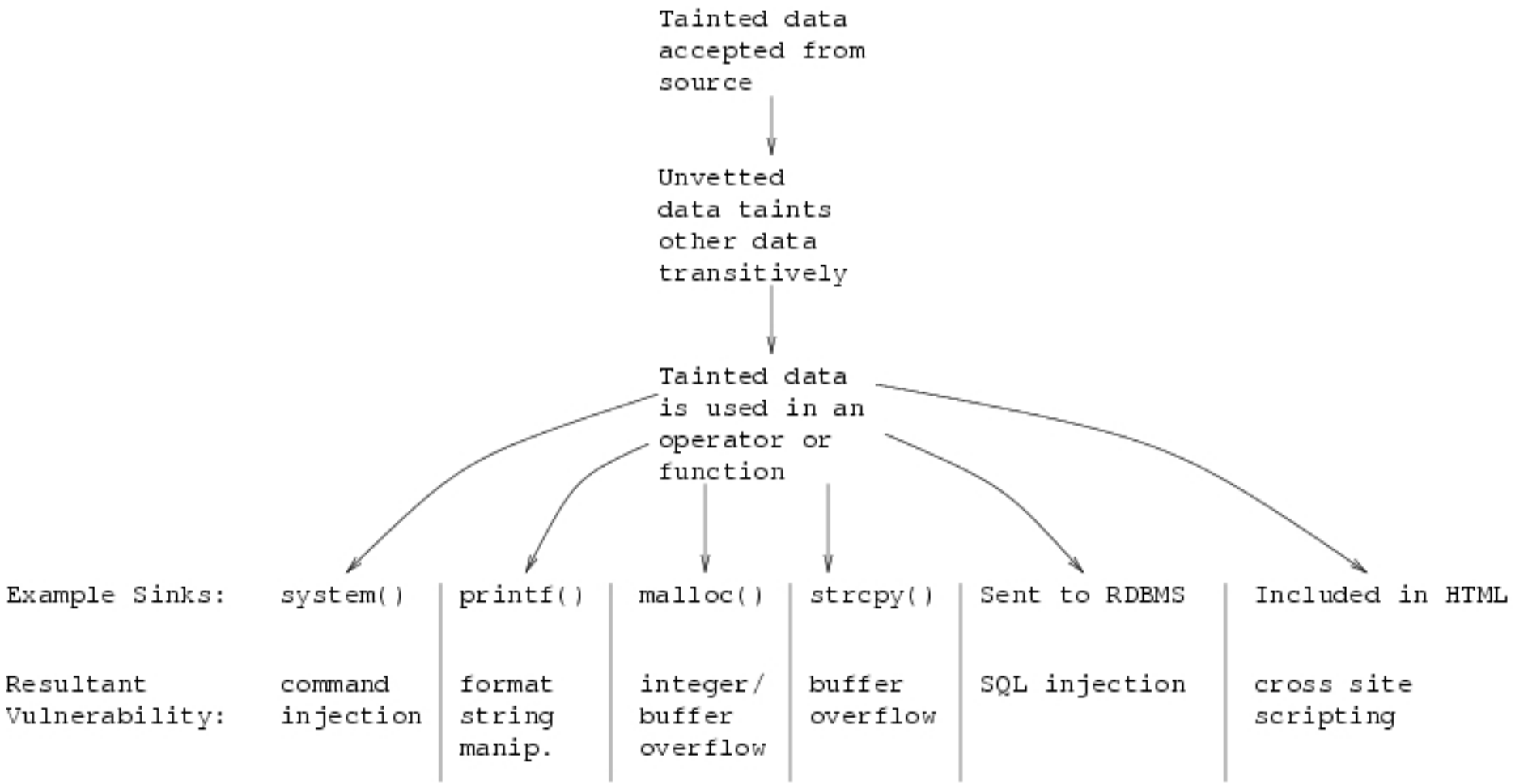


TOCTOU

- **Race condition between time of check and use**
- **Not applicable to all programs**



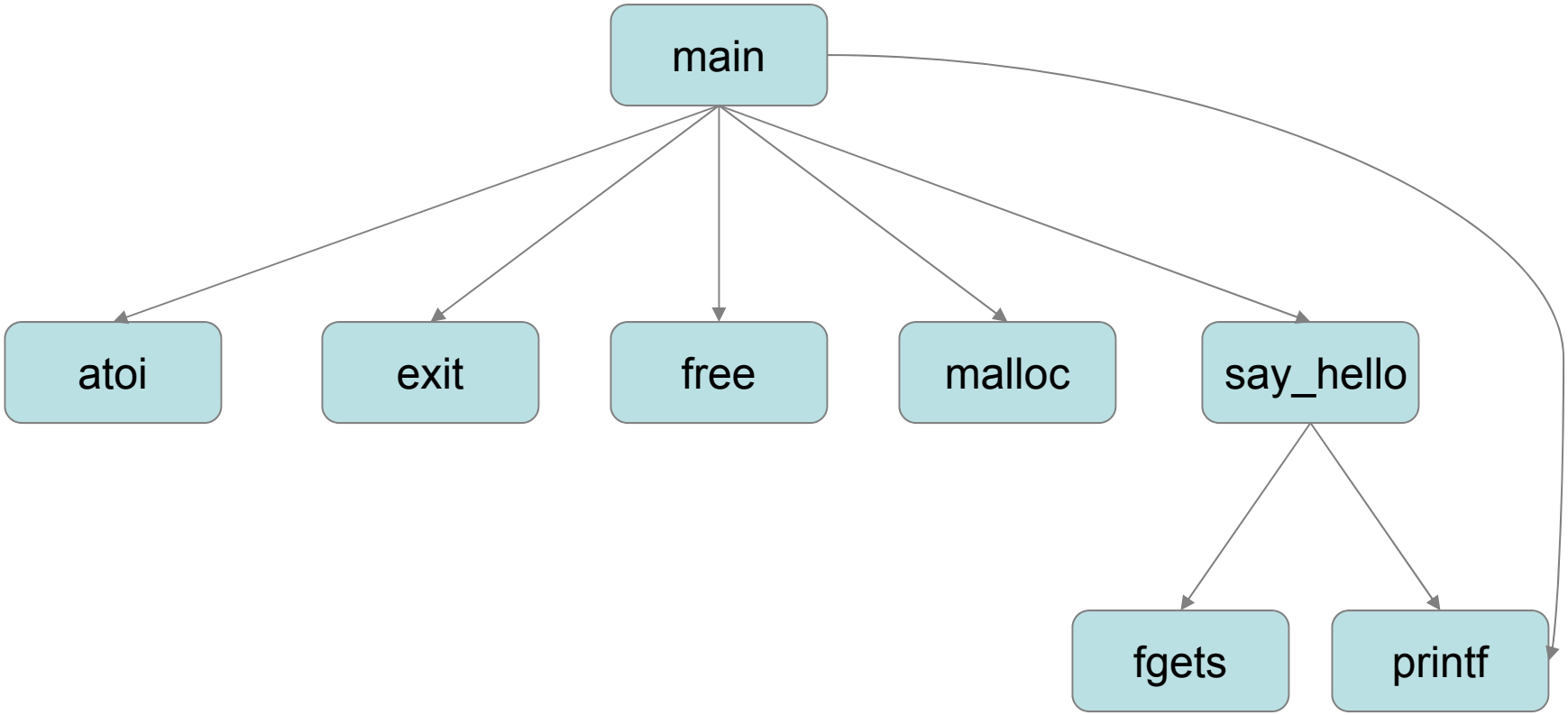
Tainting checkers



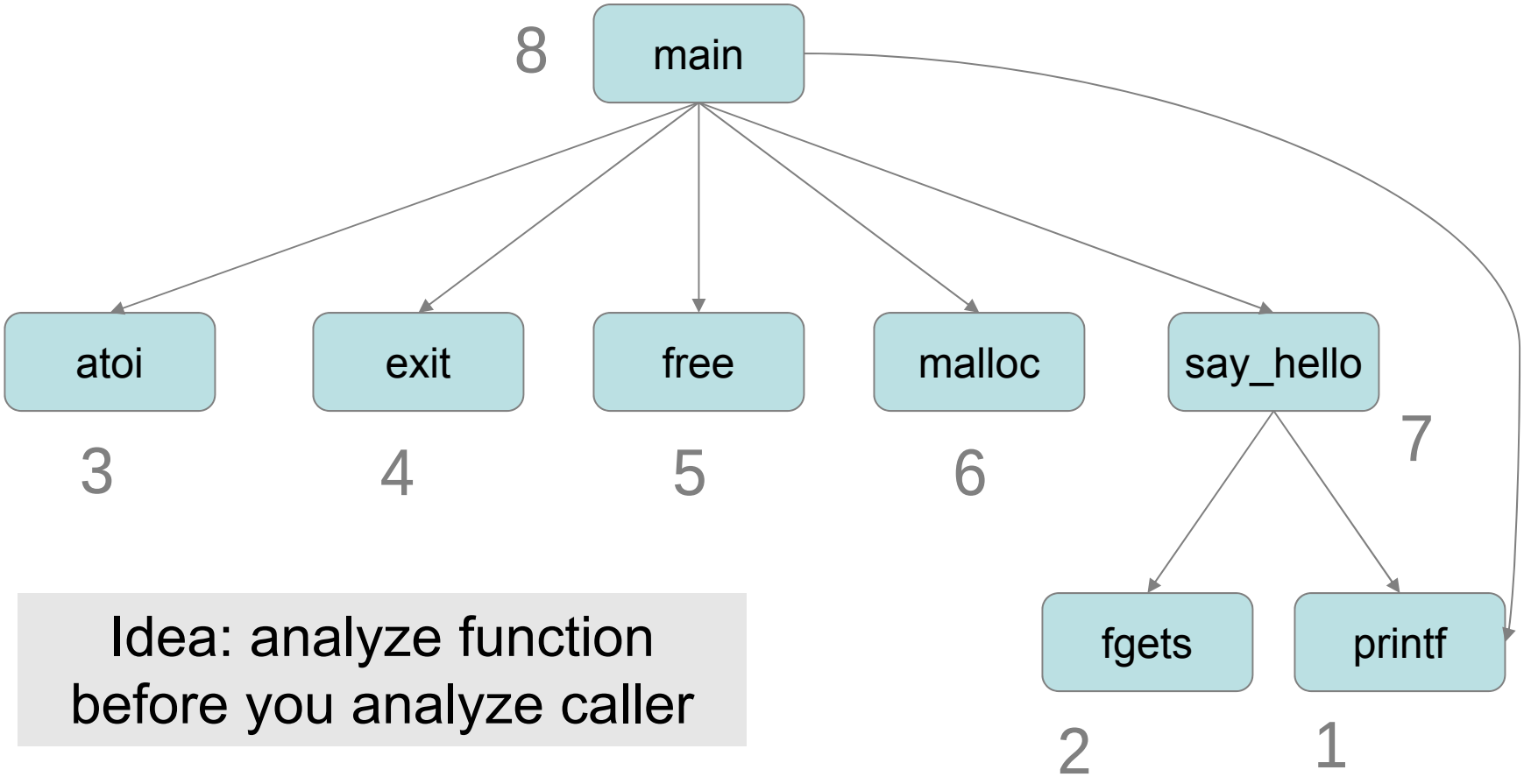
Example code with function def, calls

```
#include <stdlib.h>
#include <stdio.h>

void say_hello(char * name, int size) {
  printf("Enter your name: ");
  fgets(name, size, stdin);
  printf("Hello %s.\n", name);
}

int main(int argc, char *argv[]) {
  if (argc != 2) {
    printf("Error, must provide an input buffer size.\n");
    exit(-1);
  }
  int size = atoi(argv[1]);
  char * name = (char*)malloc(size);
  if (name) {
    say_hello(name, size);
    free(name);
  } else {
    printf("Failed to allocate %d bytes.\n", size);
  }
}
```

