Privacy-Preserving Payment Splitting

Saba Eskandarian

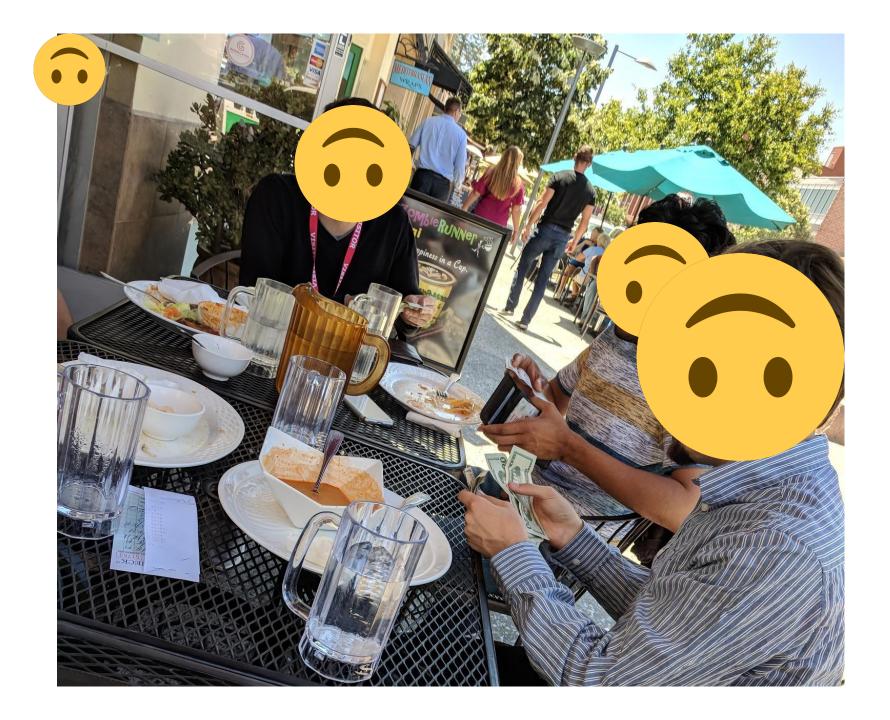
Stanford University

Mihai Christodorescu

Visa Research

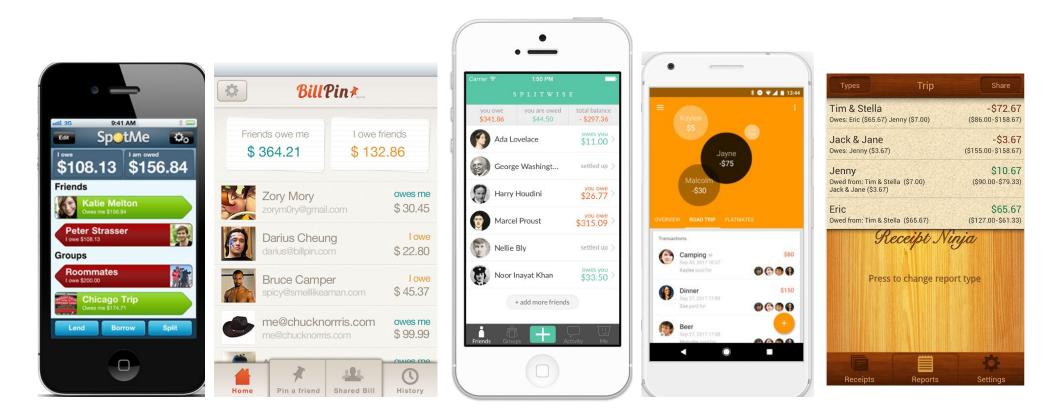
Payman Mohassel

Facebook



Splitwise, Receipt Ninja, Billpin, SpotMe, Conmigo, Settle Up, ...

Convenient way to keep track of costs and debts between groups of friends or colleagues



Privacy Policy – Data We Collect:

"This data includes, for example, group names, expense descriptions and amounts, payments and their confirmation numbers, comments and reminders, receipt images, notes, and memos, in addition to any other information that you attach or share while using ..."

