CS 346: Cryptography

Spring 2023

## Homework 1: Symmetric Cryptography

Due: January 24, 2023 at 11:59pm (Submit on Gradescope)

Instructor: David Wu

**Instructions.** You **must** typeset your solution in LaTeX using the provided template:

https://www.cs.utexas.edu/~dwu4/courses/sp23/static/homework.tex

You must submit your problem set via Gradescope (accessible through Canvas).

**Collaboration Policy.** You may discuss your general *high-level* strategy with other students, but you may not share any written documents or code. You should not search online for solutions to these problems. If you do consult external sources, you must cite them in your submission. You must include the names of all of your collaborators with your submission. Refer to the official course policies for the full details.

**Problem 1: Compression and Encryption** [**10 points**]. Compression is a common technique used for data storage and transmission. Suppose you want to use compression in conjunction with encryption to try and get the efficiency properties of compression with the security properties of encryption. We consider two proposals here. For concreteness, you can assume the data of interest is English text (which is highly compressible).

- (a) Suppose we first compress the message before encrypting it with a semantically-secure encryption scheme. Is this approach beneficial for reducing the size of the ciphertext? Does this approach yield a semantically-secure encryption scheme (for equal-length input messages)? Give a brief justification (1-2 sentences) for each property.
- (b) Suppose we first encrypt the message using a semantically-secure encryption scheme and then apply the compression algorithm to the ciphertext. Is this approach beneficial for reducing the size of the ciphertext? Does this approach yield a semantically-secure encryption scheme (for equal-length input messages)? Give a brief justification (1-2 sentences) for each property.

**Problem 2: Understanding Advantage [24 points].** In this problem, we will get some practice with the notion of advantage. Consider the following two experiments between a challenger and an adversary A:

- Experiment 0: In this experiment, the challenger samples a bit β ← {0, 1} and gives β to the adversary. The adversary A outputs a bit b' ∈ {0, 1} which is the output of the experiment.
- Experiment 1: In this experiment, the challenger gives β = 0 to the adversary. The adversary A outputs a bit b' ∈ {0, 1} which is the output of the experiment.

For an adversary  $\mathcal{A}$  and a bit  $b \in \{0, 1\}$ , we define  $W_b = \Pr[\mathcal{A} \text{ outputs } 1 \text{ in Experiment } b]$ . We then define the advantage of  $\mathcal{A}$  to be  $\operatorname{Adv}[\mathcal{A}] = |W_0 - W_1|$ . If the advantage is 1, then that means algorithm  $\mathcal{A}$  can perfectly distinguish between the two experiments. If the advantage is 0 for all adversaries  $\mathcal{A}$ , then the two experiments are identical. If the advantage of every efficient adversary is negligible (with respect to some security parameter), then we say that the two experiments are computationally indistinguishable.

- (a) Compute the advantage for each of the following adversaries:
  - Algorithm  $\mathcal{A}$  always outputs 1.
  - Algorithm  $\mathcal{A}$  outputs  $b' \stackrel{\mathbb{R}}{\leftarrow} \{0, 1\}$ .
  - Algorithm  $\mathcal{A}$  outputs  $\beta$ .
  - Algorithm  $\mathcal{A}$  outputs  $1 \beta$ .
  - Algorithm  $\mathcal{A}$  outputs 1 if  $\beta = 1$  and  $b' \stackrel{\mathbb{R}}{\leftarrow} \{0, 1\}$  if  $\beta = 0$ .
- (b) What is the maximum possible advantage of any adversary for distinguishing between these two experiments. Give a formal proof of this.

**Problem 3: Ciphertexts and Pseudorandomness [15 points].** Let (Encrypt, Decrypt) be a symmetric encryption scheme with key-space  $\{0, 1\}^{\lambda}$ , message space  $\{0, 1\}^{n}$ , and ciphertext space  $\{0, 1\}^{t}$ . In this problem, we say that (Encrypt, Decrypt) has pseudorandom ciphertexts if for all messages  $m \in \{0, 1\}^{n}$ , no efficient adversary can distinguish between the following two distributions with non-negligible probability:

$$\{\mathsf{Encrypt}(k,m): k \xleftarrow{\mathbb{R}} \{0,1\}^{\lambda}\} \text{ and } \{\mathsf{ct} \xleftarrow{\mathbb{R}} \{0,1\}^{t}\}.$$

- (a) Give an example of a semantically-secure symmetric encryption scheme where the ciphertexts are pseudorandom. Prove both properties of your scheme.
- (b) Give an example of a semantically-secure symmetric encryption scheme where the ciphertexts are *not* pseudorandom (i.e., there exists an efficient distinguisher that breaks pseudorandomness). Prove both properties of your scheme.

In both constructions, you can specify the values of *n* and *t* (as arbitrary functions of the key-length  $\lambda$ ). You may also cite results from lecture without proof.

**Problem 4: Pseudorandom Generators [25 points].** Let  $G: \{0,1\}^{\lambda} \to \{0,1\}^{n}$  be a secure PRG. For each of the following functions G', indicate whether it is a secure PRG or not. If it is secure, give a *formal* proof. If not, describe an explicit attack.

- (a)  $G'(s) := G(s) || (G(s) \oplus 1^n)$ , where  $1^n$  denotes the all-ones string of length *n*.
- (b)  $G'(s_1 || s_2) := G(s_1) \oplus G(s_1 \oplus s_2).$
- (c)  $G'(s_1 || s_2) := s_1 || G(s_2)$ , where  $|s_1| = \text{poly}(\lambda)$ .

Recall that for two bit-strings  $s_1, s_2 \in \{0, 1\}^*$ , we write  $s_1 || s_2$  to denote the *concatenation* of  $s_1$  and  $s_2$ . Please refer to this handout for examples of how to formally show whether a construction is secure or not.

**Problem 5: Time Spent [1 point].** How long did you spend on this problem set? This is for calibration purposes, and the response you provide does not affect your score.

**Optional Feedback.** Please answer the following *optional* questions to help us design future problem sets. You do not need to answer these questions. However, we do encourage you to provide us feedback on how to improve the course experience.

- (a) What was your favorite problem on this problem set? Why?
- (b) What was your least favorite problem on this problem set? Why?
- (c) Do you have any other feedback for this problem set?
- (d) Do you have any other feedback on the course so far?