Callgraph

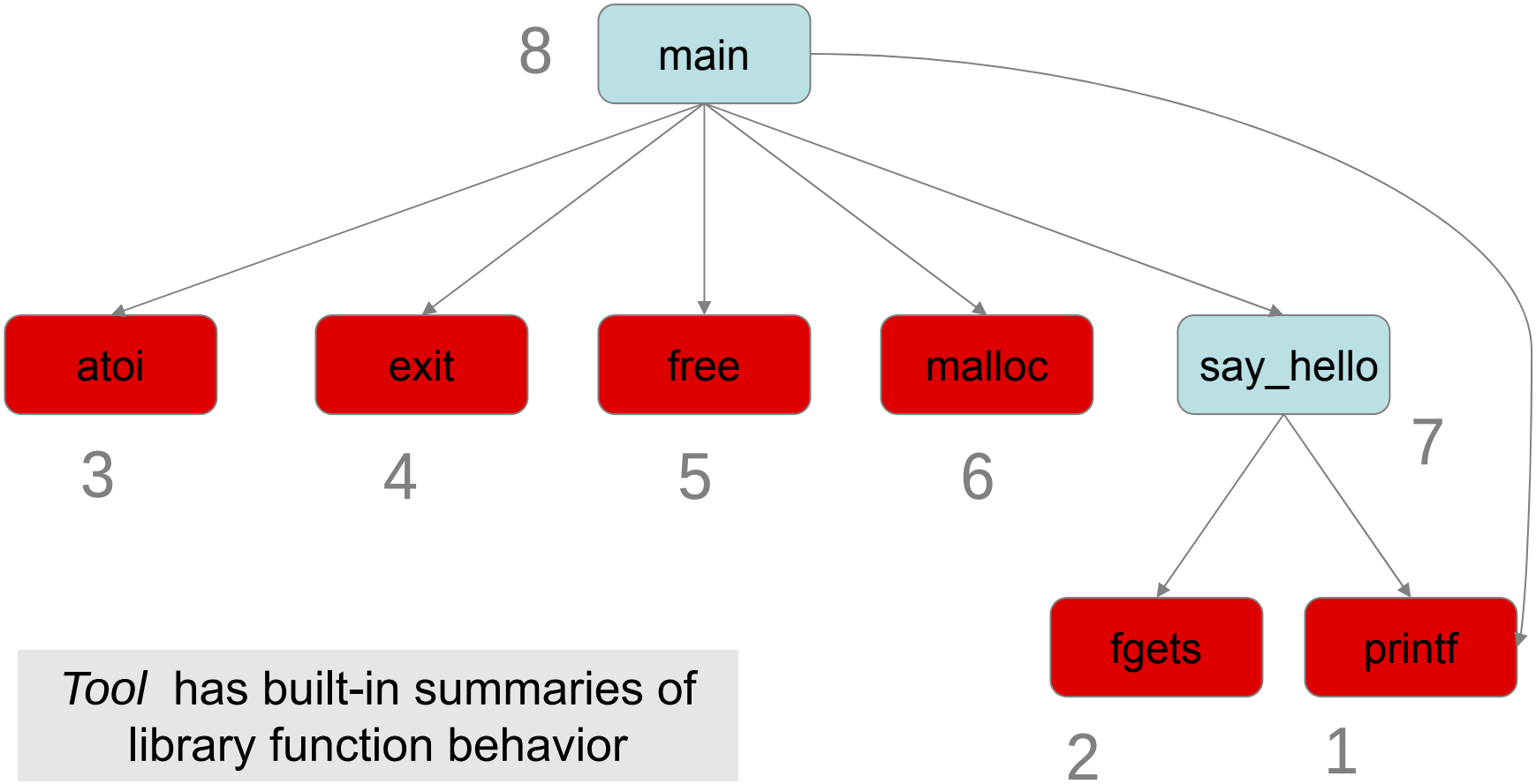


Reverse Topological Sort



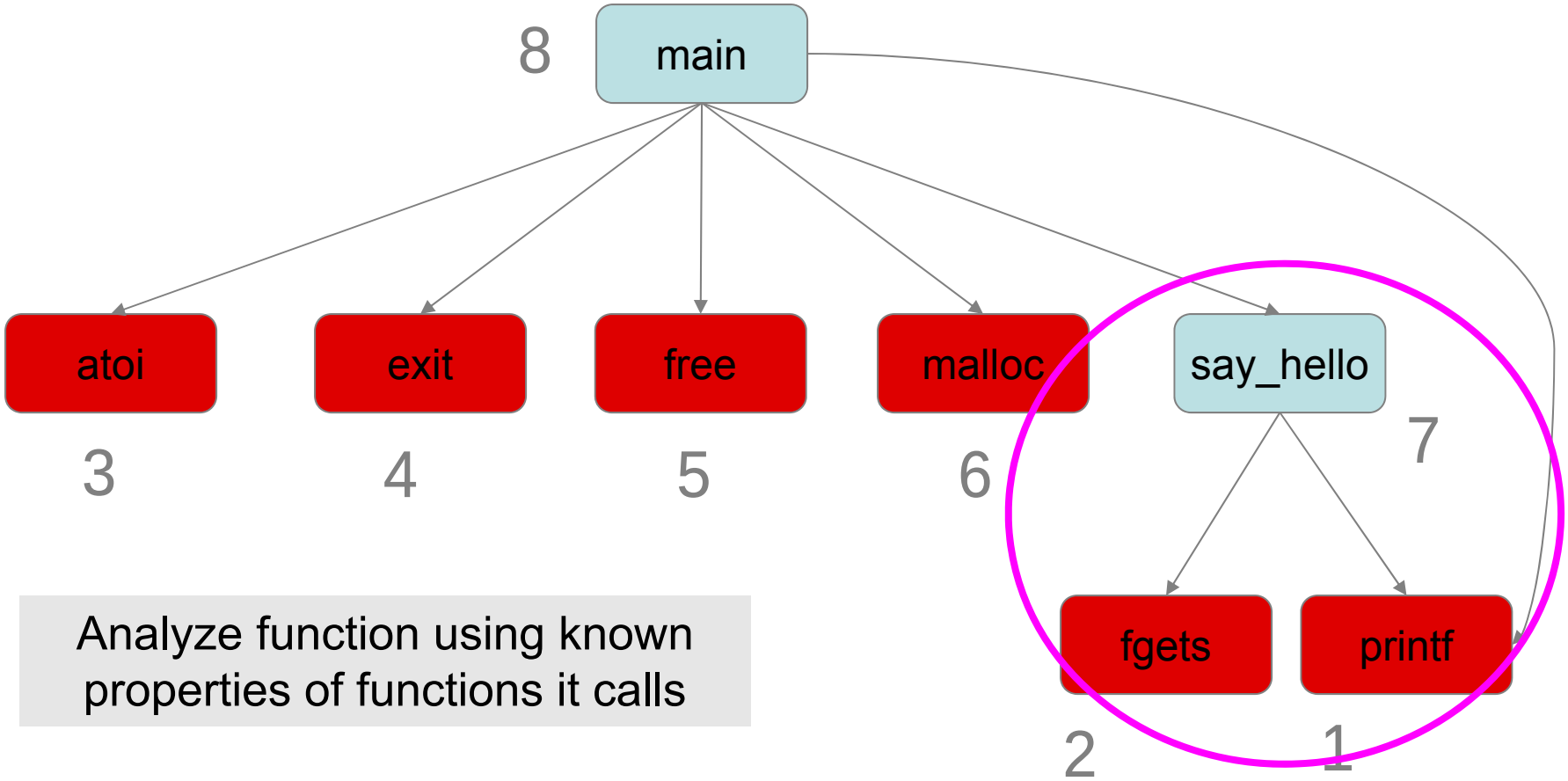
Idea: analyze function before you analyze caller

Apply Library Models



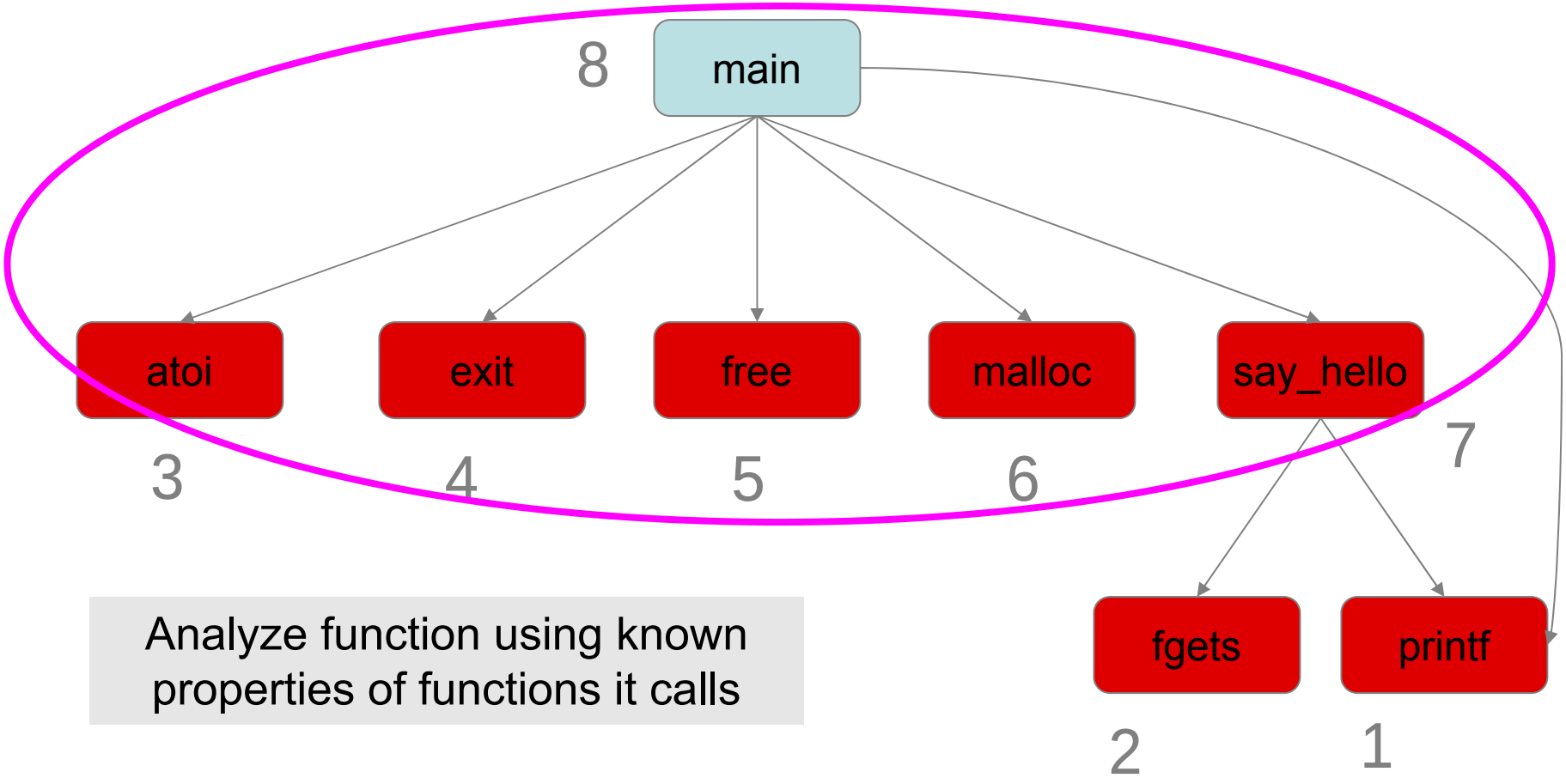
Tool has built-in summaries of library function behavior