"... the types of expenses you add, the features you use, the actions you take, and the time, frequency and duration of your activities"

owes me

\$ 30.45

\$ 22.80

\$45.37

owes me

\$ 99.99 owes me

History

lowe

George Washingt..

Harry Houdini

Marcel Proust

Noor Inayat Khan

+ add more friends

Nellie Bly

settled up

\$26.77

you owe \$315.09

settled up

\$33.50

\$132.86

nds or colleagues



-\$75

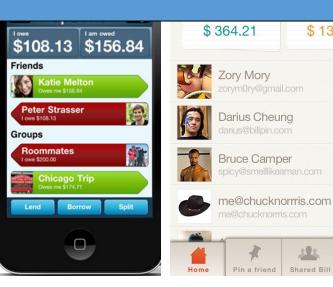
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\$150

Camping

Dinner



Privacy Policy – Data We Collect:

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zip code, e-mail address, cell phone number, occupation, hometown, "... the types of expenses you add, the features time, frequency and duration of your activities" about personal finances"



"... we may collect and process information about your actual location, like GPS signals sent by a mobile device. We may also use various technologies to determine location, such as sensor data ..."

"We also use this information to offer you tailored content – like giving you more relevant search results and ads."

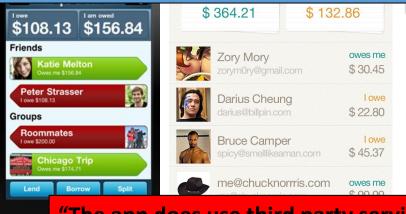




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"We also use this information to offer you tailored content – like giving you more relevant search results and ads."

"The app does use third party services that may collect information used to identify you."

Goal: cash-like privacy for payment splitting

Generic Solutions

Homomorphic encryption based solutions [e.g. Gen09, BGV11, GSW13]

Server-aided MPC solutions

[e.g. FKN94, KMR11/12, HLP11]

Zero-Knowledge Log Server

[e.g. zkLedger (NVV'18)]

Metadata-hiding anonymous group messaging?

[e.g. Riposte, Vuvuzela, Stadium, Pung, Atom]

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Goal 2: Strong performance and scalability

Our Solution

Same functionality as today's payment splitting apps

Hides user data from provider

Runs very fast: <50ms/round on phone <300µs/round on server (for realistic group sizes)

Consists mainly of AES and addition

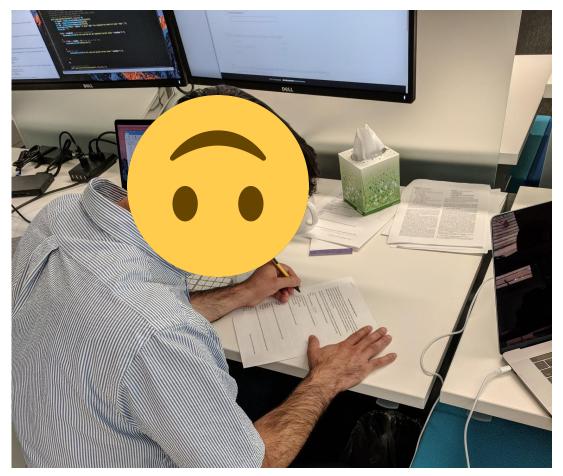
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Silent Splitter - Group Page	
Friends	
Balance: \$11	
NEW CHARGE	SETTLE
Recent Transactions You charged Bob \$2. Carol charged you \$1.	
REJECT CHARGE	
You charged Bob \$6. You charged Carol \$4.	

