

Groth-Ostrowsky-Schae (GOS) construction:

1. "Commit" to all of the wire values in the circuit
2. Prove that each output wire is the NAND of the input wires.
3. Open the output wire to a 1 (and the input wires associated with the statement)

How to commit? Use a BGN encryption scheme!

Formally, let  $C: \{0,1\}^n \times \{0,1\}^h \rightarrow \{0,1\}$  be the circuit

1. Let  $s$  be the number of wires in the circuit, Index them topologically.
2. Let  $t_1, \dots, t_s \in \{0,1\}$  be the value of the wires in  $C(x,w)$
3. Prover commits to each wire by constructing a BGN ciphertext:
  - Sample  $r_i \xleftarrow{R} \mathbb{Z}_N$  and set  $c_i = g^{t_i} h^{r_i}$
  - For each NAND gate in the circuit (with wires  $i, j, k$ ), construct a NIZK proof that  $t_k = \text{NAND}(t_i, t_j)$  with respect to  $c_i, c_j, c_k$  and  $t_i, t_j, t_k \in \{0,1\}$ .

Proof consists of commitments  $c_1, \dots, c_s$ , NIZK proofs for each NAND gate and the openings for the statement  $(r_1, \dots, r_n)$  and for the output  $r_s$ .

To verify, check NIZK proofs all verify and that

$$c_i = g^{x_i} h^{r_i} \text{ for all } i \in [n] \text{ and } c_s = g h^{r_s}$$

Suffices to construct NIZK proof that  $t_k = \text{NAND}(t_i, t_j)$  and  $t_i, t_j, t_k \in \{0, 1\}$

Suppose  $c = g^t h^r$ . How to prove in zero-knowledge that  $t \in \{0, 1\}$  (i.e., without revealing  $t$ )?

Idea:  $t \in \{0, 1\}$  if and only if  $t(t-1) = 0$ . Use pairing to compute  $g^{t(t-1)}$ .

$$\begin{aligned} e(c, cg^{-1}) &= e(g^t h^r, g^{t-1} h^r) \\ &= \underbrace{e(g, g)^{t(t-1)}}_{\text{vanishes}} \cdot \underbrace{e(g^t, h^r) \cdot e(h^r, g^{t-1}) \cdot e(h^r, h^r)}_{e(h, (g^{2t-1} h^r)^r)} \end{aligned}$$

Proof is  $u = (g^{2t-1} h^r)^r$ .

Soundness. Suppose  $c \neq g^t h^r$  for some  $t \in \{0, 1\}$  and  $r \in \mathbb{Z}_N$ . Then,

$$\begin{aligned} e(c, cg^{-1}) &= e(g^t h^r, g^{t-1} h^r) \\ &= \underbrace{e(g, g)^{t(t-1)}}_{\substack{t(t-1) \neq 0 \text{ so} \\ \text{this component is non-zero} \\ \text{in the mod-}p \text{ subgroup of } \mathbb{G}_T}} \cdot \underbrace{e(g^t, h^r) \cdot e(g^{t-1}, h^r) \cdot e(h^r, h^r)}_{\text{mod-}g \text{ subgroup of } \mathbb{G}} \end{aligned}$$

Thus, there does not exist  $u \in \mathbb{G}$  such that  $e(c, cg^{-1}) = \underbrace{e(h, u)}_{\text{zero in mod-}p \text{ subgroup}}$ .

Zero-knowledge: Proof is deterministic. To prove zero-knowledge, need to randomize (to hide values of  $t$  and  $r$ ).

Prover picks  $\alpha \xleftarrow{R} \mathbb{Z}_N^*$  Then,

$$e(h, u) = e(h^\alpha, u^{\alpha^{-1}})$$

Instead of giving out  $u$ , give out  $\pi_1 = h^\alpha$  and  $\pi_2 = u^{\alpha^{-1}}$

Check that  $e(c, cg^{-1}) \stackrel{?}{=} e(\pi_1, \pi_2) = e(h^\alpha, u^{\alpha^{-1}}) = e(h, u)$ .

Also give out  $\pi_3 = g^\alpha$  and have verifier check that

$$e(g, \pi_1) \stackrel{?}{=} e(\pi_3, h) \quad \text{[necessary for soundness]}$$

Correctness:  $e(g, \pi_1) = e(g, h^\alpha) = e(g^\alpha, h) = e(\pi_3, h)$ .

Randomization is sufficient to prove zero-knowledge.

Proving NAND relation:  $g^{t_1} h^{r_1} \quad g^{t_2} h^{r_2} \quad g^{t_3} h^{r_3}$

We can show that  $t_1, t_2, t_3 \in \{0, 1\}$ . Suffices to now show that  $t_3 = \text{NAND}(t_1, t_2)$ .

When  $t_1, t_2, t_3 \in \{0, 1\}$ , this holds if and only if  $t_1 + t_2 - 2t_3 + 2 \in \{0, 1\}$

[Can just check 8 possibilities for  $t_1, t_2, t_3$ ]

Can now use homomorphisms of BGN to prove this:

$$\begin{aligned} c_1 &= g^{t_1} h^{r_1} \\ c_2 &= g^{t_2} h^{r_2} \\ c_3 &= g^{t_3} h^{r_3} \end{aligned} \Rightarrow c_1 \cdot c_2 \cdot c_3^{-2} \cdot g^2 = \underbrace{g^{t_1+t_2-2t_3+2} h^{r_1+r_2-2r_3+2}}$$

prove this is commitment to 0/1 value