Bottom Up Analysis



Analyze function using known properties of functions it calls

Bottom Up Analysis



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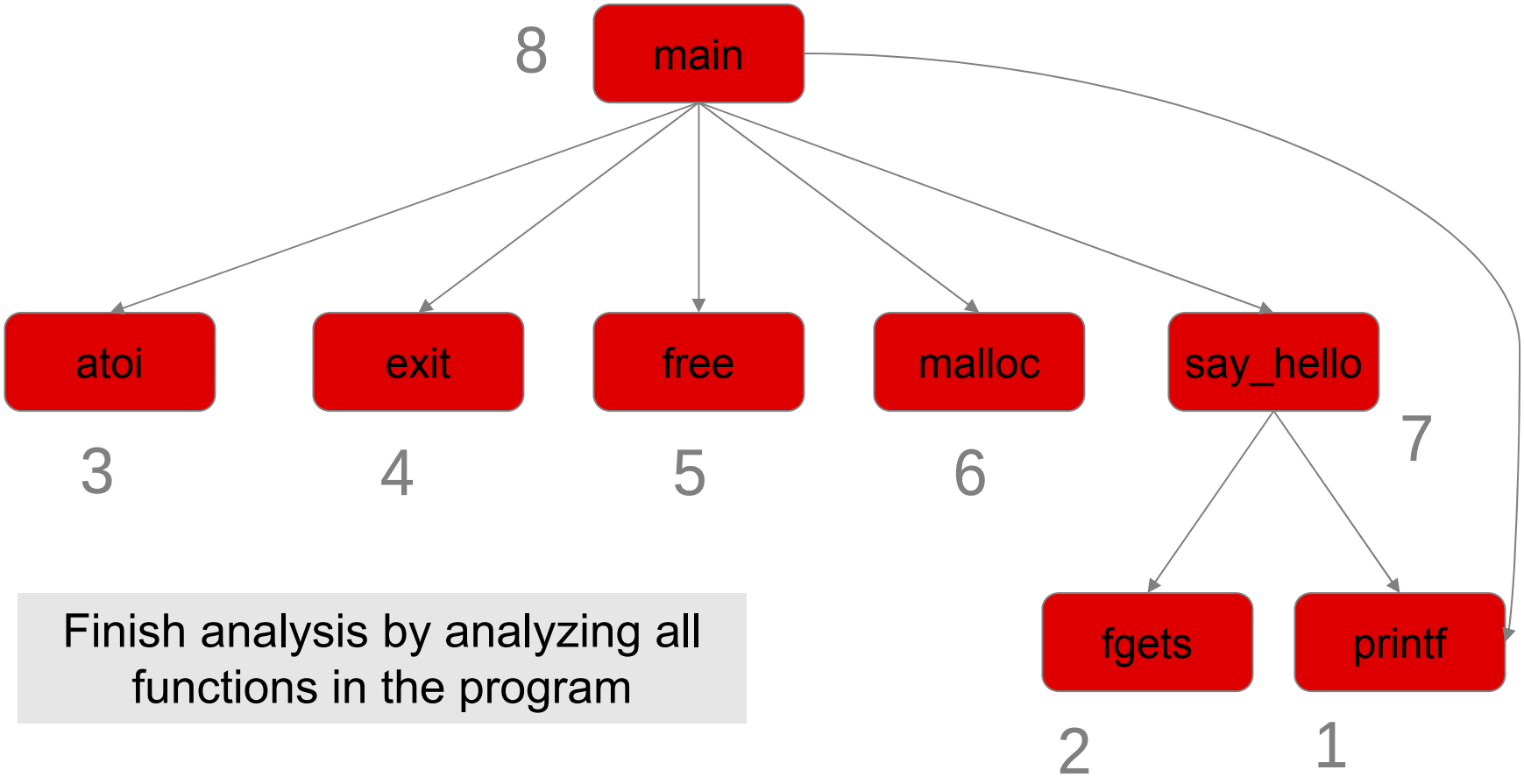
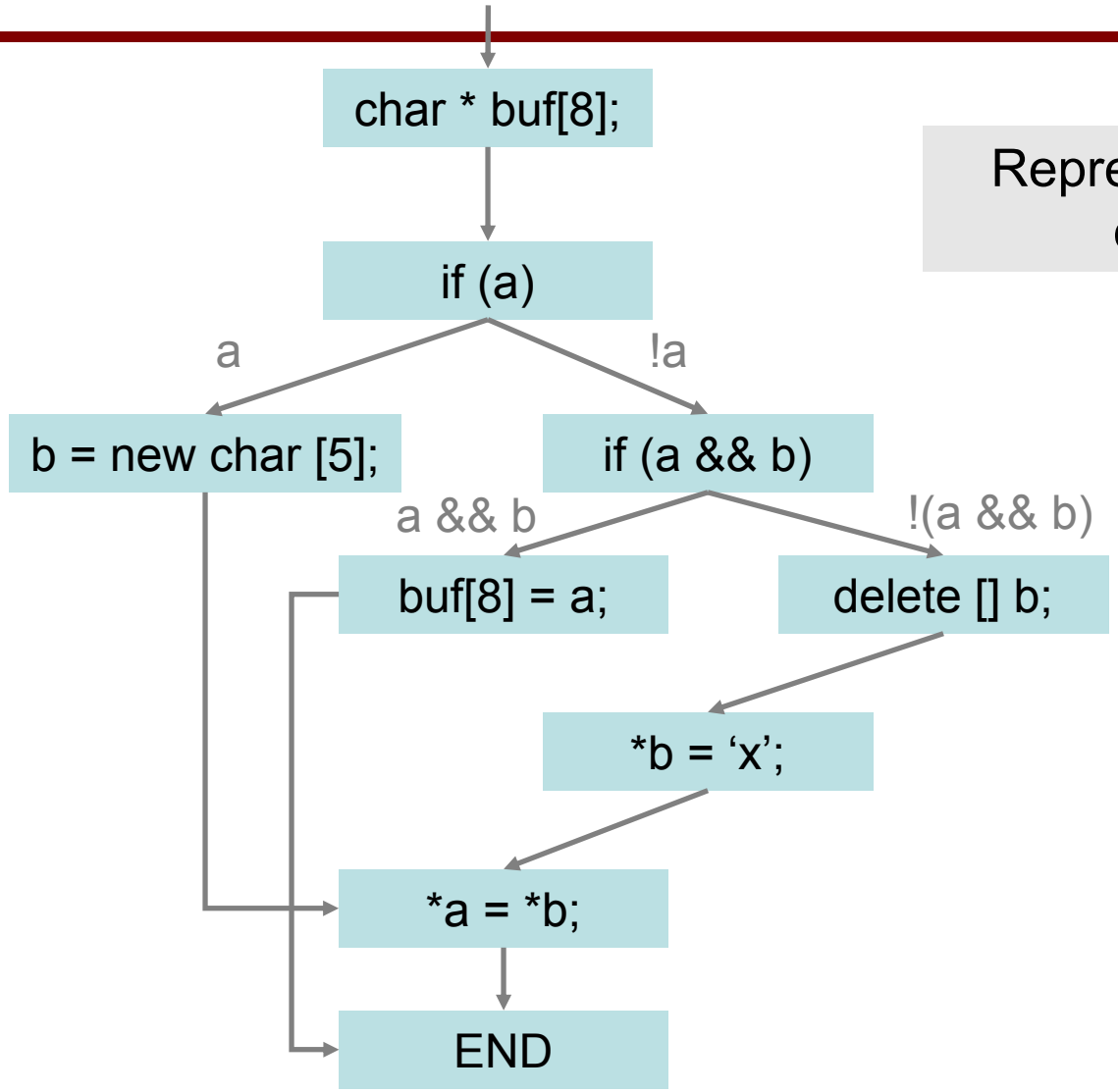


Finish analysis by analyzing all functions in the program

Finding Local Bugs

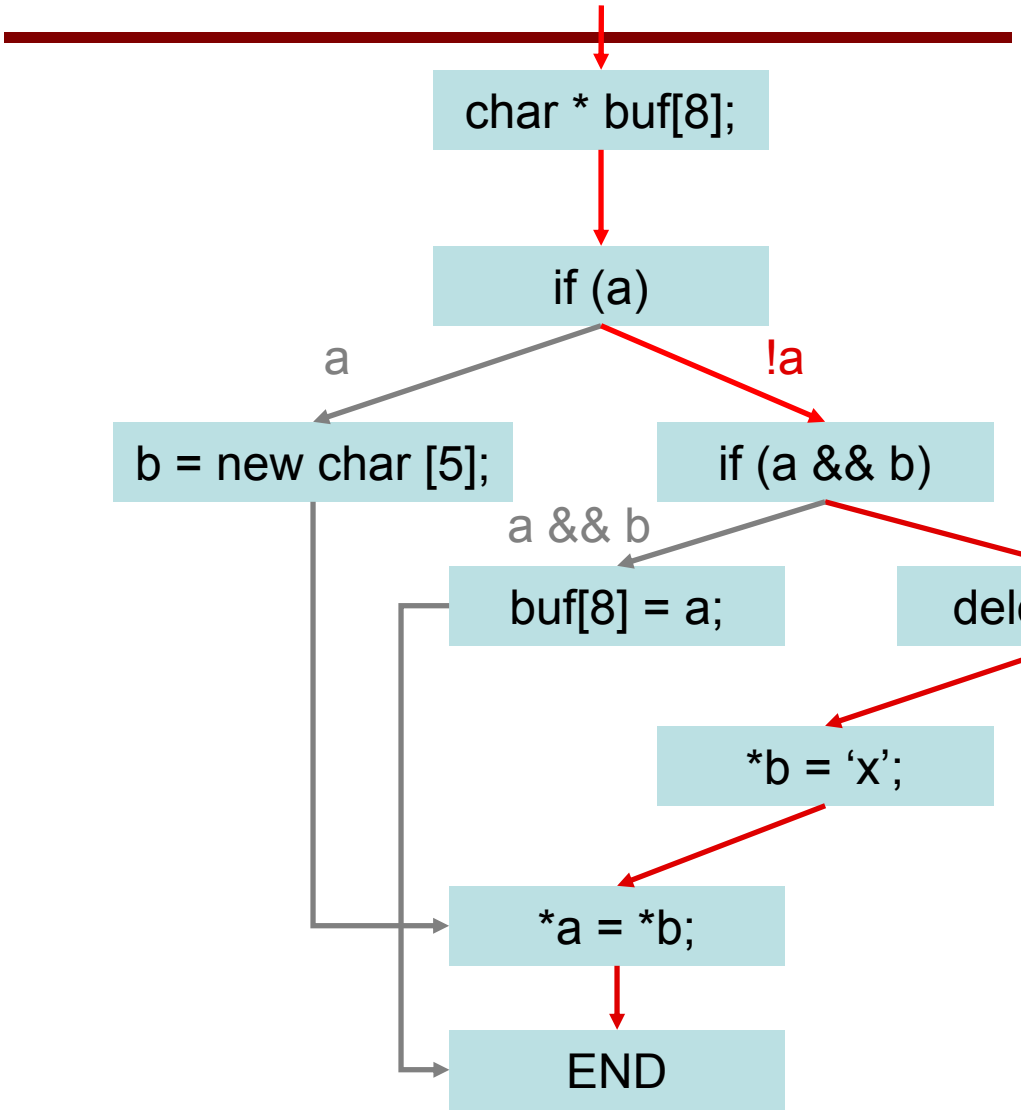
```
#define SIZE 8
void set_a_b(char * a, char * b) {
 char * buf[SIZE];
 if (a) {
     b = new char[5];
 } else {
     if (a && b) {
      buf[SIZE] = a;
      return;
     } else {
      delete [] b;
     }
     *b = ‘x’;
 }
 *a = *b;
}
```