Informal User Survey

Sent to ~250 employees in Visa Palo Alto office, got 51 responses

Some takeaways:

- Groups tend to be small
- Groups have only a few transactions a day
- Transaction amounts are usually fairly small amounts of money



(Dramatization, it was an online survey)

Group members connect to server via app







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Group members share secret key during setup

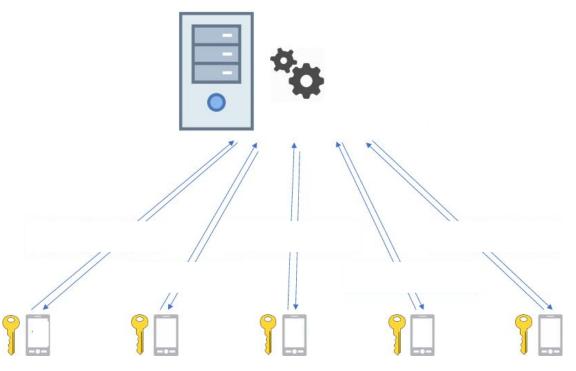




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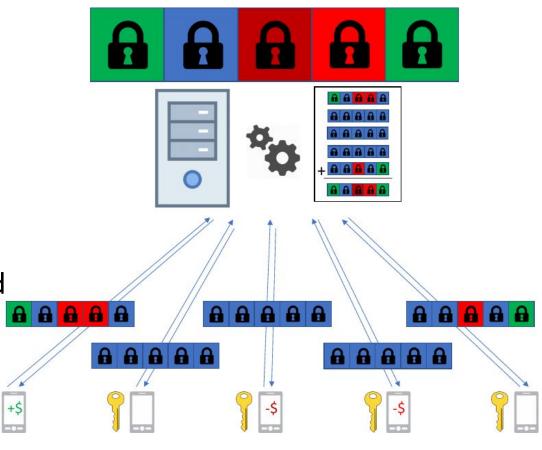
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Server *blindly* sums values and sends results (New balance, charger identity, integrity check)



Security Properties

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 No user can create or destroy money (assume >0 honest users)
No user can undetectably frame an honest user for making a charge Server Integrity: Malicious server can only cause denial of service

Security Properties

<u>Server Privacy</u>: any two sets of transactions indistinguishable to server <u>Debtor Privacy</u>: transaction hides who it puts into debt to others <u>User Integrity</u>:

1) No user can create or destroy money (assume >0 honest users)

 No user can undetectably frame an honest user for making a charge <u>Server Integrity</u>: Malicious server can only cause denial of service <u>Limitations</u>:

We do not hide group membership from the server

We do not protect against collusion between a malicious user and server

Example: Alice requests \$1 from Bob in their friend group 👧 🚳 🍲







Example: Alice requests \$1 from Bob in their friend group 👧 🚳

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Alice sets her vector to all 0s except a 1 in Bob's position



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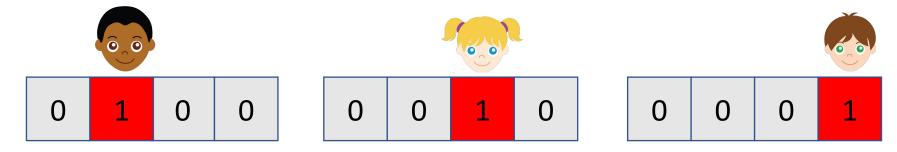


Alice sets her vector to all 0s except a 1 in Bob's position

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Anyone not making a charge puts a 1 in their own position



Example: Alice requests \$1 from Bob in their friend group

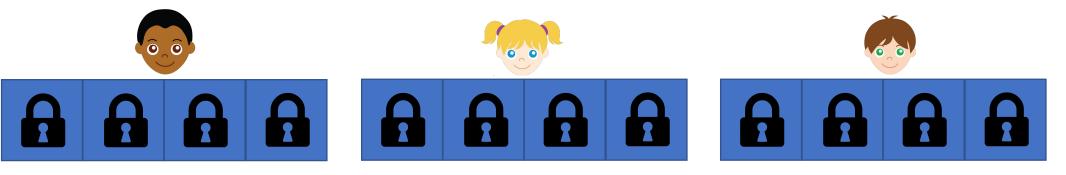
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Each user encrypts his/her vector and sends the result to the server