Control Flow Graph



Represent logical structure of code in graph form

Path Traversal

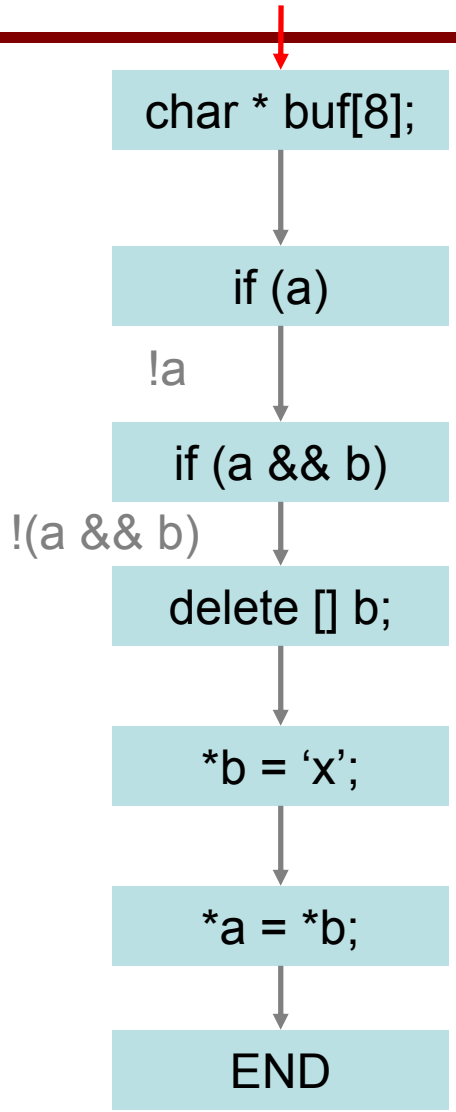


Conceptually Analyze each path through control graph separately

Actually Perform some checking computation once per node; combine paths at merge nodes

Apply Checking

Null pointers Use after free Array overrun



See how three checkers are run for this path

Checker

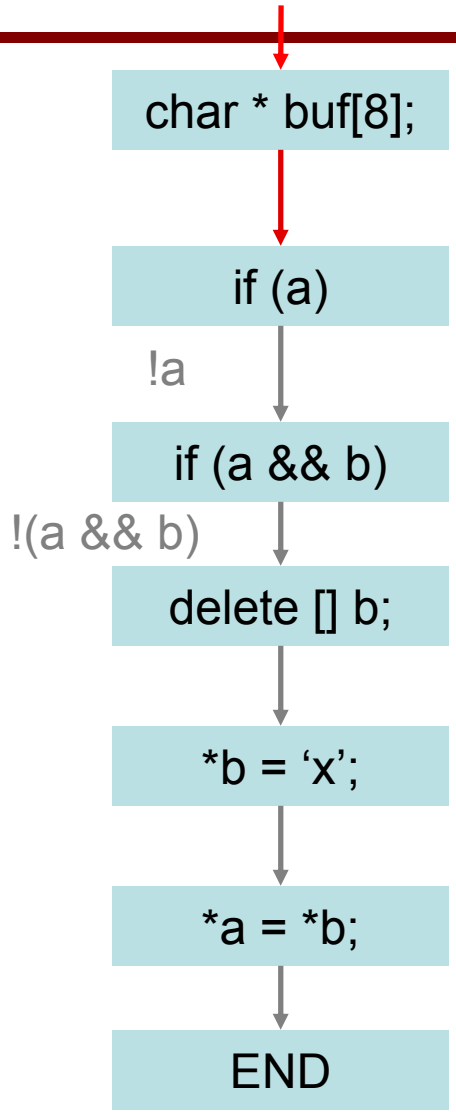
- Defined by a state diagram, with state transitions and error states

Run Checker

- Assign initial state to each program var
- State at program point depends on state at previous point, program actions
- Emit error if error state reached

Apply Checking

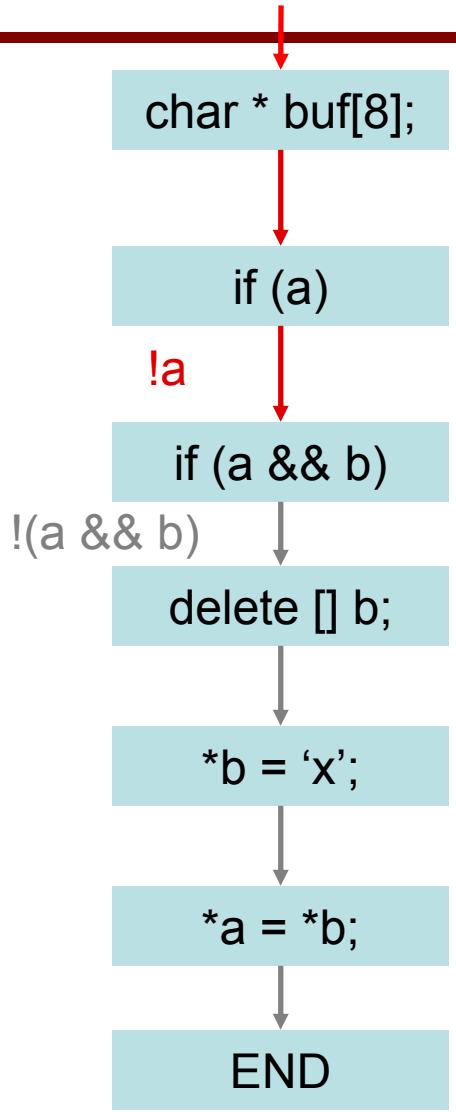
Null pointers Use after free Array overrun



"buf is 8 bytes"

Apply Checking

Null pointers Use after free Array overrun

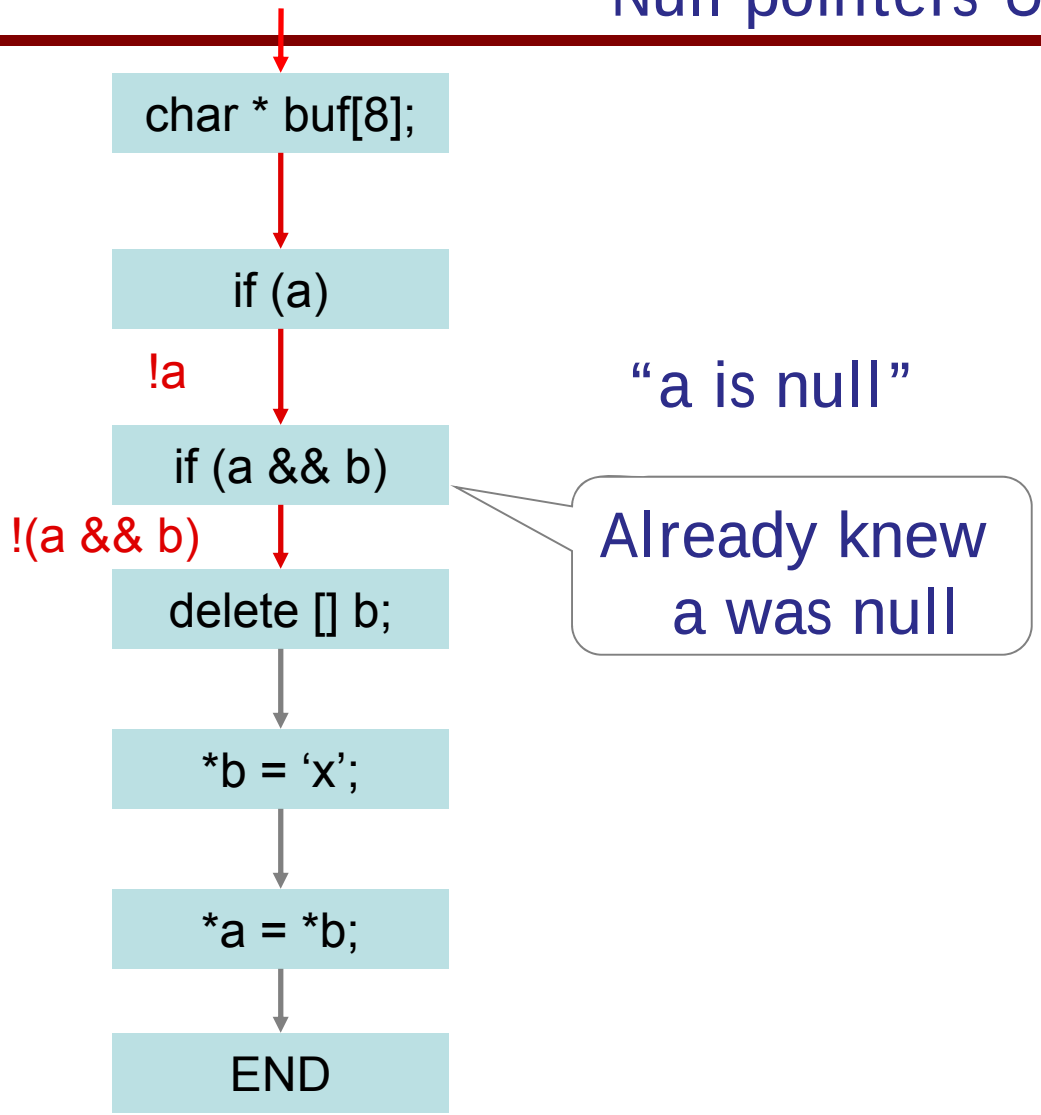


"buf is 8 bytes"

"a is null"

Apply Checking

Null pointers Use after free Array overrun



"buf is 8 bytes"

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Apply Checking

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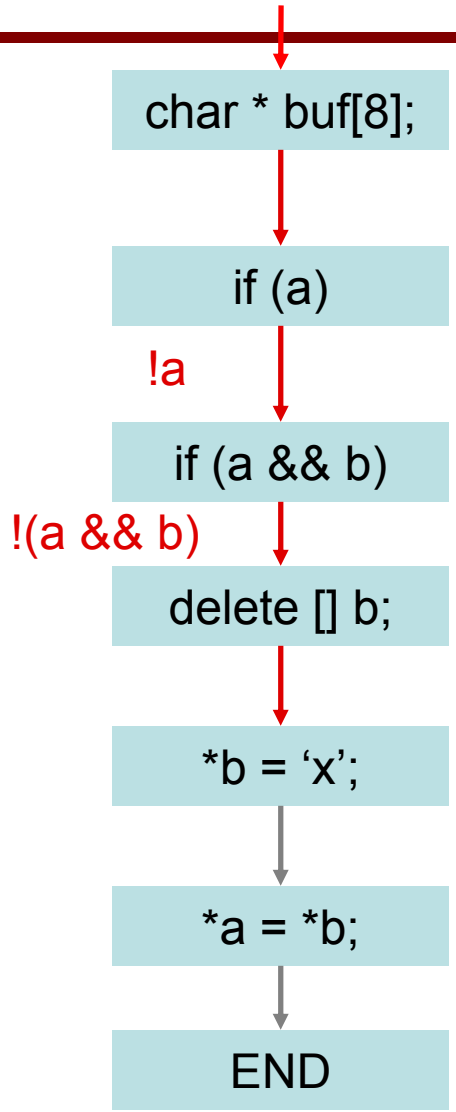


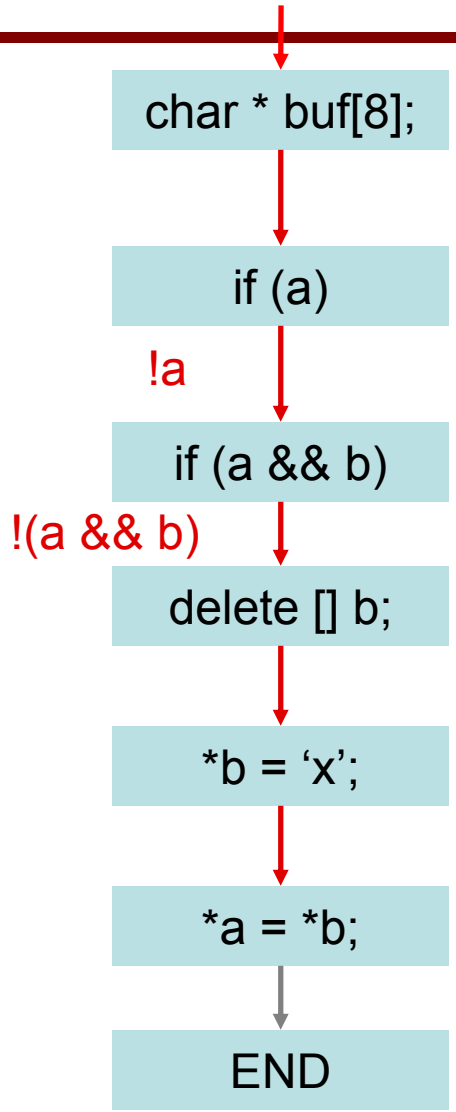
"buf is 8 bytes"

"a is null"

"b is deleted"

Apply Checking

Null pointers Use after free Array overrun



"buf is 8 bytes"

"a is null"

"b is deleted"

"b dereferenced!"

Apply Checking

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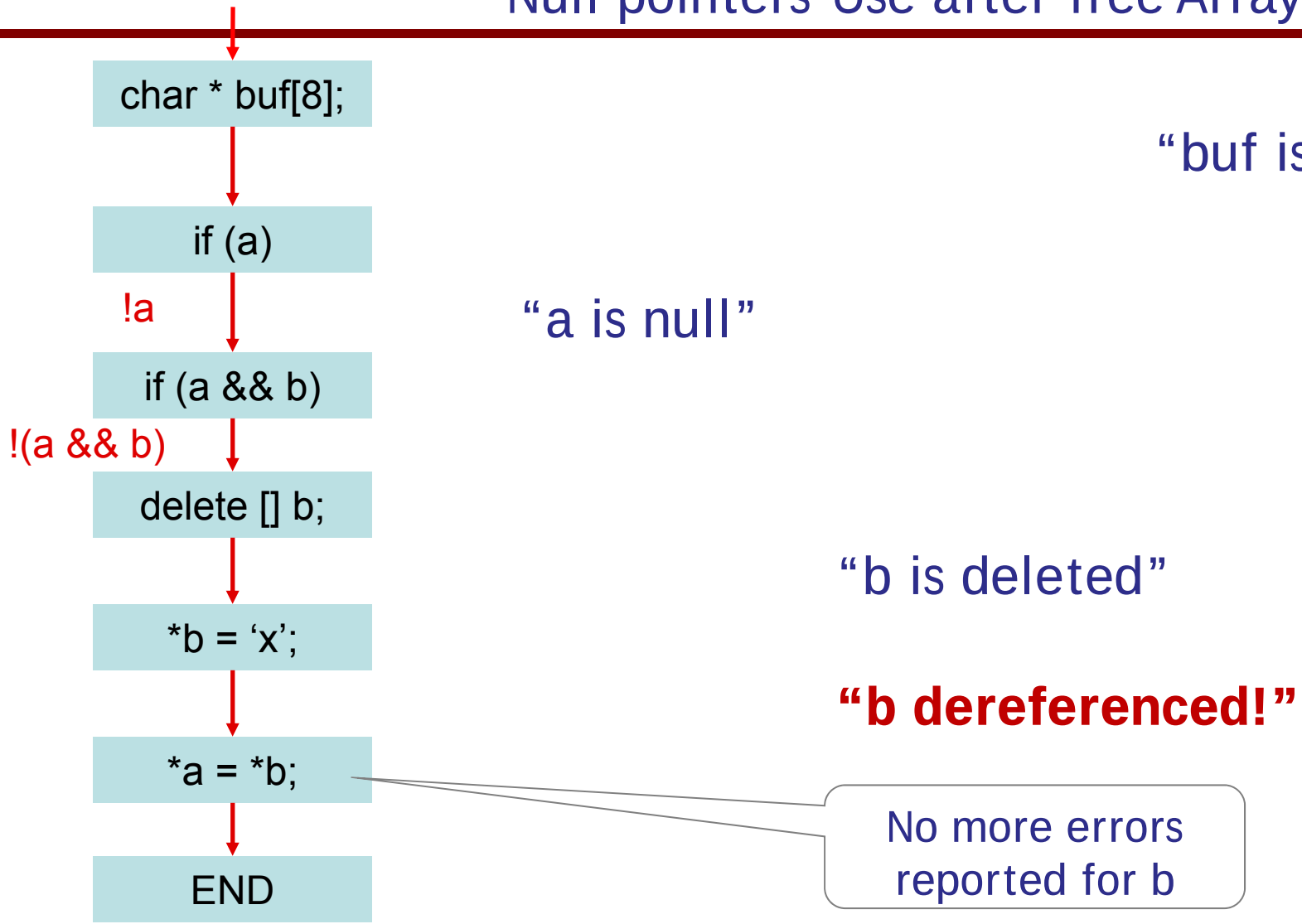
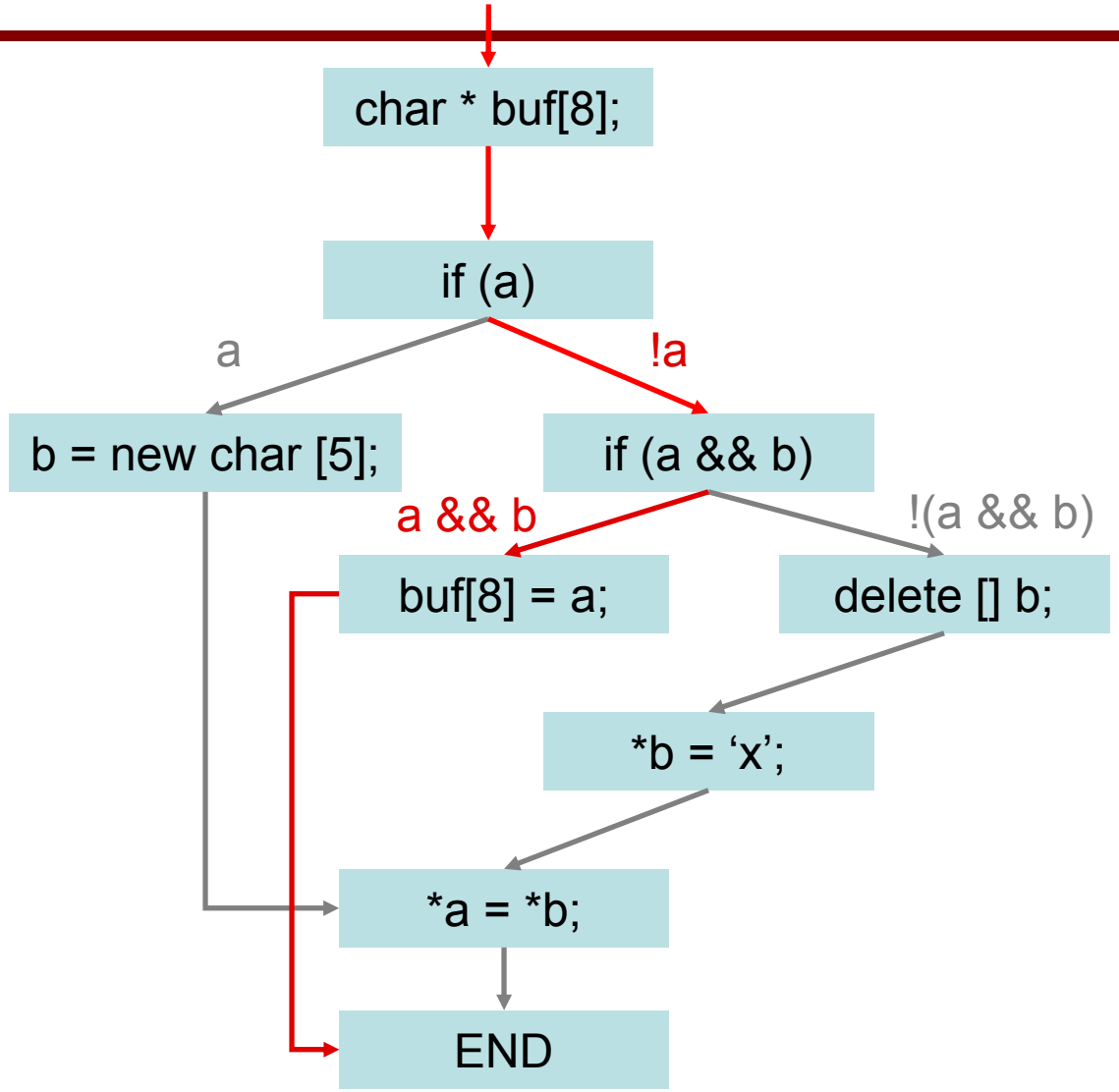


False Positives

- **What is a bug? Something the user will fix.**

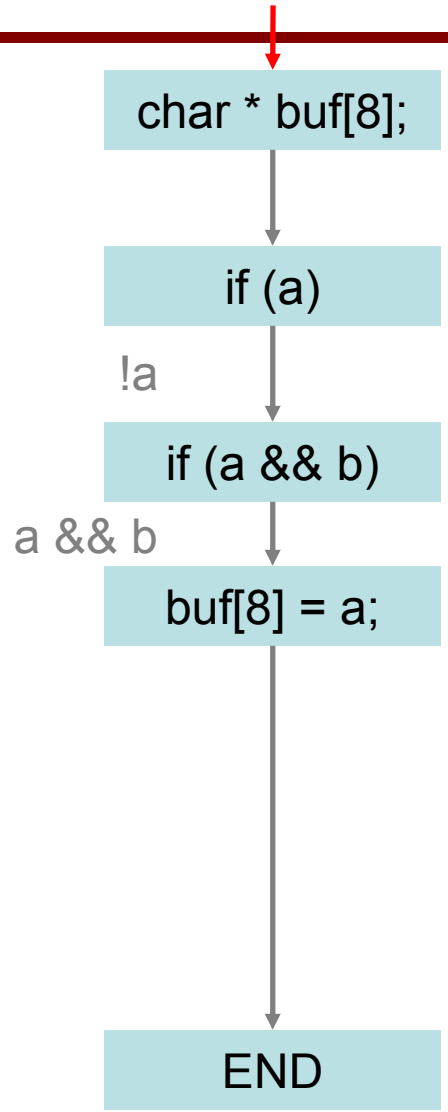
- **Many sources of false positives**
  - False paths
  - Idioms
  - Execution environment assumptions
  - Killpaths
  - Conditional compilation
  - “third party code”
  - Analysis imprecision
  - …