Anyone not making a charge puts a 1 in their own position



Faces from sweetclipart.com

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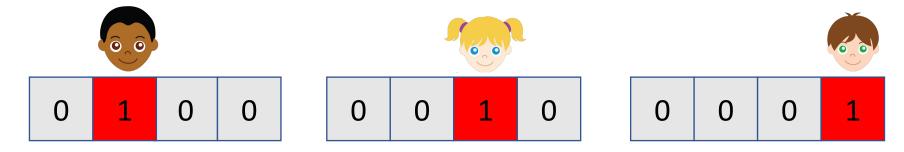
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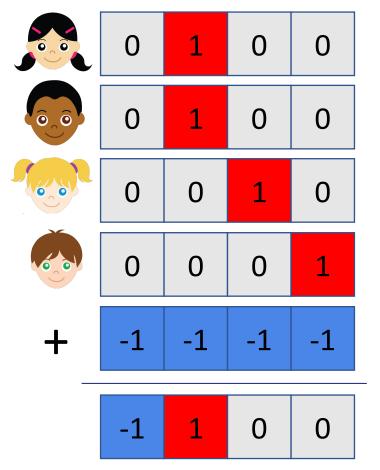
We'll start by showing the protocol without encryption



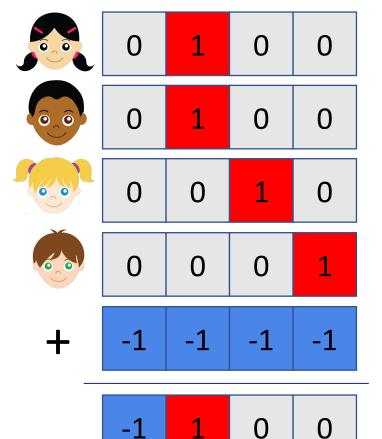
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The server adds up everyone's values and subtracts 1



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The result is added to users' existing balances



Note: server tracks *debt*, so negative is less debt

How does Bob know it was Alice who charged him?

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For each user, server takes "(input in user's own position) -1"

$$\int 0 \rightarrow -1 \quad \int 0 \quad 1 \rightarrow 0 \quad \int 0 \quad 0 \quad 1 \rightarrow 0$$

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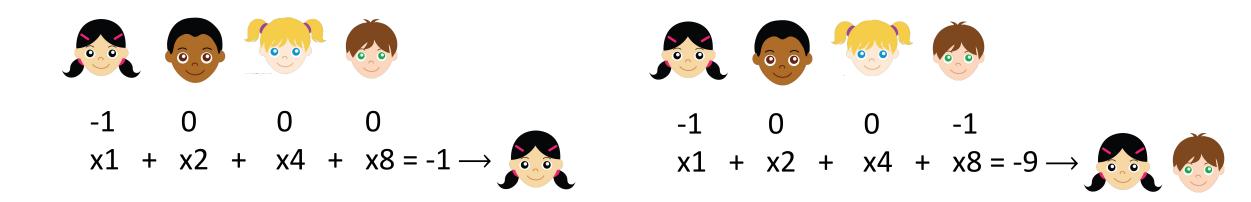
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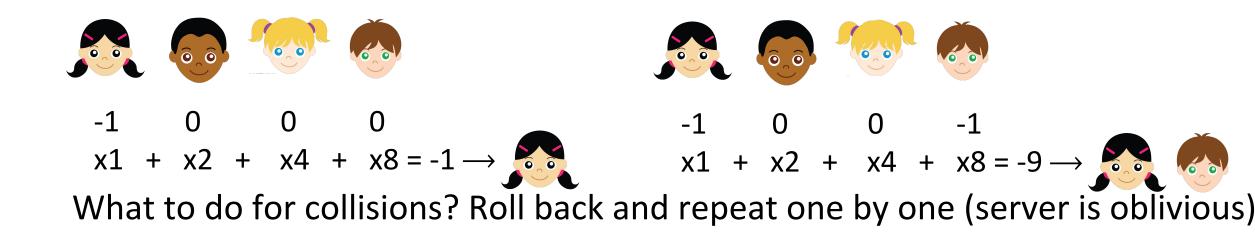
x1 = -1 + x2 = 0 + x4 = 0 + x8 = 0And sums up the results to identify the charger(s)

$$= -1 \rightarrow$$

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Adding Server Privacy

<u>Observation 1:</u> server does the same *fixed* set of additions every round.