A False Path

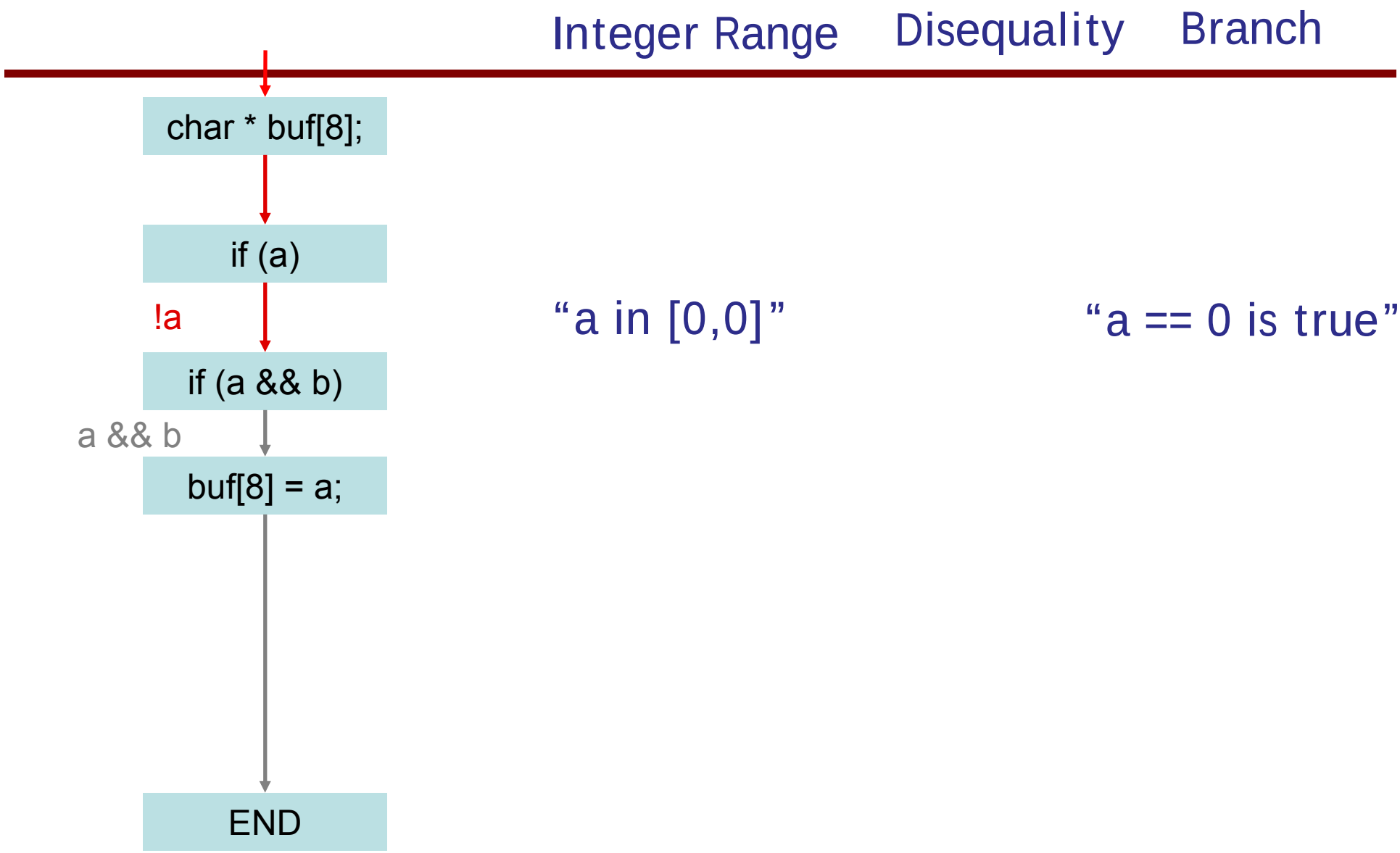


False Path Pruning

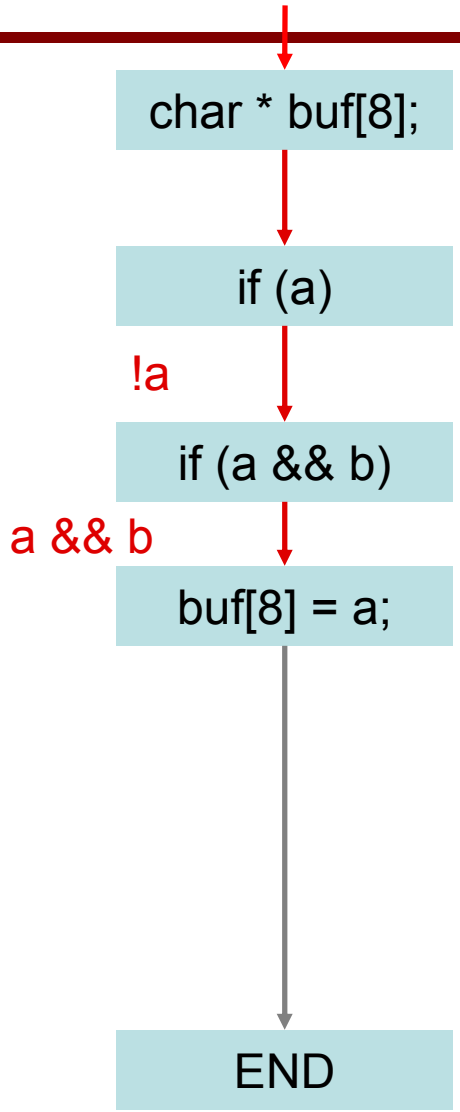
Integer Range Disequality Branch



False Path Pruning



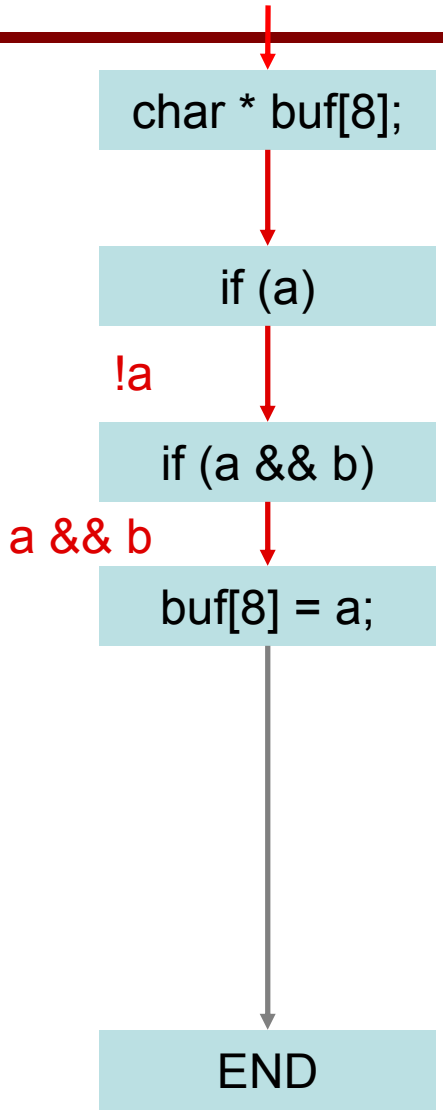
False Path Pruning



Integer Range	Disequality	Branch
“a in [0,0]”		“a == 0 is true”
	“a != 0”	

False Path Pruning

Integer Range | Disequality | Branch



Impossible

"a in [0,0]"

"a == 0 is true"

"a != 0"

Environment Assumptions

- **Should the return value of malloc() be checked?**

```
int *p = malloc(sizeof(int));
*p = 42;
```

OS Kernel:
Crash machine.

File server:
Pause filesystem.

Web application:
200ms downtime

Spreadsheet:
Lose unsaved changes.

Game:
Annoy user.

IP Phone:
Annoy user.

Library:
?

Medical device:
malloc?!

Statistical Analysis

- **Assume the code is usually right**

3/4 deref

```
int *p = malloc(sizeof(int));
*p = 42;
```

```
int *p = malloc(sizeof(int));
*p = 42;
```

```
int *p = malloc(sizeof(int));
*p = 42;
```

```
int *p = malloc(sizeof(int));
if(p)  *p = 42;
```

```
int *p = malloc(sizeof(int));
if(p) *p = 42;
```

```
int *p = malloc(sizeof(int));
if(p) *p = 42;
```

```
int *p = malloc(sizeof(int));
if(p) *p = 42;
```

```
int *p = malloc(sizeof(int));
*p = 42;
```

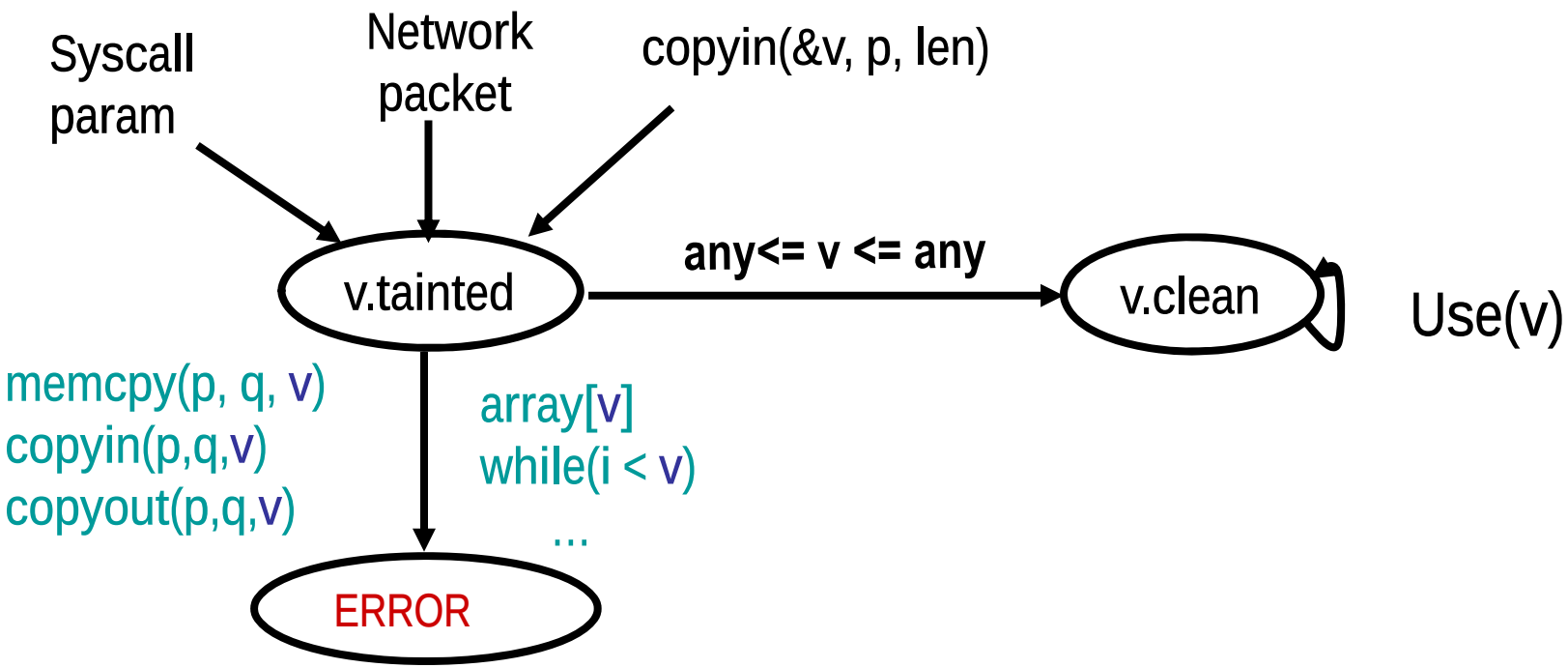
1/4 deref

Application to Security Bugs

- **Stanford research project**
  - Ken Ashcraft and Dawson Engler, Using Programmer-Written Compiler Extensions to Catch Security Holes, IEEE Security and Privacy 2002
  - Used modified compiler to find over 100 security holes in Linux and BSD
  - http://www.stanford.edu/~engler/
- **Benefit**
  - Capture recommended practices, known to experts, in tool available to all

Sanitize integers before use

Warn when unchecked integers from untrusted sources reach trusting sinks



Linux: 125 errors, 24 false; BSD: 12 errors, 4 false

Example security holes

- **Remote exploit, no checks**

```
/* 2.4.9/drivers/isdn/act2000/capi.c:actcapi_dispatch */
isdn_ctrl cmd;
...
while ((skb = skb_dequeue(&card->rcvq))) {
   msg = skb->data;
   ...
   memcpy(cmd.parm.setup.phone,
          msg->msg.connect_ind.addr.num,
          msg->msg.connect_ind.addr.len - 1);
```

Example security holes

- **Missed lower-bound check:**

```
/* 2.4.5/drivers/char/drm/i810_dma.c */

if(copy_from_user(&d, arg, sizeof(arg)))
   return –EFAULT;
if(d.idx > dma->buf_count)
   return –EINVAL;
buf = dma->buflist[d.idx];
Copy_from_user(buf_priv->virtual, d.address, d.used);
```

User-pointer inference

- **Problem: which are the user pointers?**
  - Hard to determine by dataflow analysis
  - Easy to tell if kernel *believes* pointer is from user!
- **Belief inference**
  - "*p" implies safe kernel pointer
  - "copyin(p)/copyout(p)" implies dangerous user ptr
  - Error: pointer p has both beliefs.
- **Implementation: 2 pass checker**

  inter-procedural: compute all tainted pointers

  local pass to check that they are not dereferenced

Results for BSD and Linux

- **All bugs released to implementers; most serious fixed**

Violation	Linux Bug	Linux Fixed	BSD Bug	BSD Fixed
Gain control of system	18	15	3	3
Corrupt memory	43	17	2	2
Read arbitrary memory	19	14	7	7
Denial of service	17	5	0	0
Minor	28	1	0	0
Total	125	52	12	12

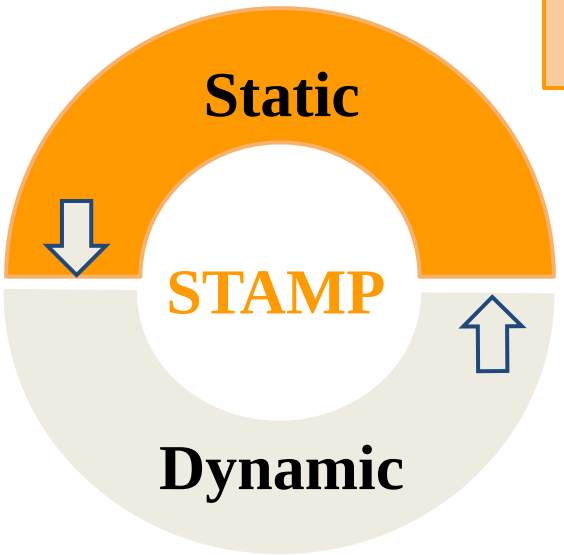
Outline

- General discussion of tools
  - Goals and limitations
  - Approach based on abstract states
- More about one specific approach
  - Property checkers from Engler et al., Coverity
  - Sample security-related results
- ★ Static analysis for Android malware
  - …

Slides from: S. Bugrahe, A. Chou, I&T Dillig, D. Engler, J. Franklin, A. Aiken, …

STAMP Admission System

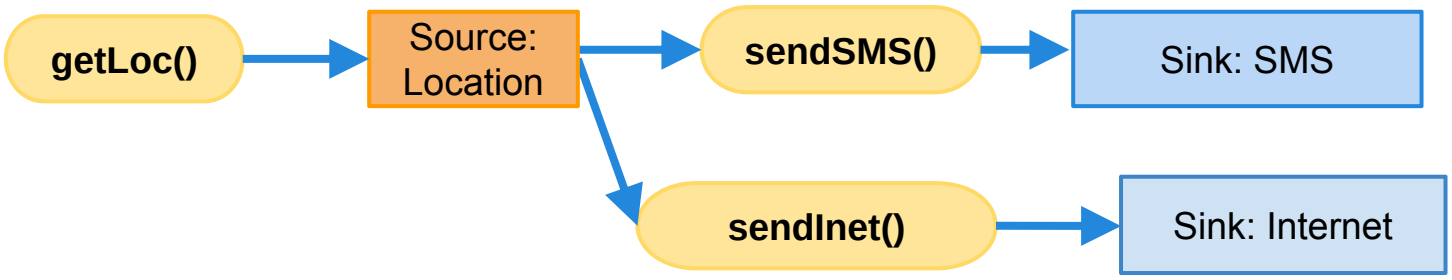
Static Analysis
More behaviors,
fewer details



Dynamic Analysis
Fewer behaviors,
more details

Alex Aiken,
John Mitchell,
Saswat Anand,
Jason Franklin
Osbert Bastani,
Lazaro Clapp,
Patrick Mutchler,
Manolis Papadakis

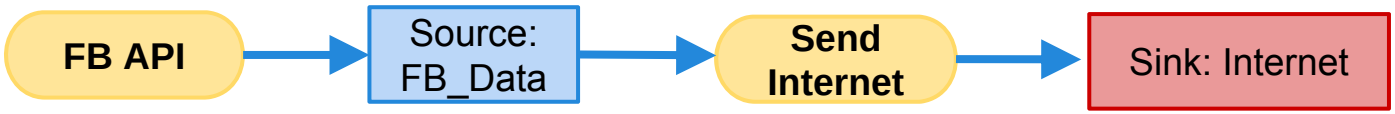
Data Flow Analysis



- Source-to-sink flows
  - Sources: Location, Calendar, Contacts, Device ID etc.
  - Sinks: Internet, SMS, Disk, etc.

Applications of Data Flow Analysis

- Malware/Greyware Analysis
  - Data flow summaries enable enterprise-specific policies
- API Misuse and Data Theft Detection



- Automatic Generation of App Privacy Policies
  - Avoid liability, protect consumer privacy

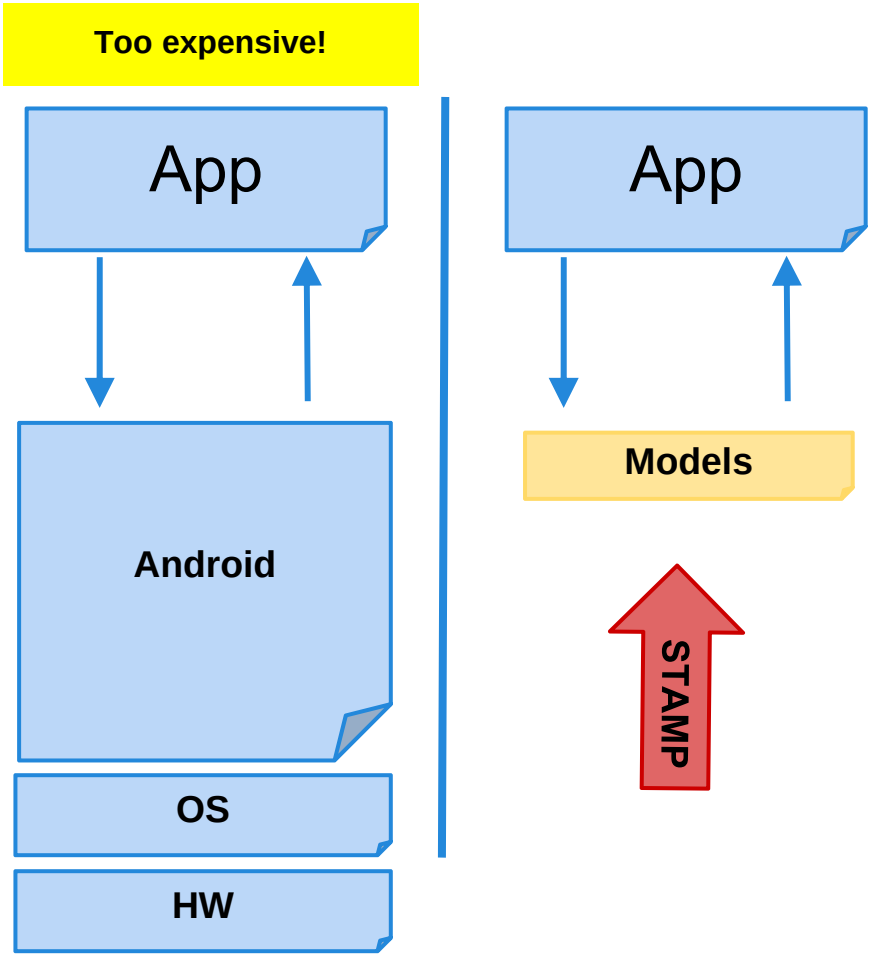
Privacy Policy
This app collects your:
Contacts
Phone Number
Address

- Vulnerability Discovery

Challenges

- Android is 3.4M+ lines of complex code
  - Uses reflection, callbacks, native code
- **Scalability:** Whole system analysis impractical
- **Soundness:** Avoid missing flows
- **Precision:** Minimize false positives

STAMP Approach



- Model Android/Java
  - Sources and sinks
  - Data structures
  - Callbacks
  - 500+ models
- Whole-program analysis
  - Context sensitive

Building Models

- 30k+ methods in Java/Android API
  - 5 mins x 30k = 2500 hours

- Follow the permissions
  - 20 permissions for sensitive sources
    - ACCESS_FINE_LOCATION (8 methods with source annotations)
    - READ_PHONE_STATE - (9 methods)
  - 4 permissions for sensitive sinks
    - INTERNET, SEND_SMS, etc.

Identifying Sensitive Data

```
android.Telephony.TelephonyManager: String getDeviceId()
```

- Returns device IMEI in String
- Requires permission GET_PHONE_STATE

```
@STAMP(
   SRC ="$GET_PHONE_STATE.deviceid",
   SINK ="@return"
)
```

Data We Track (Sources)

- Account data
- Audio
- Calendar
- Call log
- Camera
- Contacts
- Device Id
- Location
- Photos (Geotags)
- SD card data
- SMS

30+ types of sensitive data

Data Destinations (Sinks)

- Internet (socket)
- SMS
- Email
- System Logs
- Webview/Browser
- File System
- Broadcast Message