No data-dependent operations

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No data-dependent operations

Observation 2: all the clients share a secret key k.

They can independently generate the same PRF outputs

Adding Server Privacy

Solution:

Instead of actually encrypting, users mask values with a PRF output

They send $v_i + r_i$, where $r_i = PRF(k, \text{group}, \text{user}, \text{round}, i)$

Users calculate sum of masks and remove them from server responses Calculating/removing masks fast because it's just AES and addition

Extensions

- Integrity
- Larger transactions
- Multiple charges per Round
- Identifying misbehaving users
- Handling framing
- Handling users going offline
- Improving usability for charge requests
- Integration with payment systems
- Payment splitting with collateral

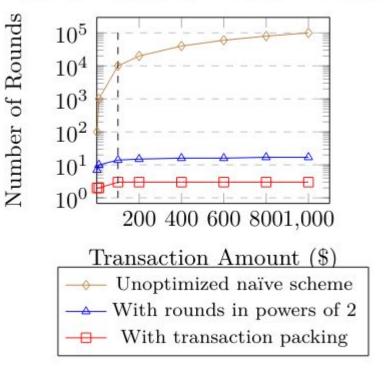
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Rounds Needed to Process Transactions



See paper for details!

Performance

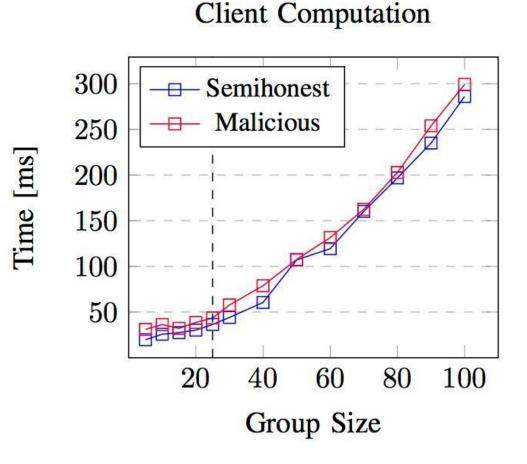
Client Performance

<50ms/round for realistic groups (realistic based on user survey)

Malicious server overhead <20ms

Only computes AES and addition

Client bandwidth for group size: 10 (≥69% of groups in survey): 160 Bytes 25 (≥92% of groups in survey): 400 Bytes 100 (≥100% of groups in survey): 1.6Kb



Server Performance

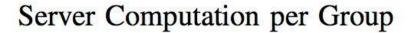
<300 *microseconds* for realistic groups (realistic based on user survey)

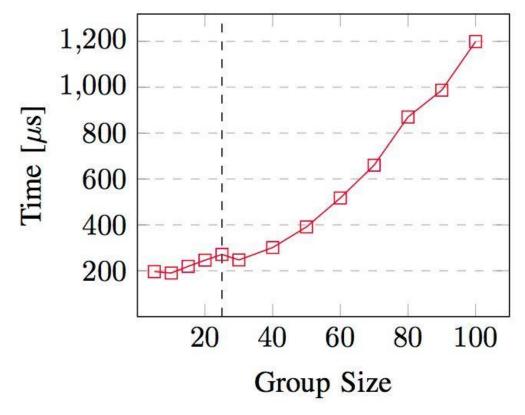
No changes for malicious security

Only computes addition

Server memory requirements small – can handle user inputs as they arrive, no need to keep in memory

See paper for more evaluation details





Summary

Our system allows payment-splitting groups to hide

- Who pays,
- Who is paid,
- How much is spent,
- When transactions are made,
- And more

From a potentially malicious server at minimal performance cost

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