10+ types of exit points

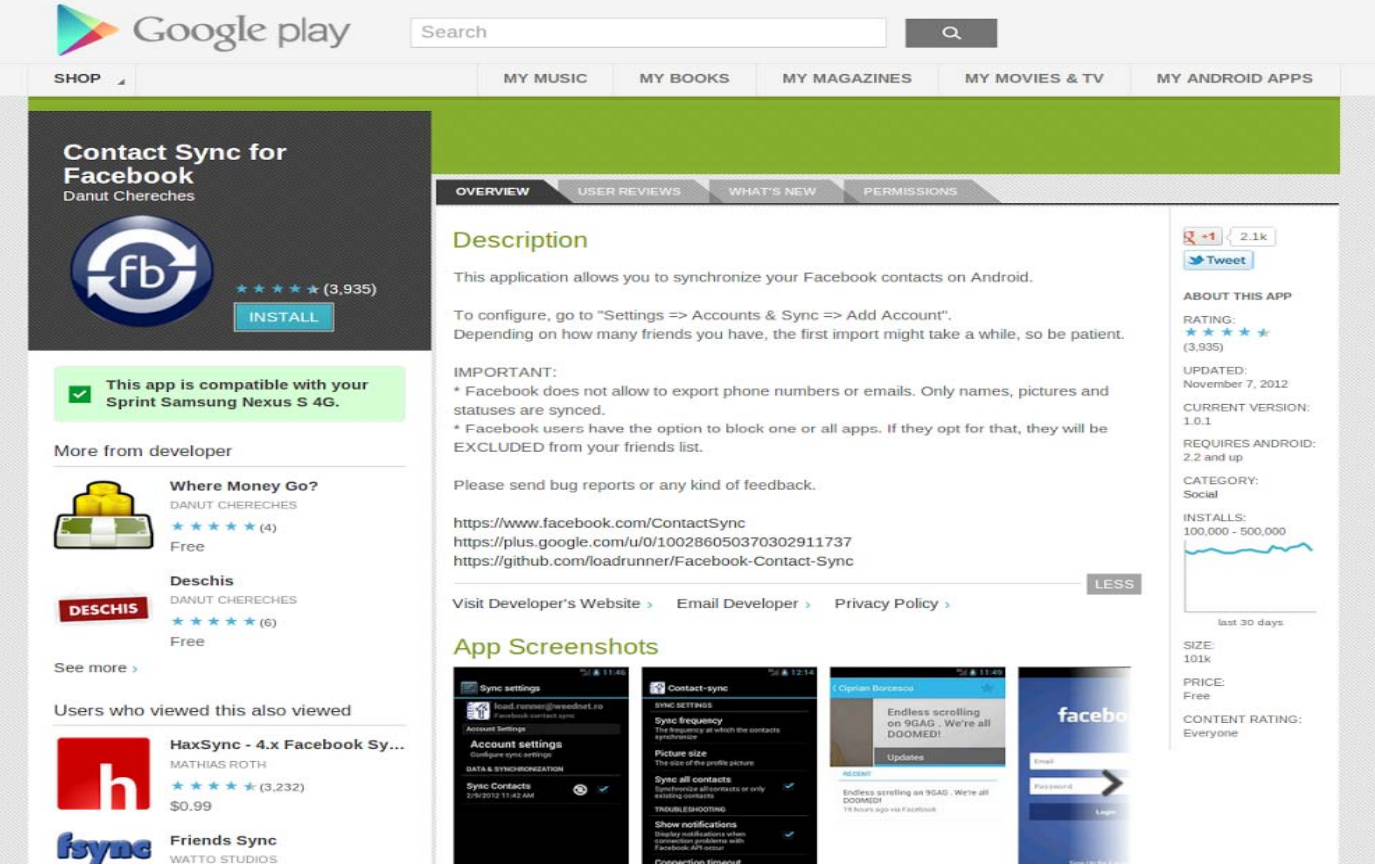
Currently Detectable Flow Types

396 Flow Types

Unique Flow Types = Sources x Sink

Example Analysis

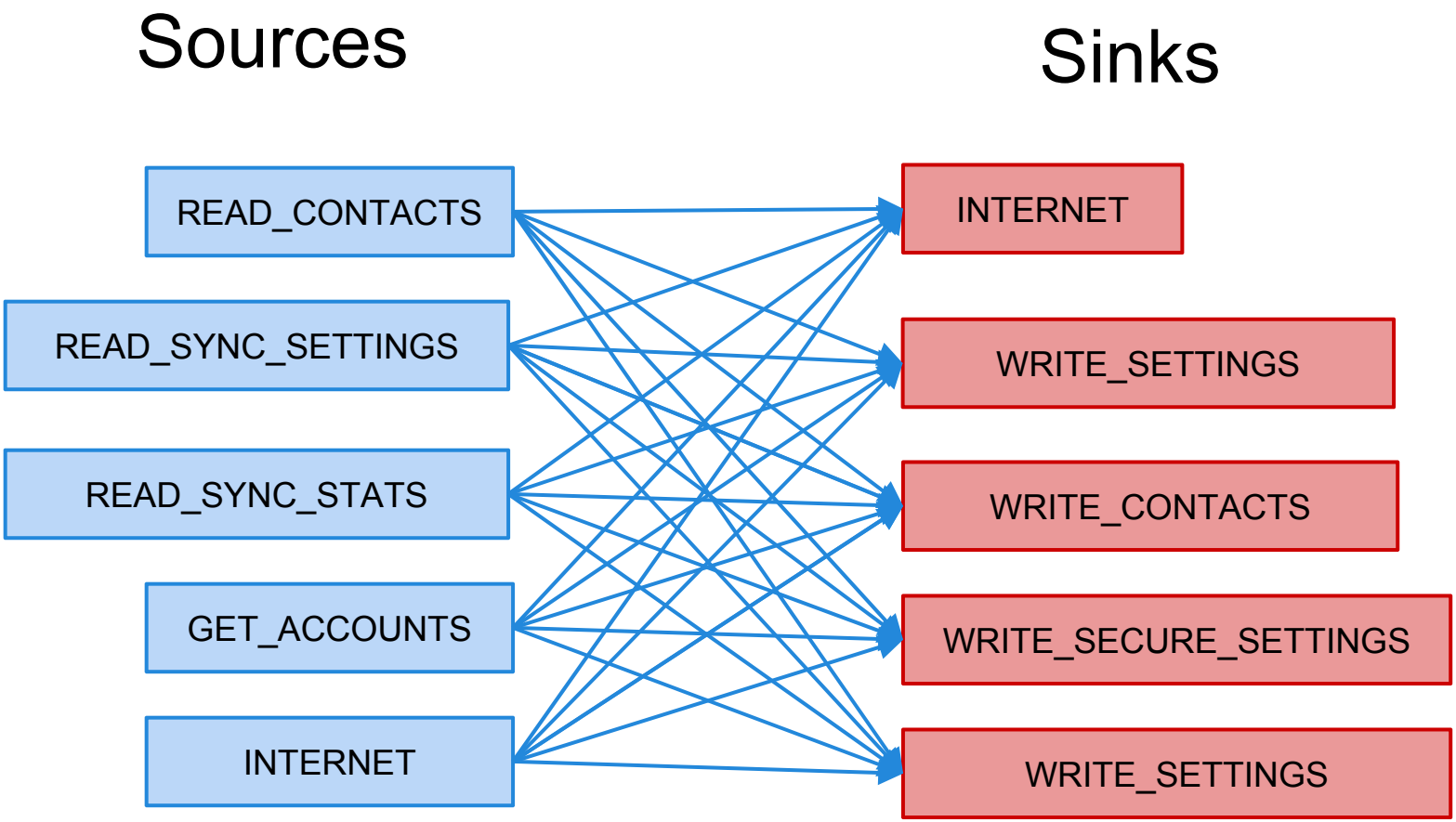
Contact Sync for Facebook (unofficial)



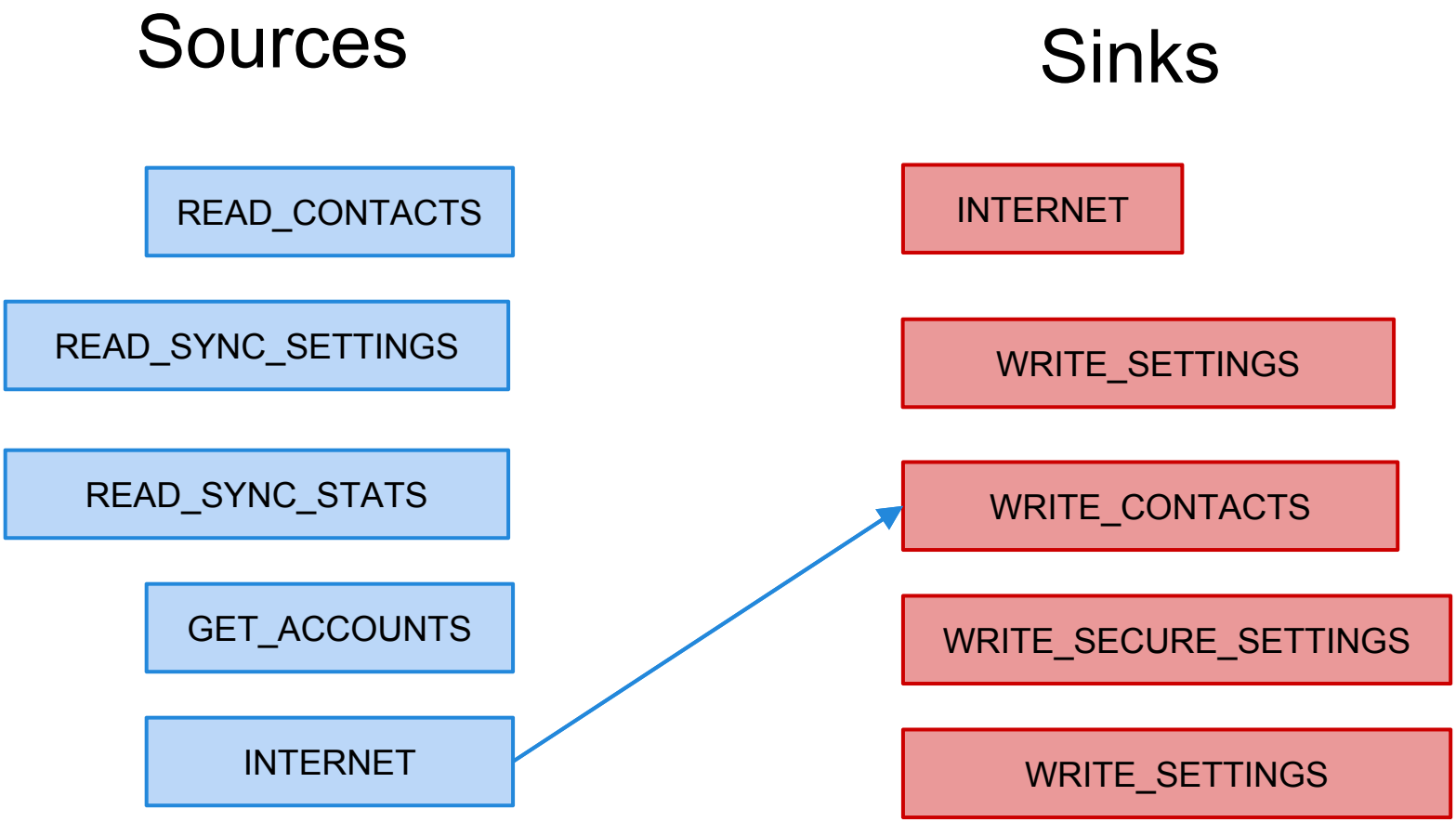
Contact Sync Permissions

Category	Permission	Description
Your Accounts	AUTHENTICATE_ACCOUNTS	Act as an account authenticator
	MANAGE_ACCOUNTS	Manage accounts list
	USE_CREDENTIALS	Use authentication credentials
Network Communication	**INTERNET**	**Full Internet access**
	ACCESS_NETWORK_STATE	View network state
Your Personal Information	**READ_CONTACTS**	**Read contact data**
	WRITE_CONTACTS	Write contact data
System Tools	WRITE_SETTINGS	Modify global system settings
	WRITE_SYNC_SETTINGS	Write sync settings (e.g. Contact sync)
	READ_SYNC_SETTINGS	Read whether sync is enabled
	READ_SYNC_STATS	Read history of syncs
Your Accounts	**GET_ACCOUNTS**	**Discover known accounts**
Extra/Custom	WRITE_SECURE_SETTINGS	Modify secure system settings

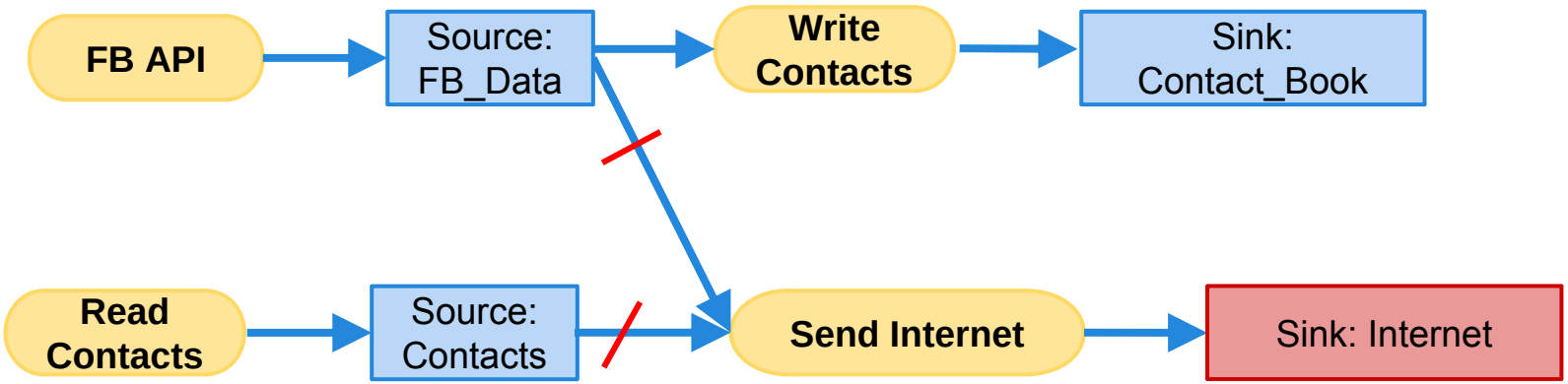
Possible Flows from Permissions



Expected Flows



Observed Flows



Example Study: Mobile Web Apps

- Goal

  Identify security concerns and vulnerabilities specific to mobile apps that access the web using an embedded browser

- Technical summary
  - WebView object renders web content
  - methods loadUrl, loadData, loadDataWithBaseUrl, postUrl
  - addJavascriptInterface(obj, name) allows JavaScript code in the web content to call Java object method name.foo()

Sample results

Analyze 998,286 free web apps from June 2014

Mobile Web App Feature	% Apps
JavaScript Enabled	97
JavaScript Bridge	36
shouldOverrideUrlLoading	94
shouldInterceptRequest	47
onReceivedSslError	27
postUrl	2
Custom URL Patterns	10

Vuln	% Relevant	% Vulnerable
Unsafe Navigation	15	34
Unsafe Retrieval	40	56
Unsafe SSL	27	29
Exposed POST	2	7
Leaky URL	10	16

Summary

- Static vs dynamic analyzers
- General properties of static analyzers
  - Fundamental limitations
  - Basic method based on abstract states
- More details on one specific method
  - Property checkers from Engler et al., Coverity
  - Sample security-related results
- Static analysis for Android malware
  - STAMP method, sample studies

Slides from: S. Bugrahe, A. Chou, I&T Dillig, D. Engler, J. Franklin, A. Aiken, …